

On the Maximum Number of Crossings in Star-Simple Drawings of K_n with No Empty Lens*

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Abstract. A star-simple drawing of a graph is a drawing in which adjacent edges do not cross. In contrast, there is no restriction on the number of crossings between two independent edges. When allowing empty lenses (a face in the arrangement induced by two edges that is bounded by a 2-cycle), two independent edges may cross arbitrarily many times in a star-simple drawing. We consider star-simple drawings of K_n with no empty lens. In this setting we prove an upper bound of $3((n-4)!)$ on the maximum number of crossings between any pair of edges. It follows that the total number of crossings is finite and upper bounded by $n!$.

Keywords: star-simple drawings · topological graphs · edge crossings.

1 Introduction

A *topological drawing* of a graph G is a drawing in the plane where vertices are represented by pairwise distinct points, and edges are represented by Jordan arcs with their vertices as endpoints. Additionally, edges do not contain any other vertices, every common point of two edges is either a proper crossing or a common endpoint, and no three edges cross at a single point. A *simple drawing* is a topological drawing in which adjacent edges do not cross, and independent edges cross at most once.

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We study a broader class of topological drawings, which are called *star-simple* drawings, where adjacent edges do not cross, but independent edges may cross any number of times; see Fig. 1 for illustration. In such a drawing, for every vertex v the induced substar centered at v is simple, that is, the drawing restricted to the edges incident to v forms a plane drawing. In the literature (e.g., [1,2]) these drawings also appear under the name *semi-simple*, but we prefer star-simple because the name is much more descriptive.

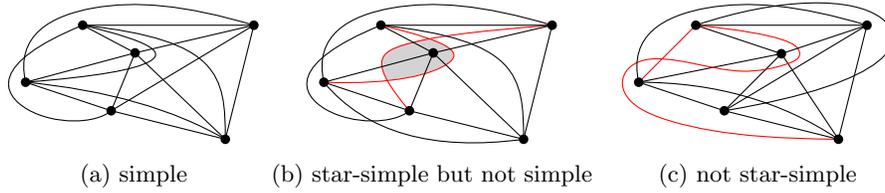


Fig. 1: Topological drawings of K_6 and a (nonempty) lens (shaded in (b)).

In contrast to simple drawings, star-simple drawings can have regions or cells whose boundary consists of two continuous pieces of (two) edges. We call such a region a *lens*; see Fig. 1b. A lens is *empty* if it has no vertex in its interior. If empty lenses are allowed, the number of crossings in star-simple drawings of graphs with at least two edges is unbounded (twisting), as illustrated in Fig. 2a. We restrict our attention to star-simple drawings with no empty lens. This restriction is—in general—not sufficient to guarantee a bounded number of crossings (spiraling), as illustrated in Fig. 2b. However, we will show that star-simple drawings of the complete graph K_n with no empty lens have a bounded number of crossings.

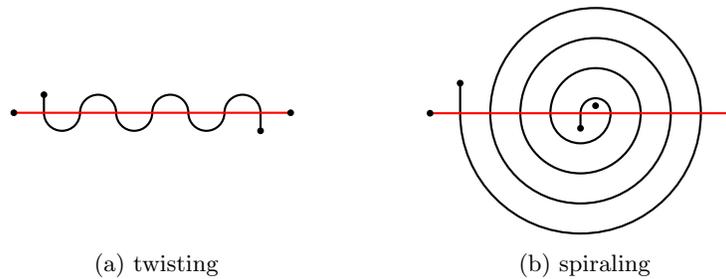


Fig. 2: Constructions to achieve an unbounded number of crossings.

Empty lenses also play a role in the context of the crossing lemma for multigraphs [5]. This is because a group of arbitrarily many parallel edges can be drawn without a single crossing. Hence, for general multigraphs there is no hope

to get a lower bound on the number of crossings as a function of the number of edges. However, if we forbid empty lenses, we cannot draw arbitrarily many parallel edges.

Kynčl [3, Section 5 "Picture hanging without crossings"] proposed a construction of two edges in a graph on n vertices with an exponential number (2^{n-4}) of crossings and no empty lens; see Fig. 3. This configuration can be completed to a star-simple drawing of K_n , cf. [6]. For $n = 6$ it is possible to have one more crossing while maintaining the property that the drawing can be completed to a star-simple drawing of K_6 ; see Fig. 4. Repeated application of the doubling construction of Fig. 3 leads to two edges with $2^{n-4} + 2^{n-6}$ crossings in a graph on n vertices. This configuration can be completed to a star-simple drawing of K_n . We suspect that this is the maximum number of crossings of two edges in a star-simple drawing of K_n .

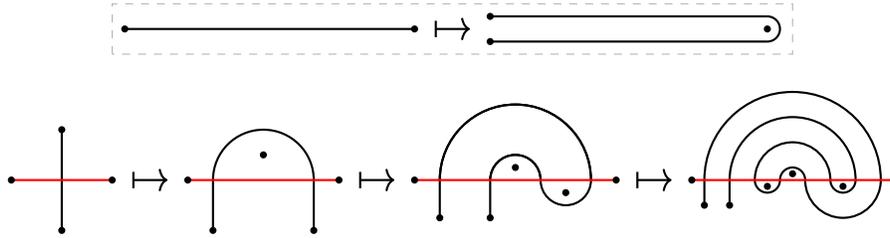


Fig. 3: The doubling construction yields an exponential number of crossings.

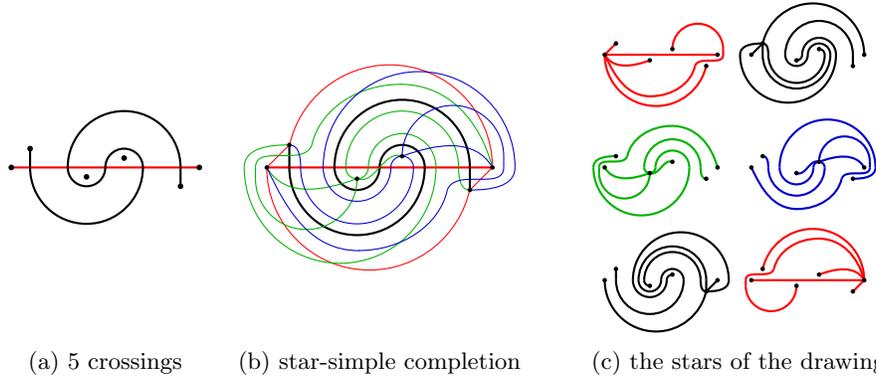


Fig. 4: Two edges with $2^{n-4} + 2^{n-6}$ crossings in a star-simple drawing of K_n , for $n = 6$.

2 Crossing patterns

In this section we study the induced drawing $D(e, e')$ of two independent edges e and e' in a star-simple drawing D of the complete graph. We start by observing that the endpoints of e and e' must lie in the same region of $D(e, e')$. This fact was also used in earlier work by Aichholzer et al. [1] and by Kynčl [4].

Lemma 1. *The four vertices incident to e and e' belong to the same region of $D(e, e')$.*

Proof. Assuming that the two edges cross at least two times, the drawing $D(e, e')$ has at least two regions. Otherwise, the statement is trivial. If the four vertices do not belong to the same region of $D(e, e')$, then there is a vertex u of e and a vertex v of e' that belong to different regions. Now consider the edge uv in the drawing D of the complete graph. This edge has ends in different regions of $D(e, e')$, whence it has a crossing with either e or e' . This, however, makes a crossing in the star of u or v . This contradicts the assumption that D is a star-simple drawing. \square

Lemma 1 implies that the deadlock configurations as shown in Fig. 5a do not occur in star-simple drawings of complete graphs. Formally, a *deadlock* is a pair e, e' of edges such that not all incident vertices lie in the same region of the drawing $D(e, e')$.

Now suppose that D is a star-simple drawing of a complete graph with no empty lens. In this case we can argue that e and e' do not form a configuration as the black edge e and the red edge e' in Fig. 5b. Indeed, that configuration has an interior lens L and by assumption this lens is non-empty, i.e., L contains a vertex x . Let e and e' be the black and the red edge in Fig. 5b, respectively, and let u be a vertex of e . The edge xu (the green edge in the figure) has no crossing with e , hence it follows the "tunnel" of the black edge. This yields a deadlock configuration of the edges xu and e' . Note that if in Fig. 5b instead of drawing the green edge xu we connect x with an edge f to one of the vertices of the red edge e' such that f and the red edge have no crossing, then f and the black edge e form a deadlock.

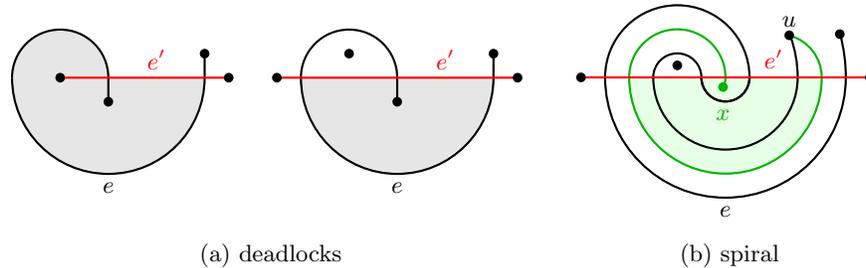


Fig. 5: Constructions to achieve an unbounded number of crossings.

We use this intuition to formally define a spiral. Two edges e, e' form a *spiral* if they form a lens L such that if we place a vertex x in L and draw a curve γ connecting x to a vertex u of e so that γ does not cross e , then γ and e' form a deadlock. The discussion above proves the following lemma:

Lemma 2. *A star-simple drawing of a complete graph with no empty lens has no pair e, e' of edges that form a spiral.*

3 Crossings of pairs of edges

In this section we derive an upper bound for the number of crossings of two edges in a star-simple drawing of K_n with no empty lens.

Theorem 1. *Consider a star-simple drawing of K_n with no empty lens. If $C(k)$ is the maximum number of crossings of a pair of edges that (a) form no deadlock and no spiral and such that (b) all lenses formed by the two edges can be hit by k points, then $C(k) \leq e \cdot k!$, where $e \approx 2.718$ is Euler's number.*

Proof. Due to Lemma 1 we can assume that all four vertices of e and e' are on the outer face of the drawing $D(e, e')$. We think of e' as being drawn red and horizontally and of e as being a black meander edge. Let p_1, \dots, p_k be points hitting all the lenses of the drawing $D(e, e')$. Let u be one of the endpoints of e . For each $i = 1, \dots, k$ we draw an edge e_i connecting p_i to u such that e_i has no crossing with e and, subject to this, the number of crossings with e' is minimized. Fig. 6 shows an example.

Note that we do not claim that all these edges e_1, \dots, e_k together with e and e' can be extended to a star-simple drawing of a complete graph. Therefore, we cannot use Lemma 2 directly but state the assumption (a) instead.

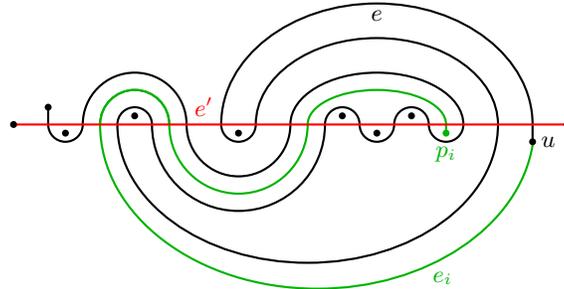


Fig. 6: The drawing $D(e, e')$ and an edge e_i connecting p_i to u .

We claim the following three properties:

- (P1) The edges e_i and e' form no deadlock and no spiral.

- (P2) All lenses of e_i and e' are hit by the $k - 1$ points $p_1, \dots, p_{i-1}, p_{i+1}, \dots, p_k$.
- (P3) Between any two crossings of e and e' from left to right, i.e., in the order along e' , there is at least one crossing of e' with one of the edges e_i .

Before proving the properties, we show that they imply the statement of the theorem by induction on k . The base case $1 = C(0) \leq e \cdot 0! = e$ is obvious. From (P1) and (P2) we see that the number X_i of crossings of e_i and e' is upper bounded by $C(k - 1)$. From (P3) we obtain that $C(k) \leq 1 + \sum_i X_i$. Combining these we get

$$C(k) \leq k \cdot C(k - 1) + 1 \leq k! \cdot \sum_{s=0}^k \frac{1}{s!} \leq k! \cdot e. \quad \square$$

For the proof of the three claims we need some notation. Let $\xi_1, \xi_2, \dots, \xi_N$ be the crossings of e and e' indexed according to the left to right order along the horizontal edge e' . Let g_i and h_i be the pieces of e' and e , respectively, between crossings ξ_i and ξ_{i+1} . The bounded region enclosed by $g_i \cup h_i$ is the *bag* B_i and g_i is the *gap* of the bag. In the drawing $D(e, e')$ the bags B_i where h_i is a crossing free piece of e are exactly the inclusion-wise minimal lenses formed by e and e' . From now on when referring to a *lens* we always mean such a minimal lens. Indeed if there is no empty minimal lens, then there is no empty lens. The following observation is crucial.

Observation 2 *For two bags B_i and B_j the open interiors are either disjoint or one is contained in the other.*

Proof. Every bag is bounded by a closed Jordan curve, and the boundaries of two distinct bags do not cross (at most they may touch at a single point that is one of $\xi_1, \xi_2, \dots, \xi_N$). \square

Observation 2 implies that the containment order on the bags is a downwards branching forest. The minimal elements in the containment order are the lenses. Consider a lens L and the point p_i inside L . Since the vertex u of e is in the outer face of $D(e, e')$, the edge e_i has to leave each bag that contains L . Furthermore, by definition e_i does not cross e and therefore it has to leave a bag B containing L through the gap g of B . We now reformulate and prove the third claim (P3).

- (P3') For each pair ξ_i, ξ_{i+1} of consecutive crossings on e' there is a lens L and a point $p_j \in L$ such that e_j crosses e' between ξ_i and ξ_{i+1} .

Proof sketch for (P3'). The pair ξ_i, ξ_{i+1} is associated with the bag B_i . In the containment order of bags a minimal bag below B_i is a lens, let L be any of the minimal elements below B_i . By assumption, L contains a point p_j . Since $L \subseteq B_i$, we have that also $p_j \in B_i$. Thus, it follows that e_j has a crossing with the gap g_i , i.e., e_j has a crossing with e' between ξ_i and ξ_{i+1} . \square

Proof sketch for (P1). We have to show that e_i and e' form no deadlock and no spiral. The minimality condition in the definition of e_i implies that if $L = B_{i_1} \subset$

$B_{i_2} \subset \dots \subset B_{i_t}$ is the maximal chain of bags with minimal element L , then e_i crosses the gaps of these bags in the given order and has no further crossings with e' . If γ is a curve from L to u that avoids e_i , then in the ordered sequence of gaps crossed by γ we find a subsequence that is identical to the ordered sequence of gaps crossed by e_i . Since e and e' form no spiral, there is such a curve γ that forms no deadlock with e' . Therefore, e_i forms no deadlock with e' , either.

Now assume that e_i and e' form a spiral. Let B be the largest bag containing p_i . Think of B as a drawing of e_i with a broad pen, which may also have some extra branches that have no correspondence in e_i , see Fig. 7. The formalization of this picture is that for every bag β formed by e_i with e' there is a bag $B(\beta)$ formed by e and e' with $B(\beta) \subset \beta$. Now, if there is a lens λ formed by e_i with e' such that every e_i -avoiding⁵ curve to u is a deadlock with e' , then there is a lens $L(\lambda)$ formed by e and e' with $L(\lambda) \subset \lambda$ such that every e -avoiding curve from $L(\lambda)$ to u is also B -avoiding and hence e_i -avoiding. Thus, every such curve has a deadlock with e' , whence e and e' form a spiral, contradiction. \square

Proof sketch for (P2). We know by (P1) that e_i and e' form no deadlock. Therefore, by Lemma 1, the vertices of e_i and e' belong to the same region of $D(e_i, e')$. All crossings of e_i with e' correspond to bags of e and e' , therefore the vertices of e and e' are in the outer face of $D(e_i, e')$. Together this shows that p_i is also in the outer face of $D(e_i, e')$. Since every lens of $D(e_i, e')$ contains a lens of $D(e, e')$, it also contains one of the points hitting all lenses of $D(e, e')$. Hence, all lenses of $D(e_i, e')$ are hit by the $k - 1$ points $p_1, \dots, p_{i-1}, p_{i+1}, \dots, p_k$. \square

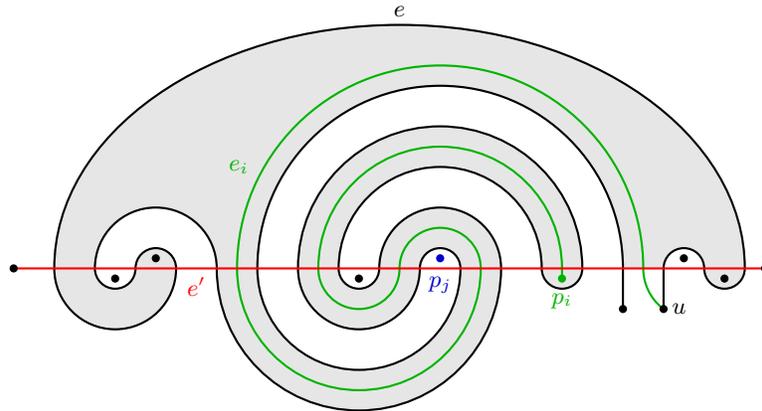


Fig. 7: An edge e_i (green) that forms a spiral with e' . The bag B in gray and the lens $L(\lambda)$ marked with the vertex p_j (blue).

⁵ that is, disjoint from e_i except for possibly a shared endpoint

4 Crossings in complete drawings

Accounting for the four endpoints of the two crossing edges we have $k \leq n - 4$ in Theorem 1. Therefore, we obtain that the number of crossings of a pair of edges in a star-simple drawing of K_n without empty lens is upper bounded by $3(n-4)!$. This directly implies that the drawing of K_n has at most $n!$ crossings. This is the first finite upper bound on the number of crossings in star-simple drawings of the complete graph K_n . We know drawings of K_n in this drawing mode that have an exponential number of crossings. Thus, it would be interesting to reduce the huge gap between the upper and the lower bound. Specifically, can a star-simple drawing of K_n have two edges with more than $2^{n-4} + 2^{n-6}$ crossings?

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