

Fast Summation of Radial Kernels via QMC Slicing

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The Problem

Given: Points $x_n, y_m \in \mathbb{R}^d$ and coefficients $w_n \in \mathbb{R}$ for $n = 1, \dots, N$, $m = 1, \dots, M$ and a **radial kernel** $K: \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}$,

$$K(x, y) = F(\|x - y\|),$$

with **basis function** $F: [0, \infty) \rightarrow \mathbb{R}$.

Task: Compute the kernel sums

$$s_m = \sum_{n=1}^N w_n K(x_n, y_m) = \sum_{n=1}^N w_n F(\|x_n - y_m\|) \quad \forall m = 1, \dots, M. \quad (1)$$

Applications: KDE, classification via SVMs, dimensionality reduction with kernelized PCA, computation of MMDs or energy distance, MMD gradient flows, Stein variational GD for Bayesian inference

Existing approaches

The direct/exact computation of (1) is unfeasible. \rightsquigarrow Complexity $\mathcal{O}(NM)$

► **Fast Fourier summation** \rightsquigarrow Complexity $\mathcal{O}(N + M + N_{ft}^d \log N_{ft})$

☺ Efficient for $d \in \{1, 2, 3\}$

☺ Becomes unfeasible for larger d

► **Random Fourier features (RFF)** \rightsquigarrow Complexity $\mathcal{O}(P(N + M))$

☺ Error bound $\mathbb{E}[\|F - \tilde{F}\|_\infty] \in \mathcal{O}(1/\sqrt{P})$

☺ Only applicable for positive definite functions (not for Riesz $F(\|x\|) = -\|x\|$)

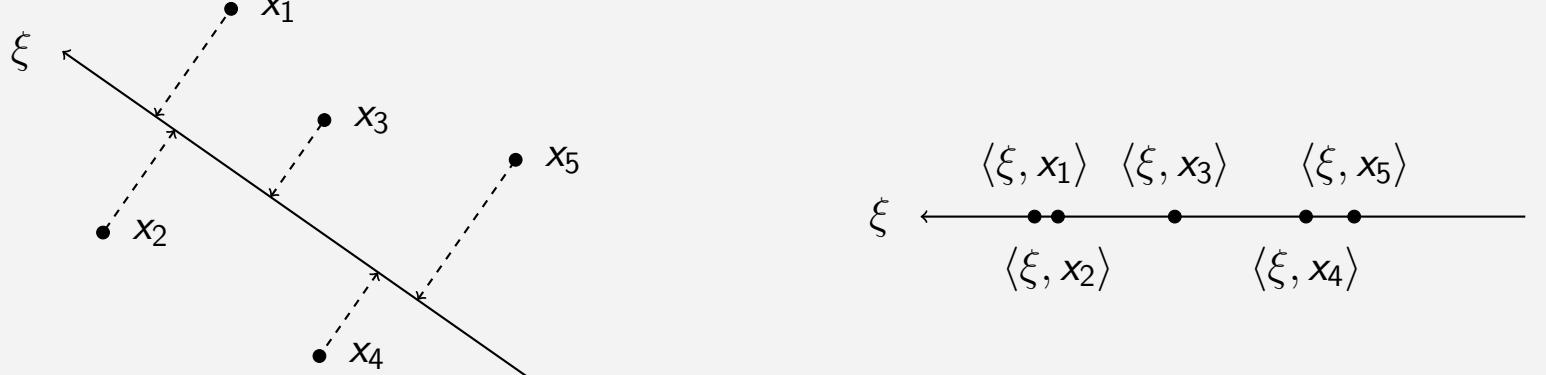
Slicing Summation

Idea: Construct $f: \mathbb{R} \rightarrow \mathbb{R}$ such that

$$F(\|x\|) = \mathbb{E}_{\xi \sim \mathcal{U}_{\mathbb{S}^{d-1}}} [f(\langle x, \xi \rangle)]. \quad (2)$$

Insert (2) into the sum (1) and sample $\xi_p \sim \mathcal{U}_{\mathbb{S}^{d-1}}$ (uniformly from \mathbb{S}^{d-1})

$$s_m = \sum_{n=1}^N w_n \mathbb{E}_{\xi \sim \mathcal{U}_{\mathbb{S}^{d-1}}} f(\langle x_n - y_m, \xi \rangle) \approx \frac{1}{P} \sum_{p=1}^P \underbrace{\sum_{n=1}^N w_n f(\langle x_n - y_m, \xi_p \rangle)}_{\text{1D kernel sum}}.$$



- Reduce the problem to P one-dimensional sums
- ☺ Complexity $\mathcal{O}(P(N + M + N_{ft} \log N_{ft}))$

Known Kernel Pairs (F, f)

Kernel	$F(x)$	$f(x)$
Gauss	$\exp(-\frac{x^2}{2\sigma^2})$	${}_1F_1(\frac{d}{2}, \frac{1}{2}; -\frac{x^2}{2\sigma^2})$
Laplace	$\exp(-\alpha x)$	${}_1F_2(\frac{d}{2}, \frac{1}{2}, \frac{1}{2}; \frac{\alpha^2 x^2}{4}) - \frac{\sqrt{\pi} \alpha x F(\frac{d+1}{2})}{\Gamma(\frac{d}{2})} {}_1F_2(\frac{d+1}{2}, 1, \frac{3}{2}, \frac{\alpha^2 x^2}{4})$
Matérn	$\frac{2^{1-\nu}}{\Gamma(\nu)} (\frac{\sqrt{2\nu}}{\beta} x)^\nu K_\nu(\frac{\sqrt{2\nu}}{\beta} x)$	[Hertrich 2025, Appx C]
Riesz for $r > -1$	$-x^r$	$-\frac{\sqrt{\pi} \Gamma(\frac{d+r}{2})}{\Gamma(\frac{d}{2}) \Gamma(\frac{r+1}{2})} x^r$
Thin Plate Spline	$x^2 \log(x)$	$dx^2 \log(x) + \frac{d}{2} (H_{d/2} - 2 + \log(4)) x^2$

Approximation Error for Random Directions

Theorem: The mean squared error is

$$\mathbb{E}_{\xi_1, \dots, \xi_P \sim \mathcal{U}_{\mathbb{S}^{d-1}}} \left[\left(\frac{1}{P} \sum_{p=1}^P f(\langle x, \xi_p \rangle) - F(\|x\|) \right)^2 \right] = \frac{\mathbb{V}_d[f](x)}{P}$$

with

Kernel	$\mathbb{V}_d[f](x)$
Any positive definite	$\mathbb{V}_d[f](x) \leq F(0)^2 - F(\ x\)^2$
Riesz $F(\ x\) = -\ x\ ^r, r > 0$	$\mathbb{V}_d[f](x) = \text{const}(d, r) F(\ x\)^2$

☺ This gives the exact rate of the expected error

☺ Error rate $1/\sqrt{P}$ for $P \rightarrow \infty$

Numerics

Slicing points ξ_p :

- Random directions
- Orthogonal directions
- Distance QMC points

Compared methods:

- Random Fourier features (RFF)
- Orthogonal Random Features (ORF)
- QMC(Sobol)-RFF for Gauss kernel

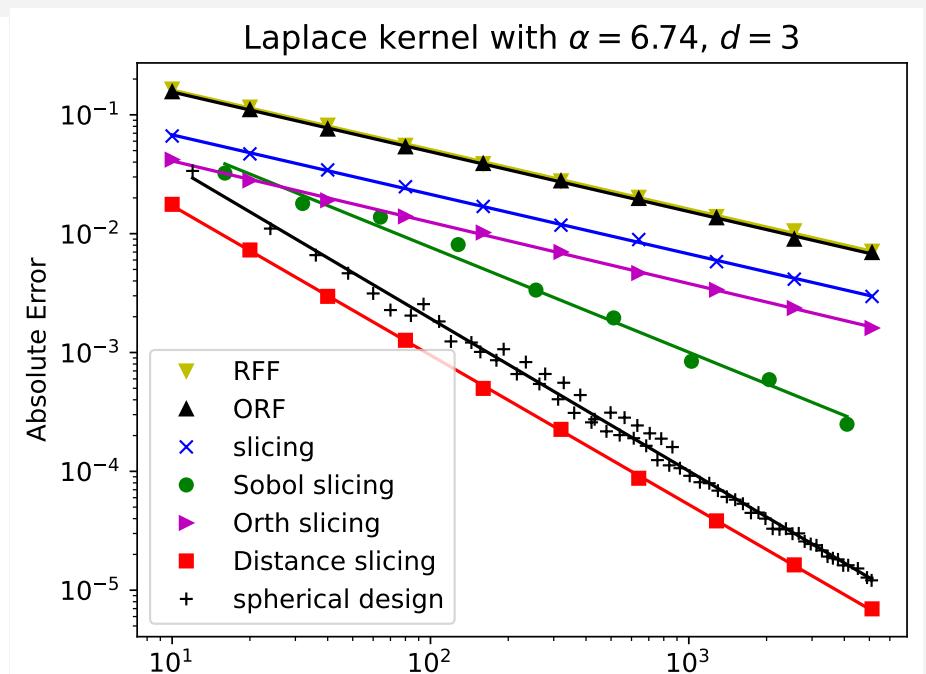
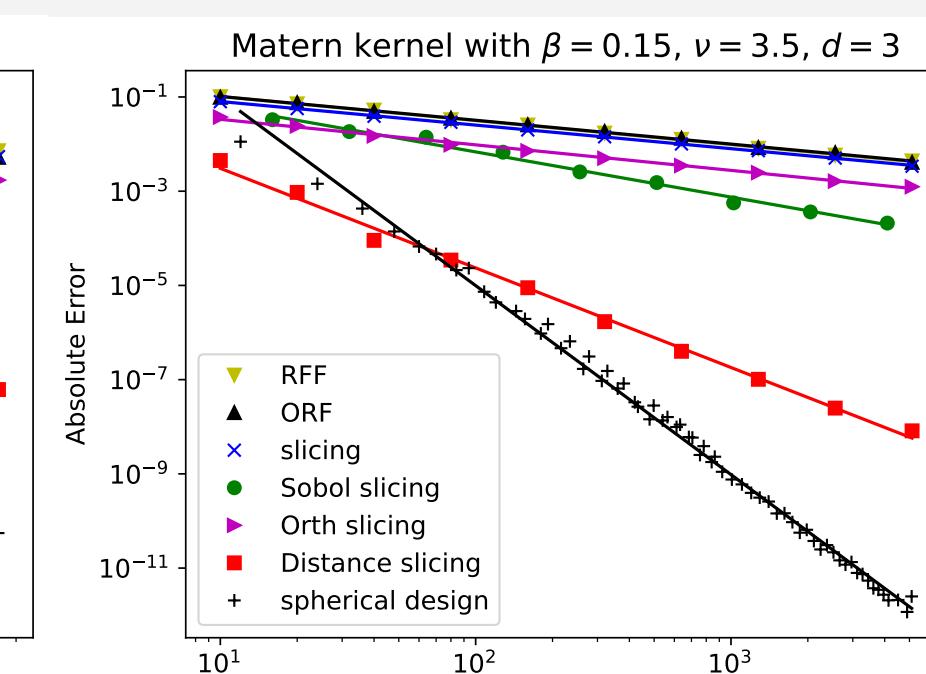
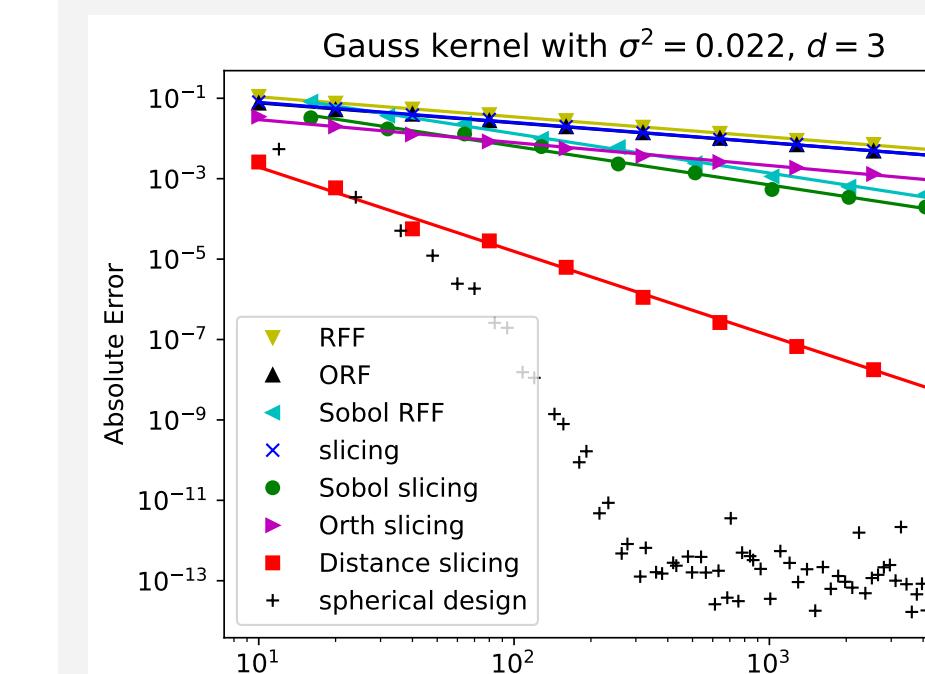


Figure: Mean approximation error $|F(\|x\|) - \frac{1}{P} \sum_{p=1}^P f(\langle x, \xi_p \rangle)|$ for dimension $d = 3$

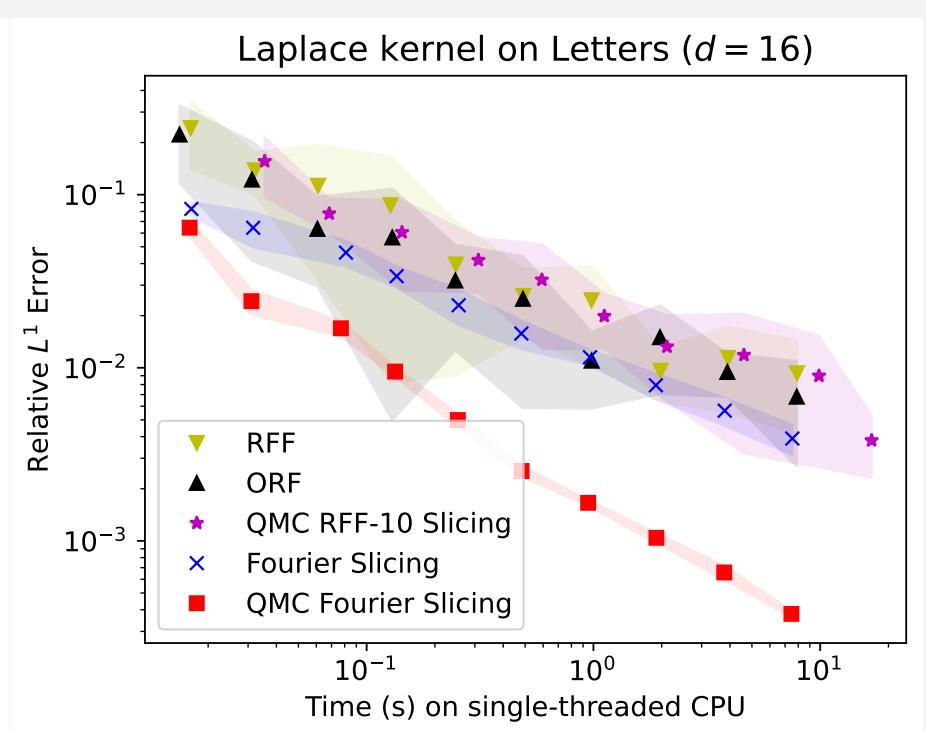
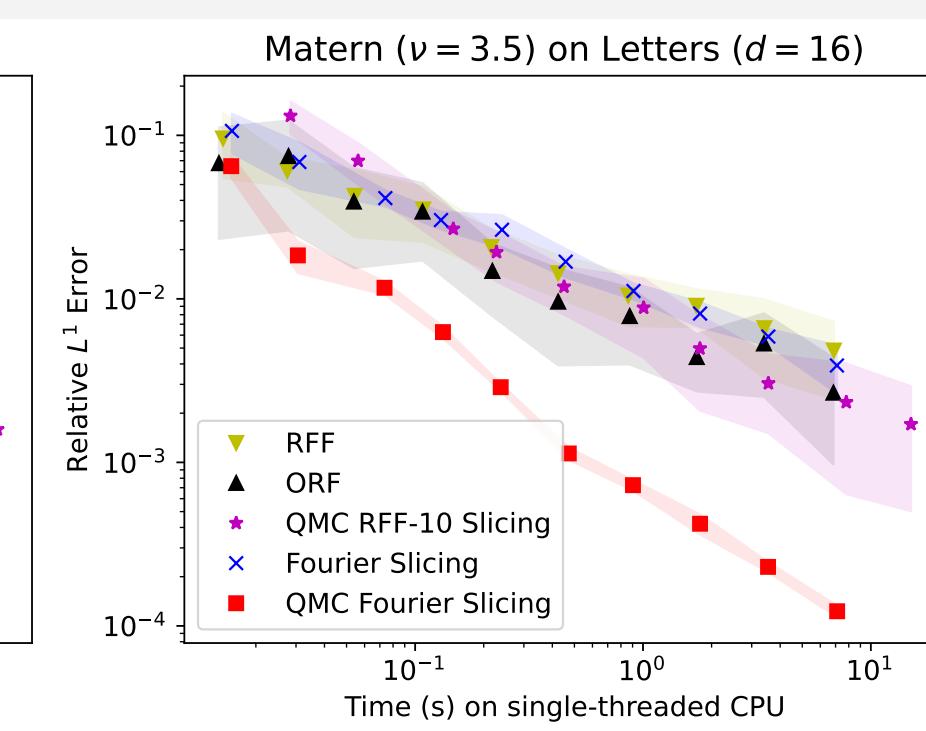
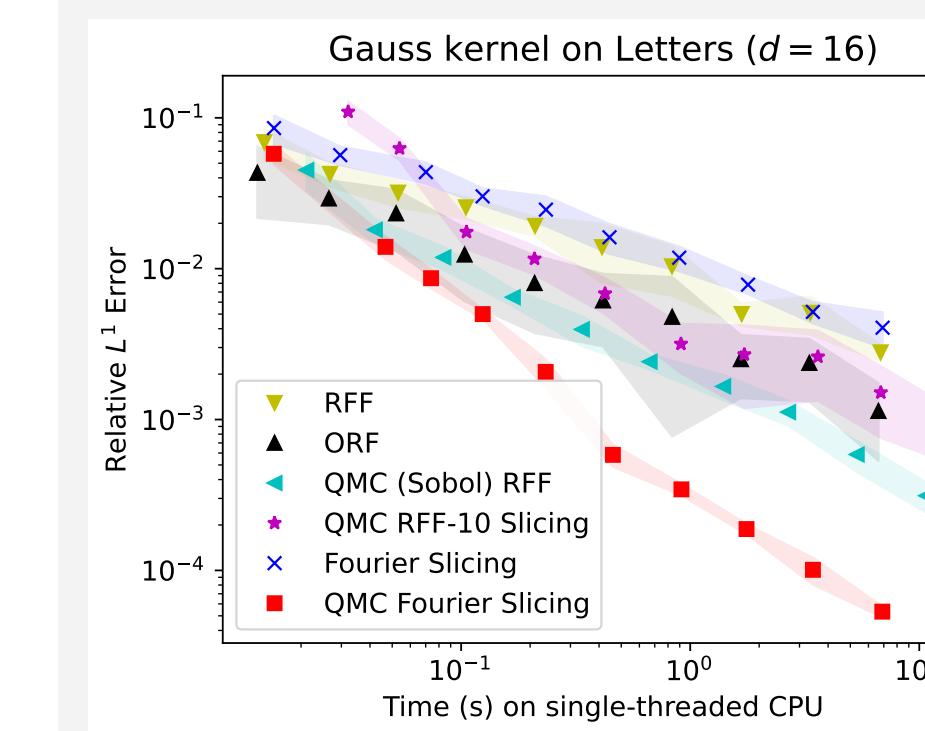


Figure: Slicing Summation on CPU (Letters dataset with $d = 16$ and $M = N = 20000$).

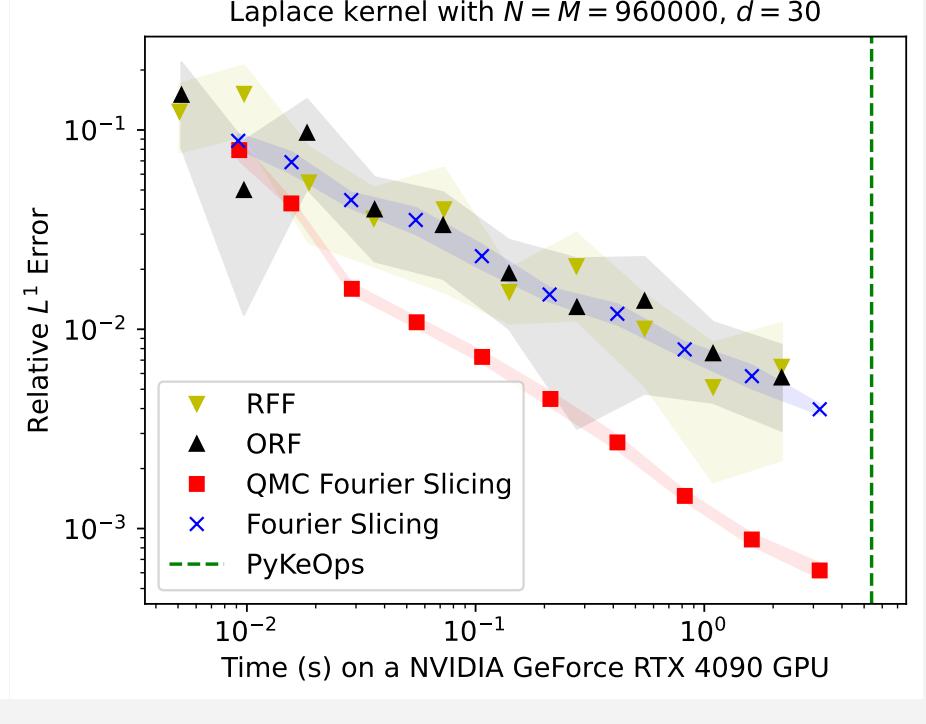
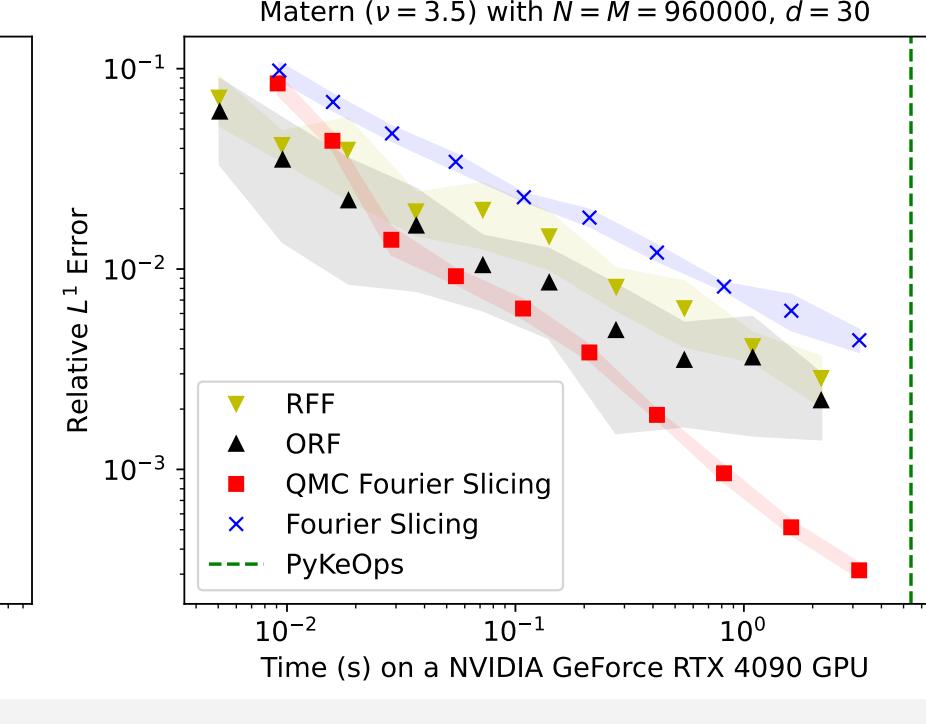
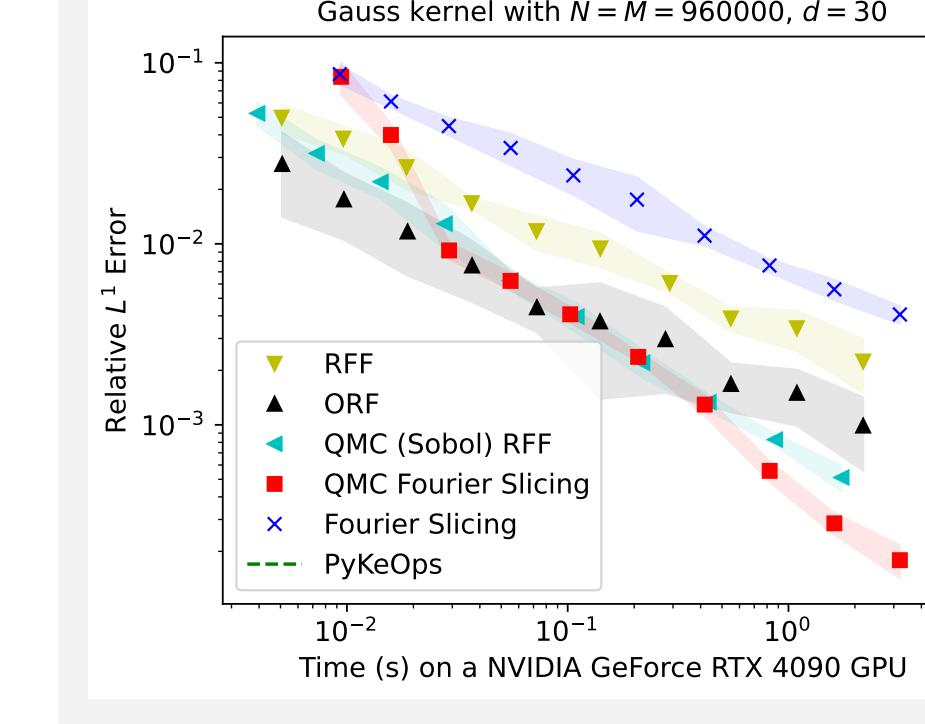


Figure: GPU times. The dashed green line is the direct computation with PyKEops.

(Computable) QMC Designs

Spherical t -designs: Integrate polynomials of degree $\leq t$ exactly

☺ Always exists and is a QMC design for $s = t$ whenever $s > \frac{d-1}{2}$

☺ Hard to compute (numerically done only for $d \leq 4$)

Distance points: Maximize

$$\mathcal{E}(\xi^P) = \sum_{p,q=1}^P \|\xi_p^P - \xi_q^P\|$$

☺ Is a QMC design for $s = \frac{d+1}{2}$

☺ Resulting worst case error rate $\mathcal{O}(1/P^{2(d-1)})$

For $d = 3$: error rate $\mathcal{O}(1/P^3)$

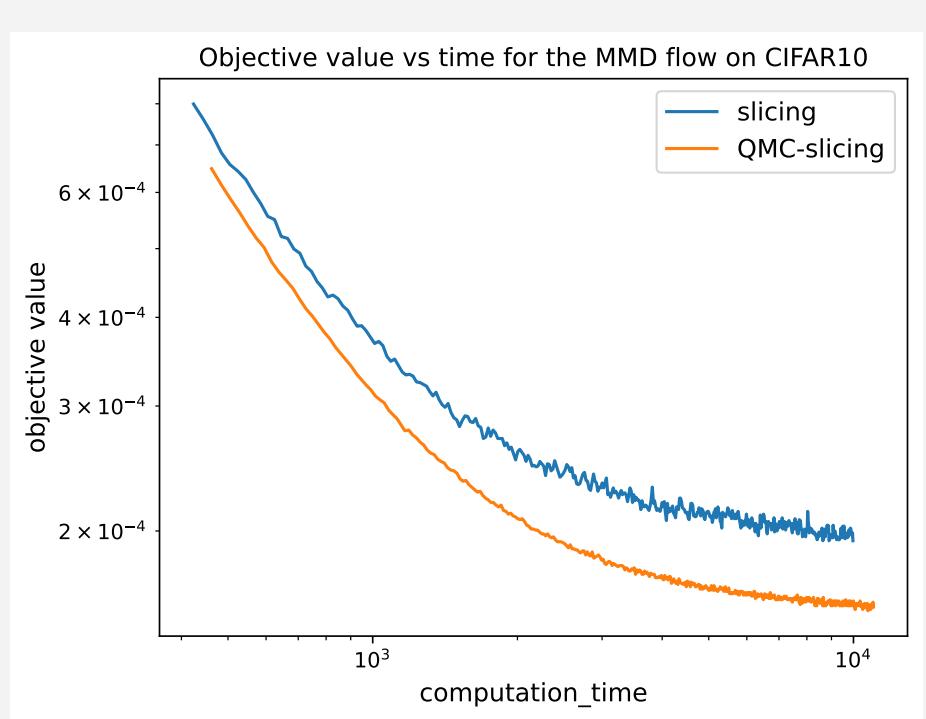
Sobol/Halton points:

☺ Sobol/Halton QMC sequences on $[0, 1]^d$ can be transformed to \mathbb{S}^{d-1}

☺ Unclear whether the transformed sequences are QMC designs on \mathbb{S}^{d-1}

Application: MMD as Loss Function

- Minimize $\text{Loss}(\mathbf{x}) = \text{MMD}(\mathbf{x}, \mathbf{y})$ via gradient descent
- use Riesz kernel $F(x) = -\|x\|$ \Rightarrow RFFs are not applicable
- Direct computation is extremely slow
- CIFAR10 dataset ($d = 3072$)



Paper:



Code:



Library:

