

ON 4-CROSSING-FAMILIES IN POINT SETS AND AN ASYMPTOTIC UPPER BOUND

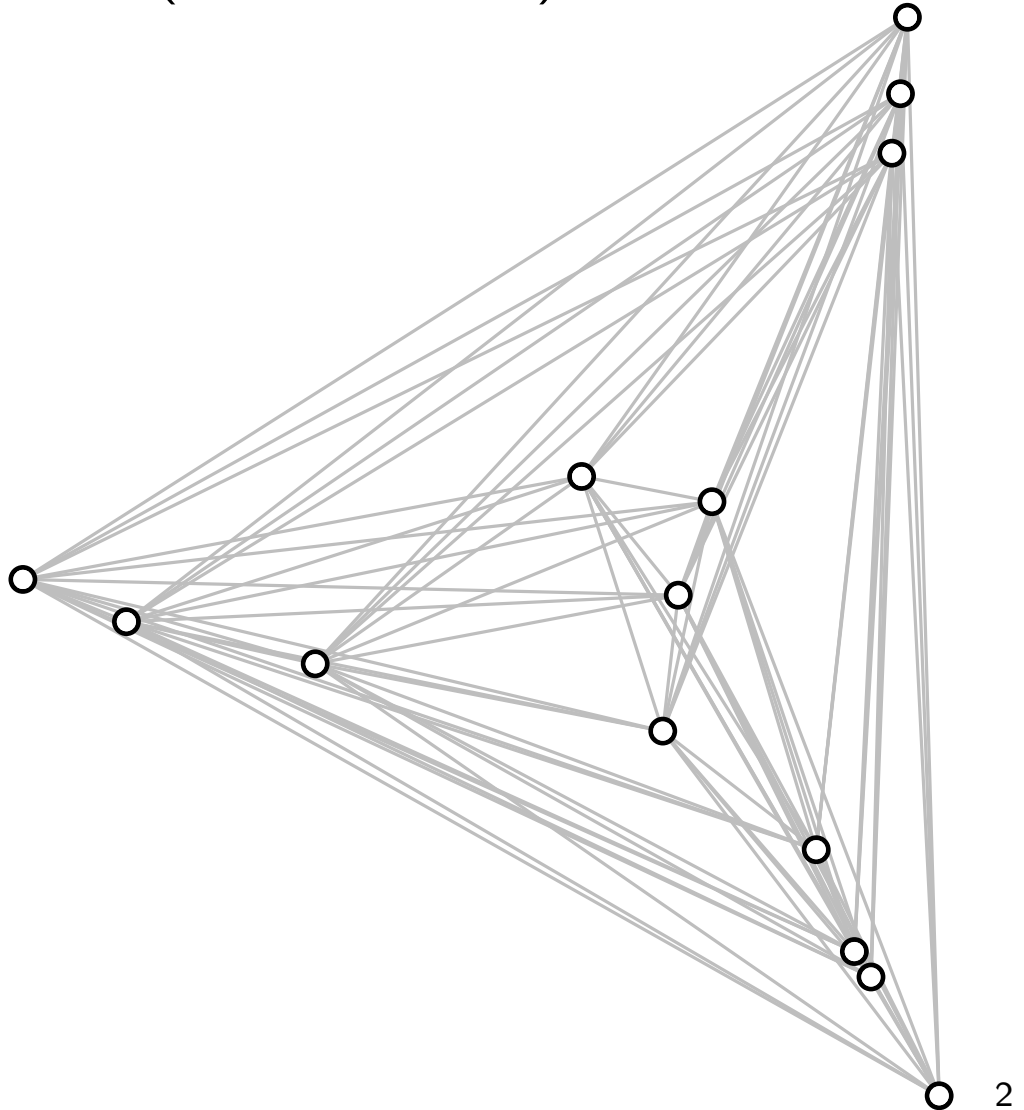
Oswin Aichholzer, Jan Kynčl,

Manfred Scheucher, and Birgit Vogtenhuber



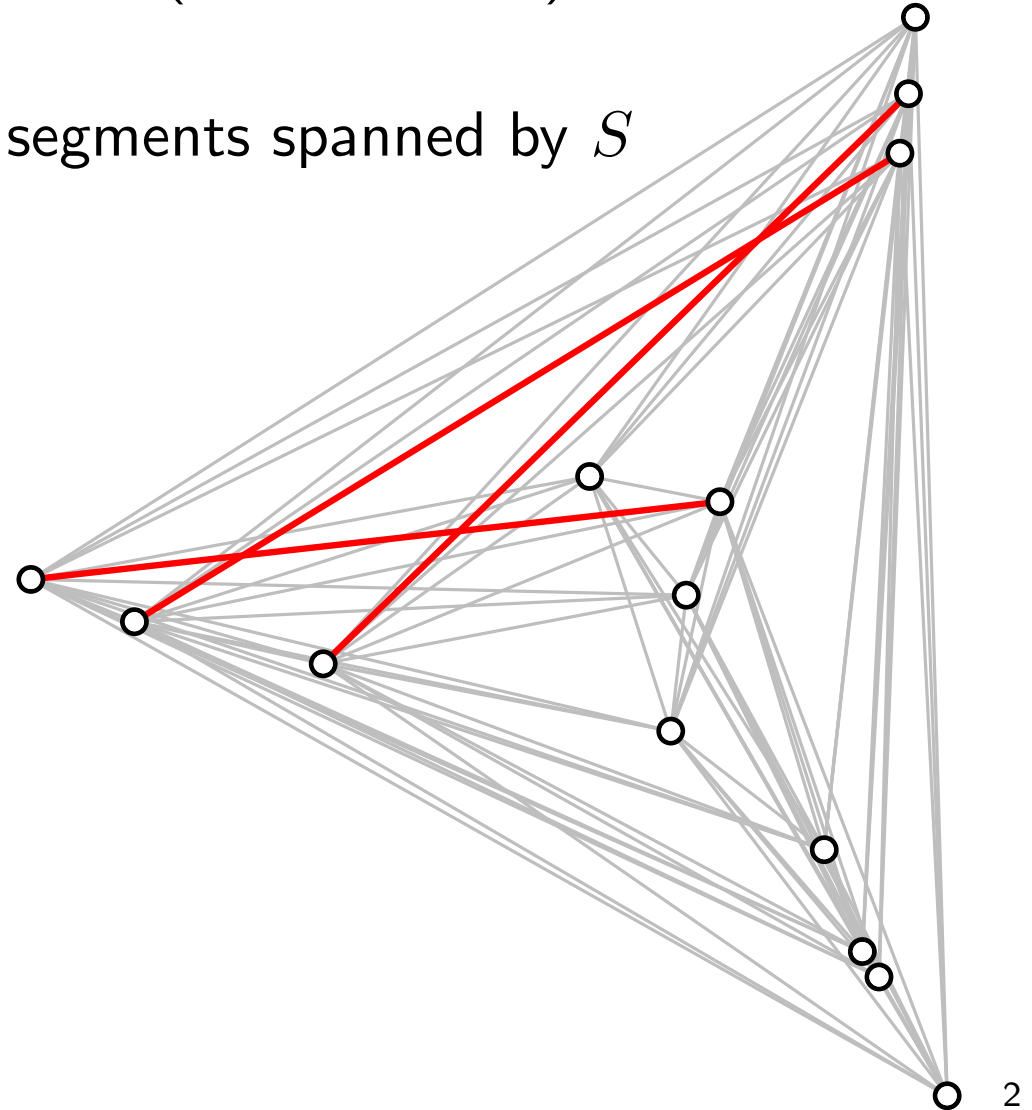
Crossing Families

- S set of n points in general position (no 3 collinear)



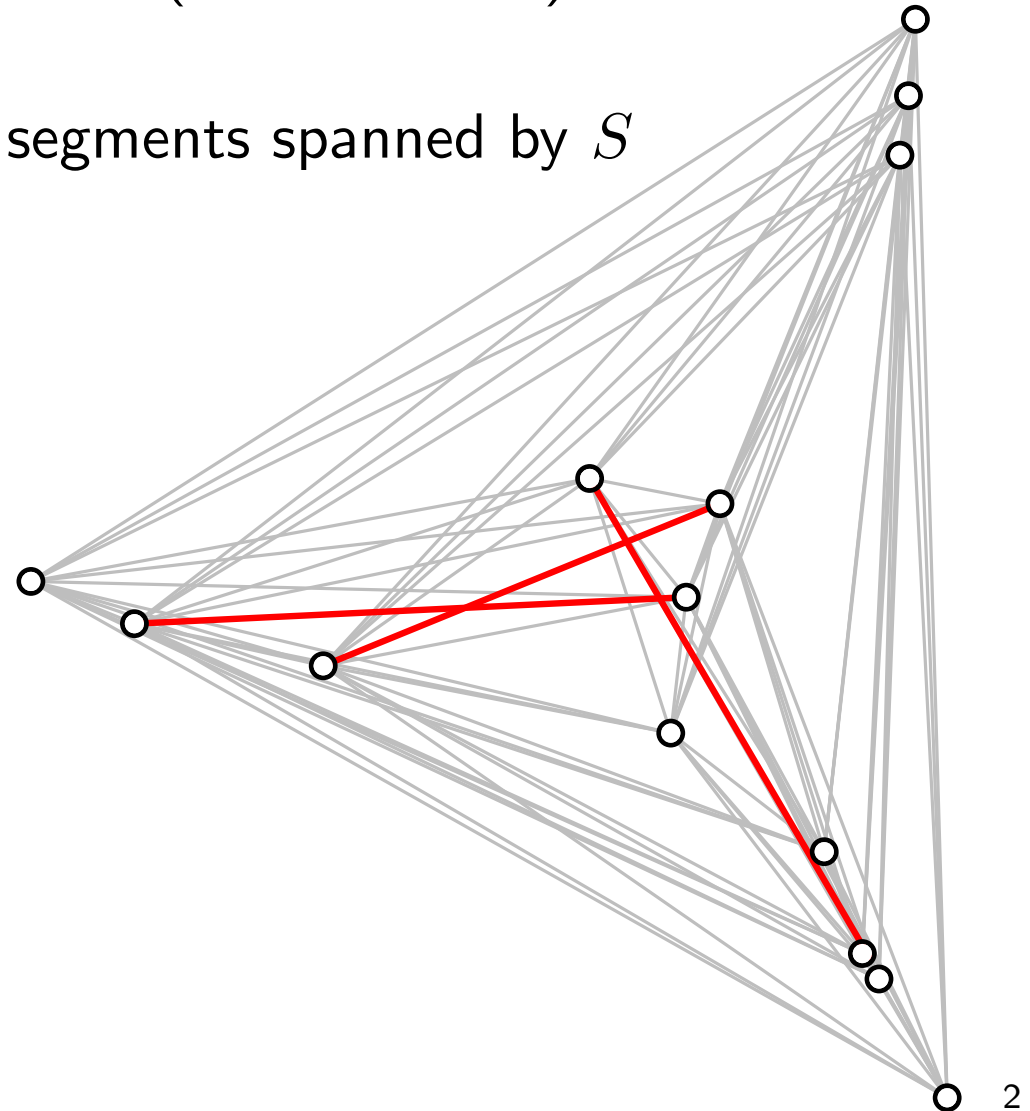
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- k -family = k pairwise crossing segments spanned by S



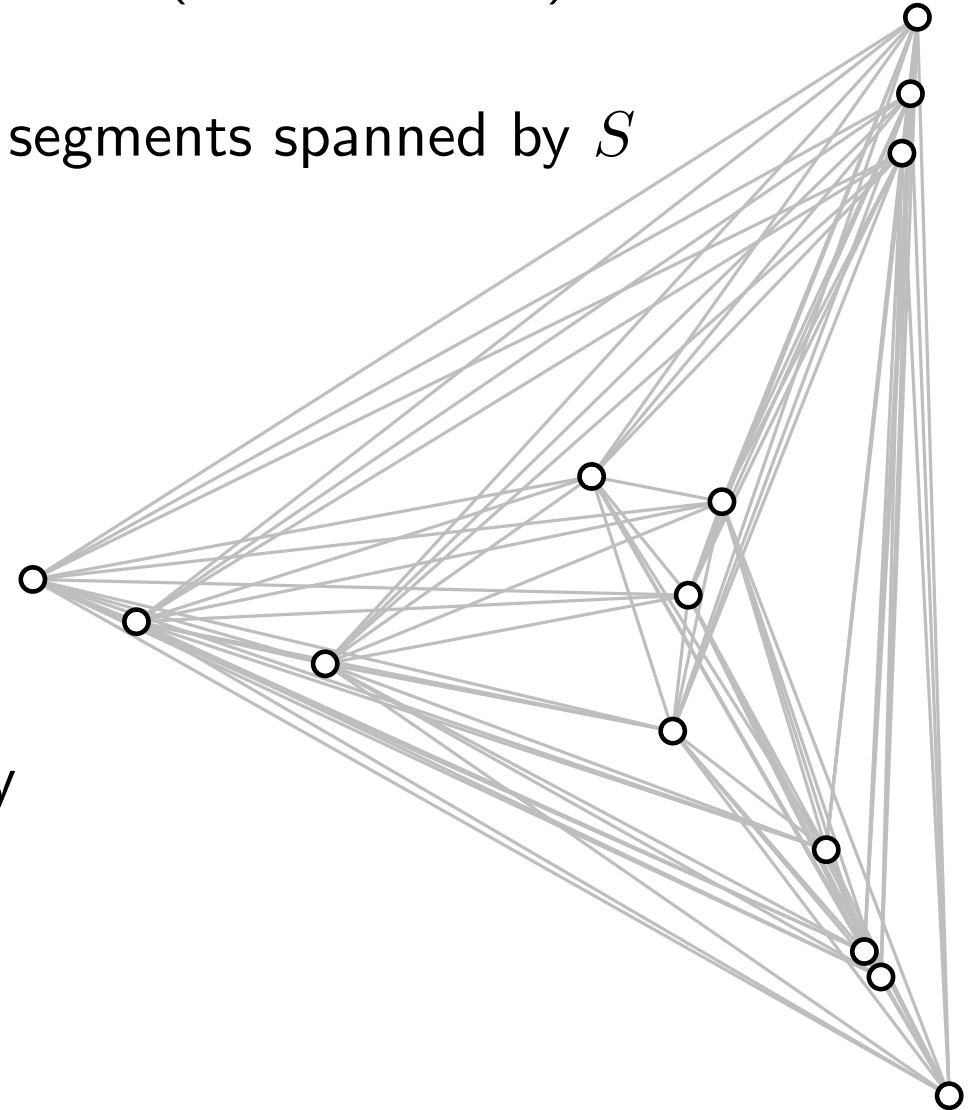
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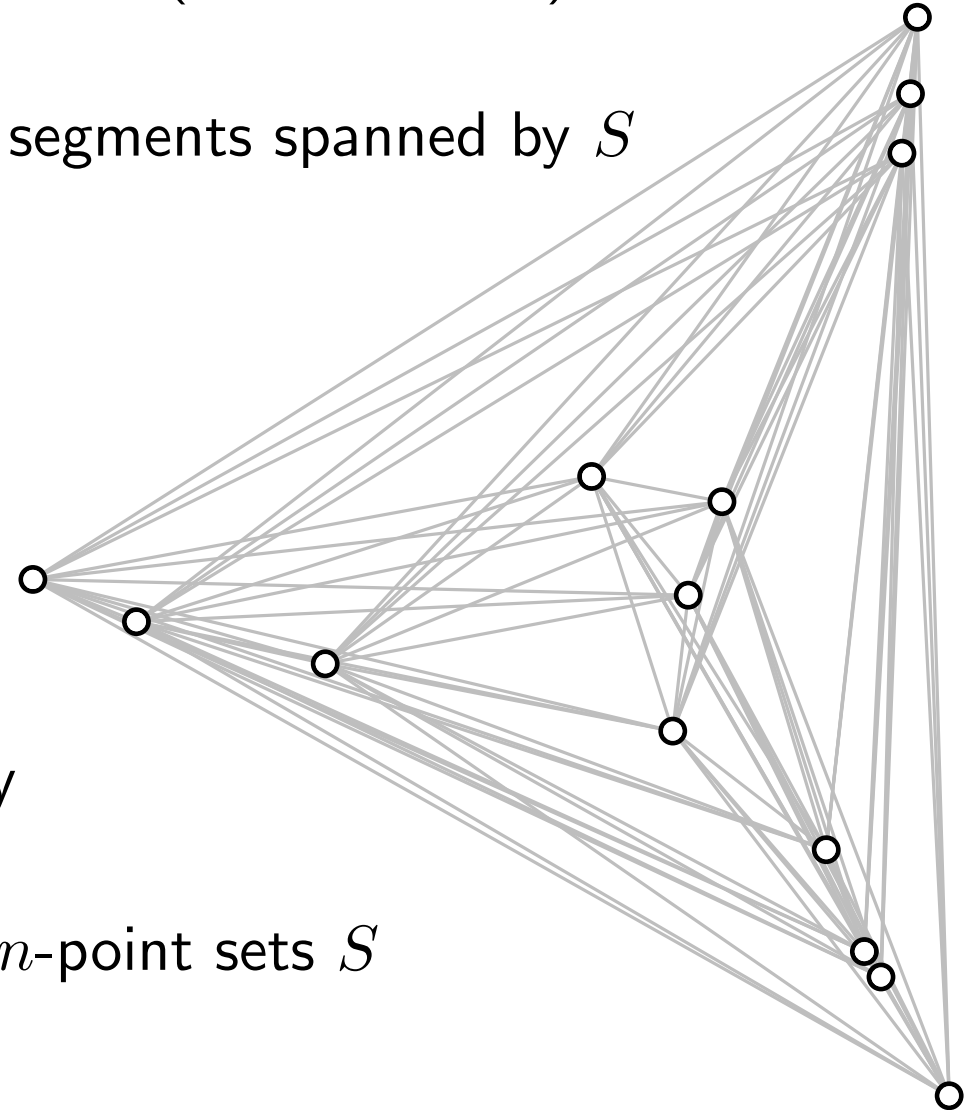
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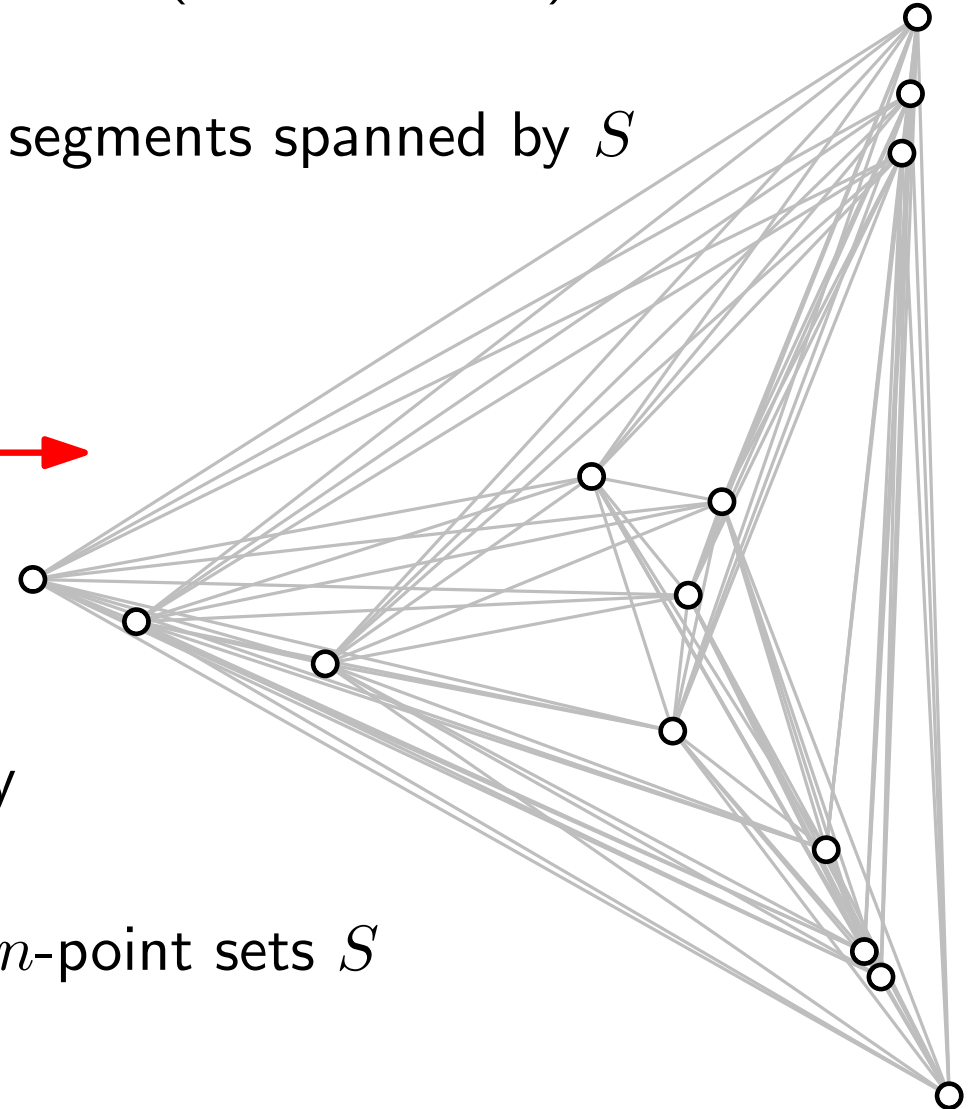


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- $cf(n)$ = min. $cf(S)$ among all n -point sets S

Crossing Families

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this set shows $cf(10) \leq 3$,
which is actually tight



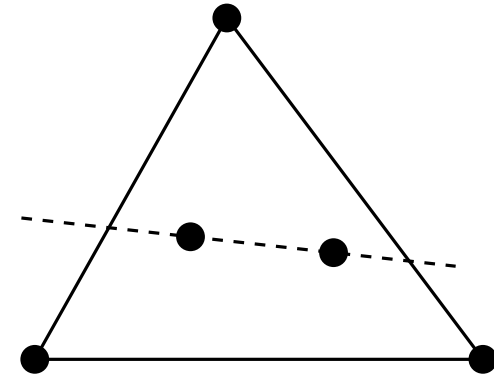
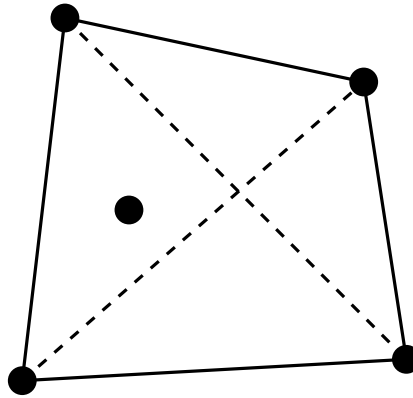
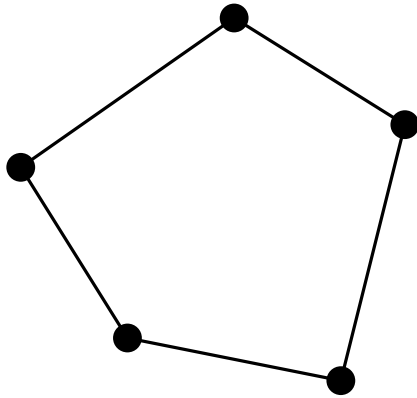
- $cf(S)$ = size of largest k -family
- $cf(n)$ = min. $cf(S)$ among all n -point sets S

A small example

- K_5 non-planar
 - \Rightarrow any set of 5 points admits a pair of crossing segments
 - $\Rightarrow \text{cf}(5) \geq 2$

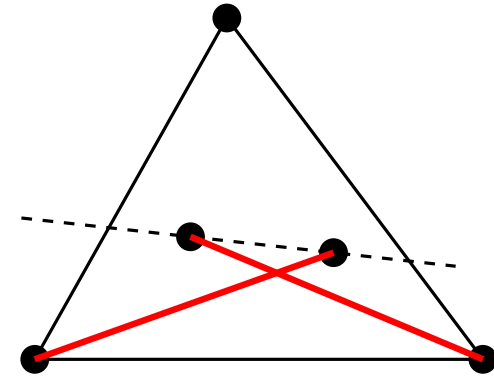
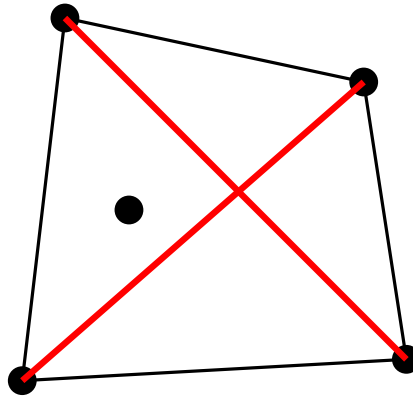
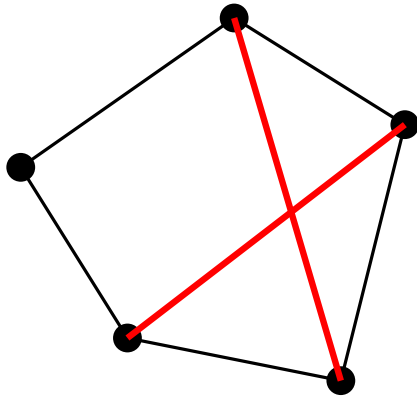
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Lower Bounds

- $\text{cf}(n) \geq \Omega(\sqrt{n})$ [Aronov, Erdős, Goddard, Kleitman, Klugerman, Pach, and Schulman '94]

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- $\text{cf}(n) \geq n^{1-o(1)}$ [Pach, Rubin, and Tardos '19]

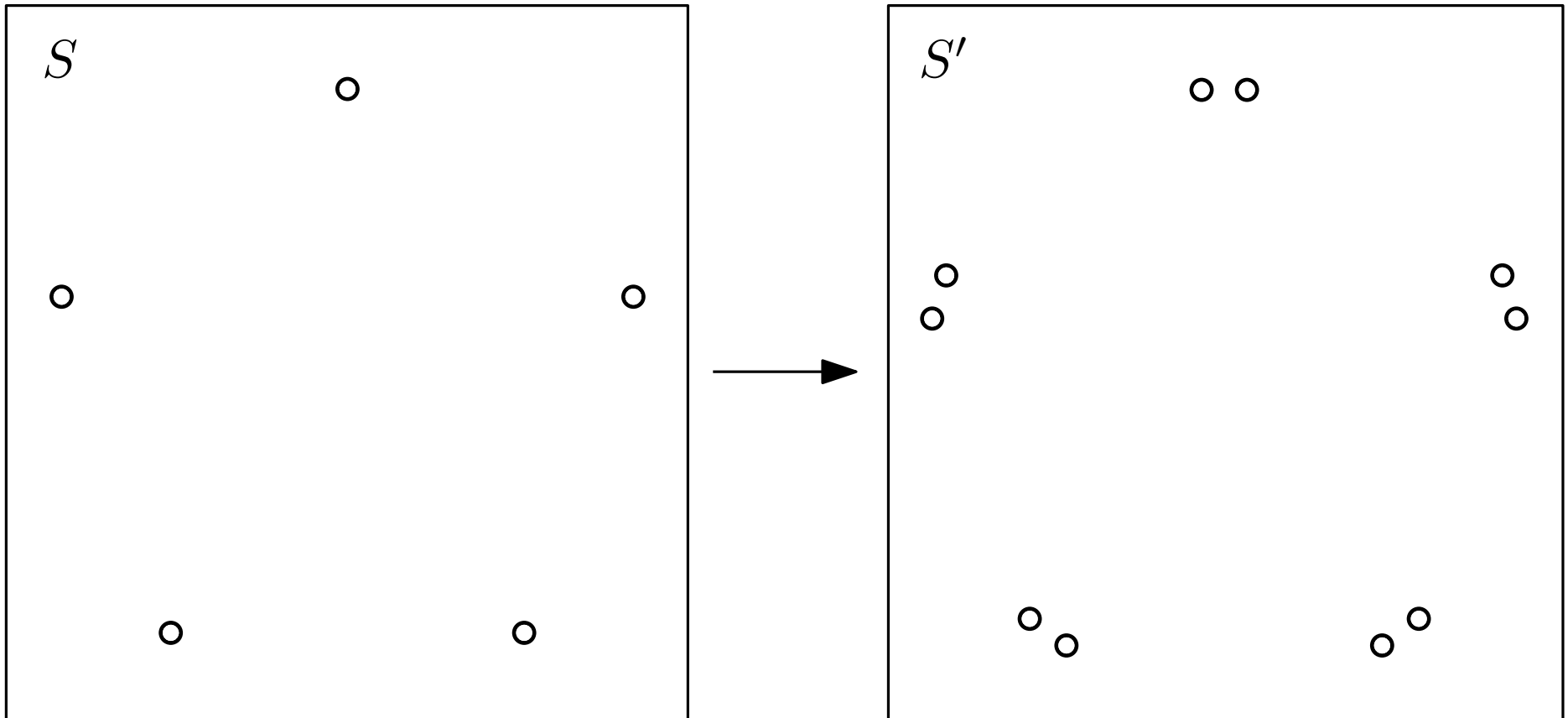
Upper Bounds

- **Theorem (Aronov, Erdős, Goddard, Kleitman, Klugerman, Pach, and Schulman '94):**

$$\text{cf}(n) \leq \text{cf}(n_0) \cdot \left\lceil \frac{n}{n_0} \right\rceil = \underbrace{\frac{\text{cf}(n_0)}{n_0}}_{\text{ratio}} \cdot n + O(1)$$

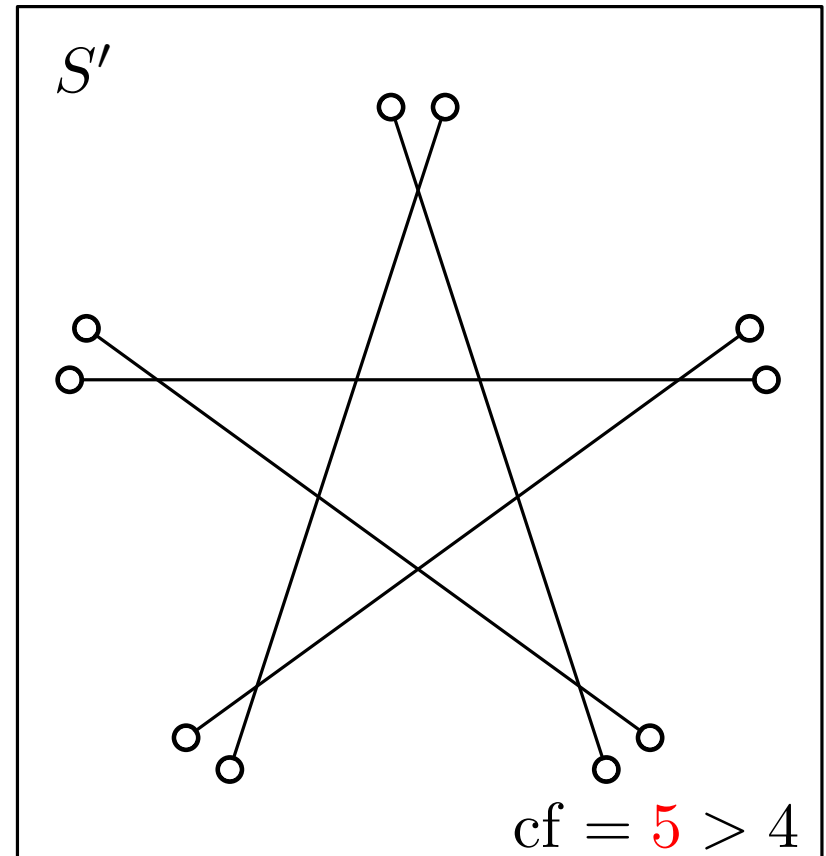
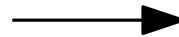
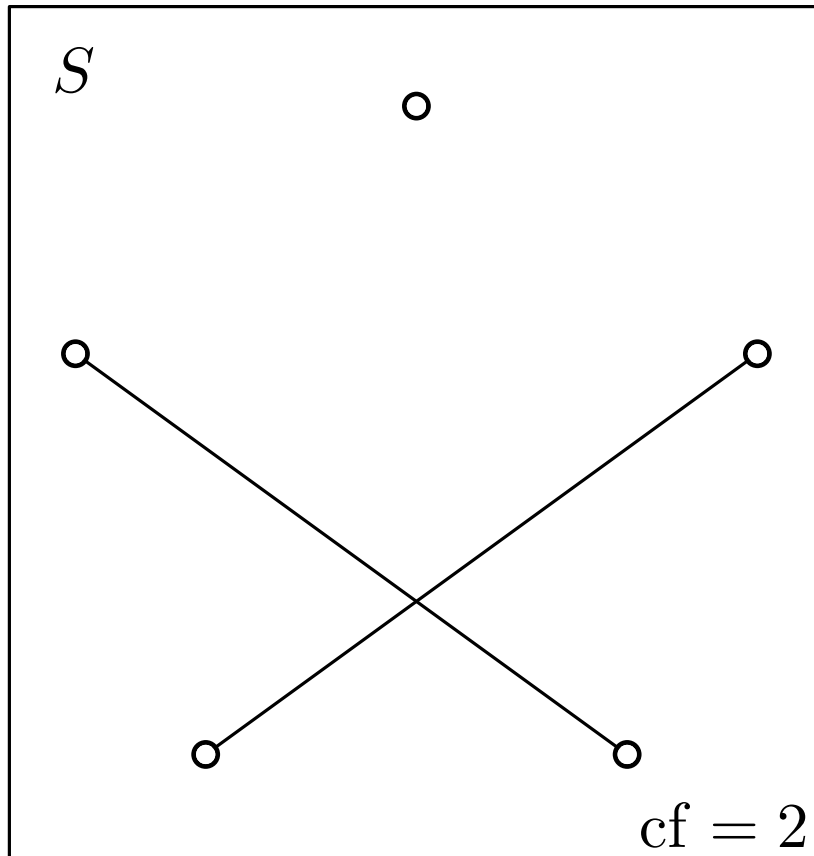
Proof Idea

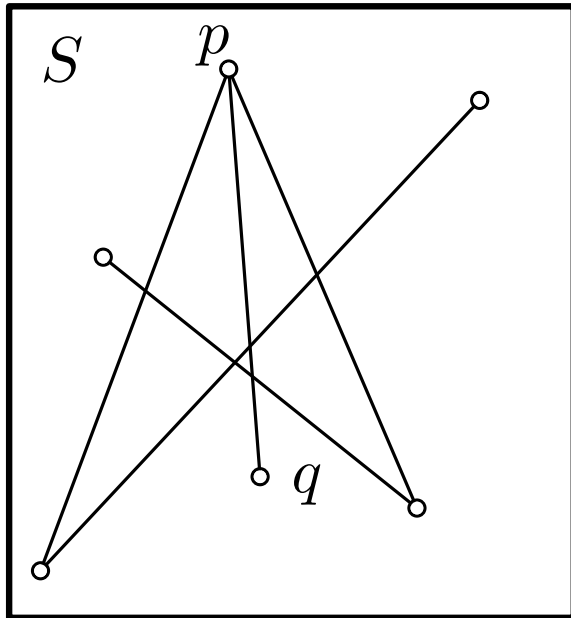
- replacing each point by k perturbed copies



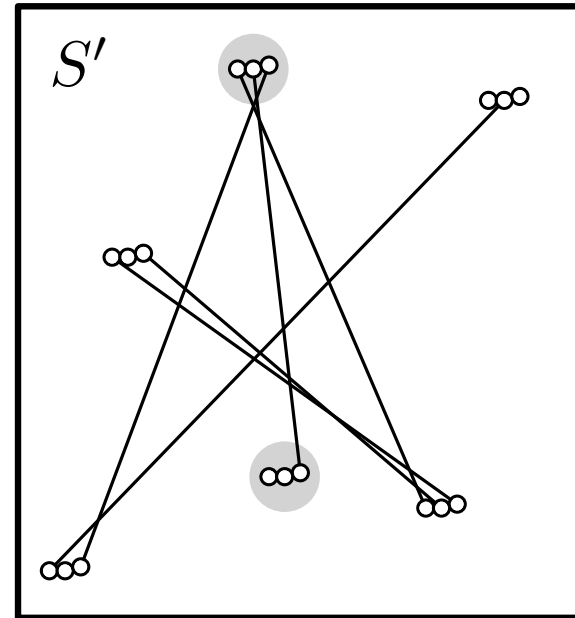
Proof Idea

- replacing each point by k perturbed copies
- need to avoid generating odd cycles of pairwise crossing edges!

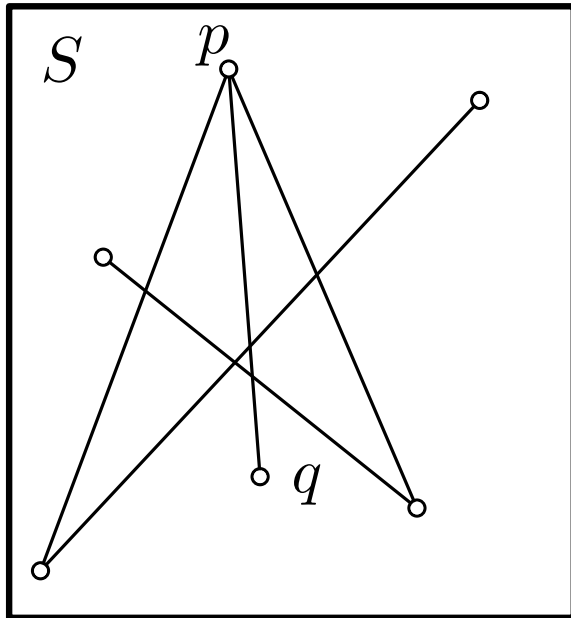




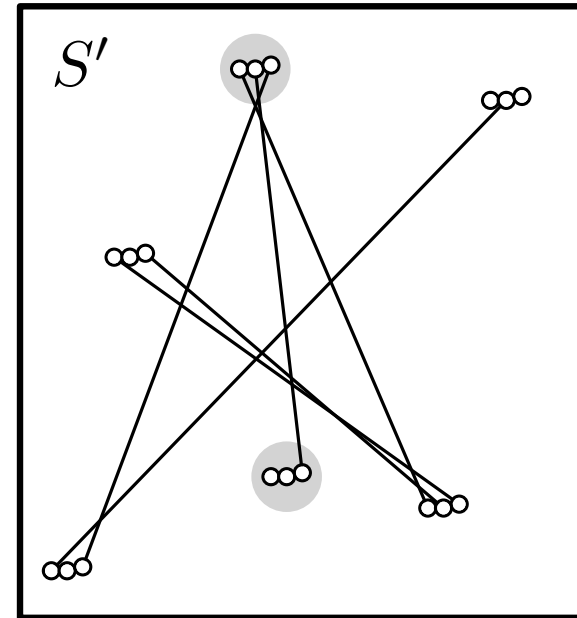
contract F'



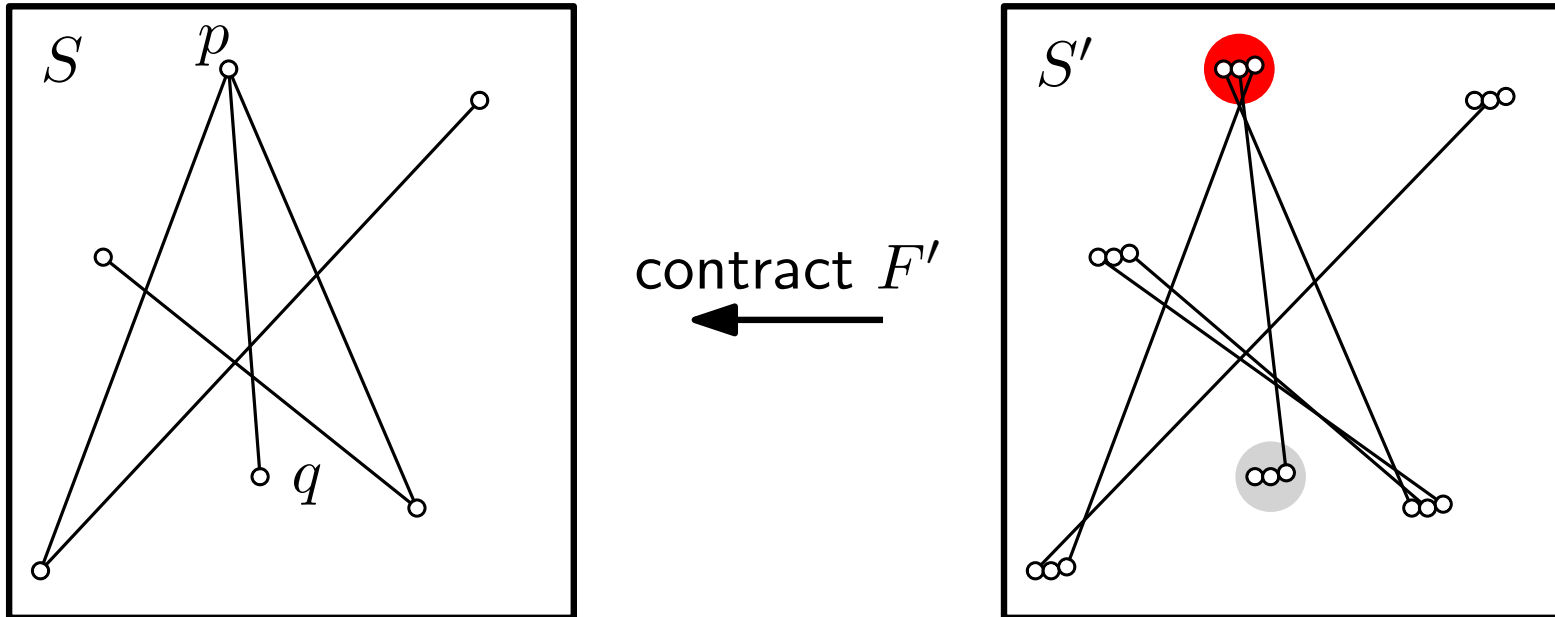
- contraction is a geometric thrackle \Rightarrow no even cycles



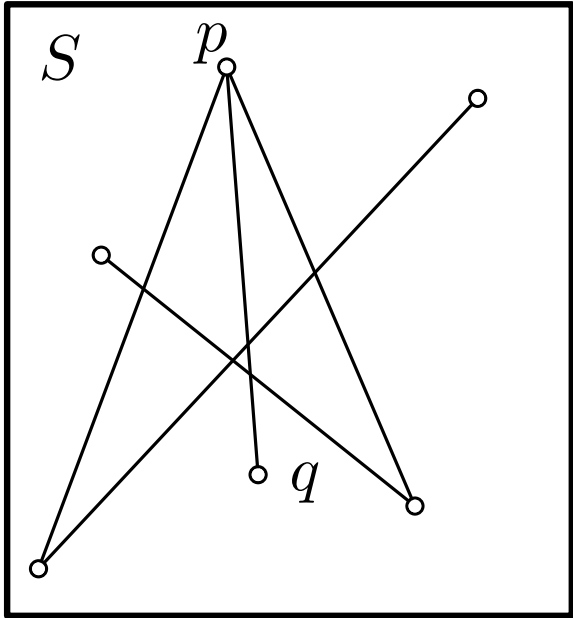
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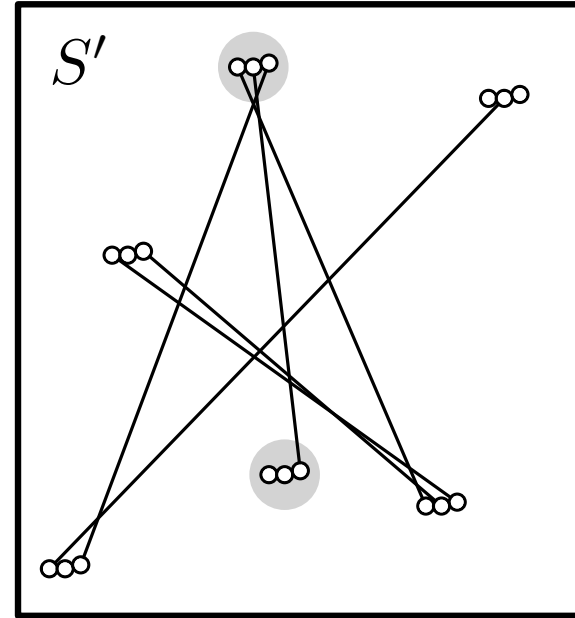
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- place all copies of a point almost on horizontal line
 \Rightarrow neighbors either all above or all below
 \Rightarrow bipartite (no odd cycles)



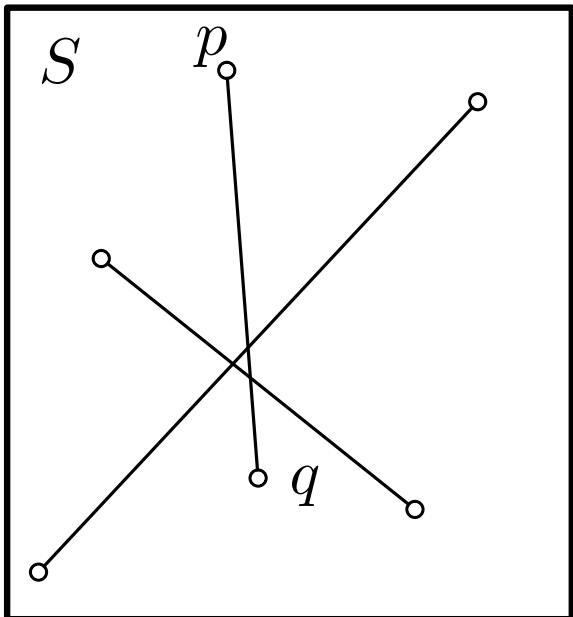
- contraction is a geometric thrackle \Rightarrow no even cycles
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- while $\deg(p) > 1$, replace edges incident to p by **one** full bundle



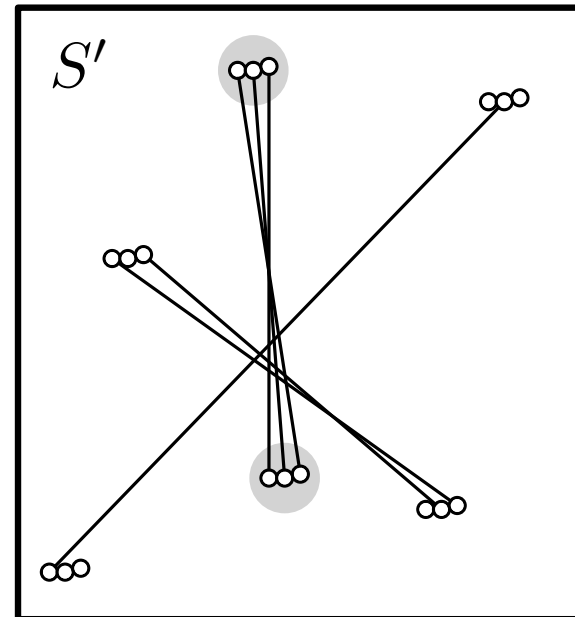
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modify F' to F''



contract F''



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- **Theorem (Aronov, Erdős, Goddard, Kleitman, Klugerman, Pach, and Schulman '94):**

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- $\text{cf}(9) = 2 \Rightarrow \text{cf}(n) \leq \frac{2}{9}n + O(1)$ [Aichholzer & Krasser '06]
- $\text{cf}(24) \leq 5 \Rightarrow \text{cf}(n) \leq \frac{5}{24}n + O(1)$ [Evans & Saeedi '19]

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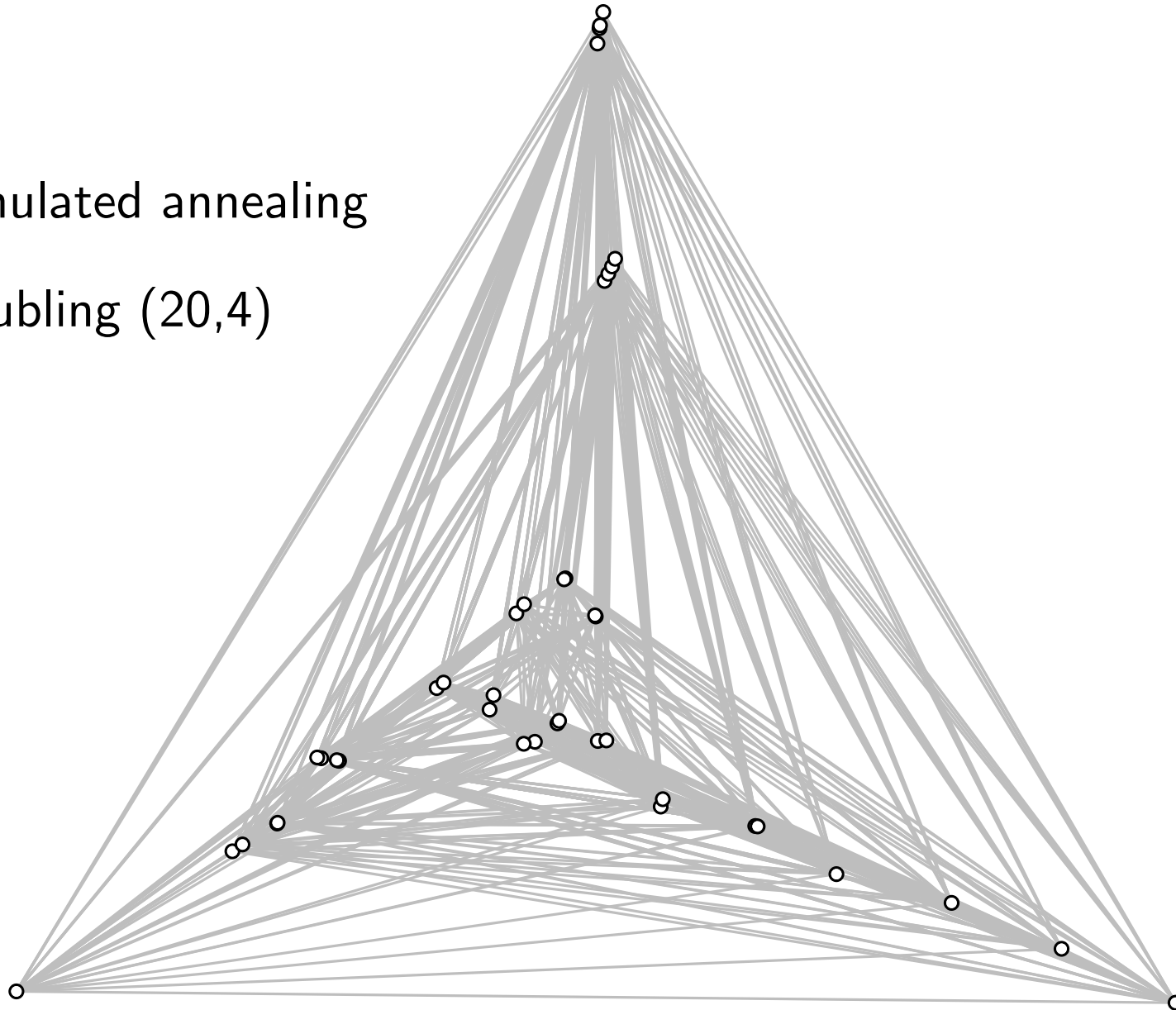
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- **NEW:** $\text{cf}(20) \leq 4 \Rightarrow \text{cf}(n) \leq \frac{1}{5}n + O(1)$
- **BRAND NEW:** $\text{cf}(41) \leq 8 \Rightarrow \text{cf}(n) \leq \frac{8}{41}n + O(1)$

41 points with $cf = 8$

- simulated annealing
- doubling (20,4)



Exact Bounds

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- $\text{cf}(4) = 1$, $\text{cf}(5) = 2$ (K_5 non-planar)
- $\text{cf}(9) = 2$, $\text{cf}(10) = 3$ [Aichholzer & Krasser '06]
- **NEW:** $\text{cf}(14) = 3$, $\text{cf}(15) = 4$

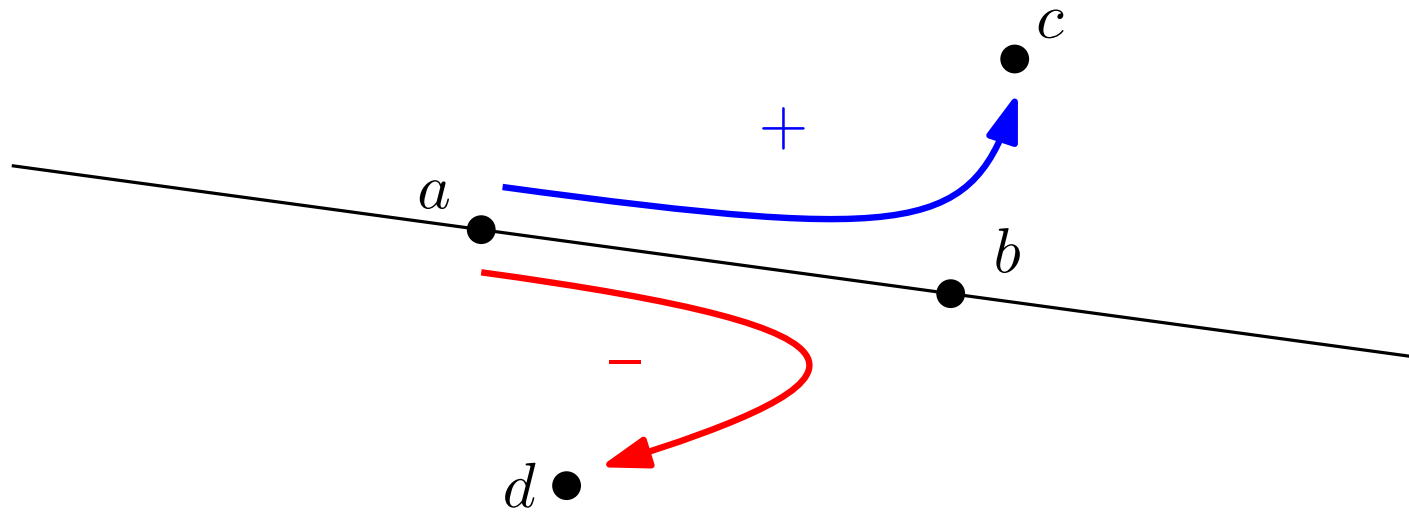
$$\begin{aligned} \text{cf}(n) &= 1 \text{ for } n \leq 4, \\ \text{cf}(n) &= 2 \text{ for } 5 \leq n \leq 9, \\ \text{cf}(n) &= 3 \text{ for } 10 \leq n \leq 14, \\ \text{cf}(n) &= 4 \text{ for } 15 \leq n \leq 20 \end{aligned}$$

k	1	2	3	4	5	6	7	8
$ S_k $	4	9	14	≥ 20	≥ 25	≥ 29	≥ 34	≥ 41

Largest point sets S_k with $\text{cf}(S_k) = k$.

Proving $cf(15) = 4$

- crossings only depend on triple-orientations



Proving $cf(15) = 4$

- crossings only depend on triple-orientations
- proof #1: via exhaustive enumeration of order types
- proof #2: via SAT model

Thank you very much for your attention!

