

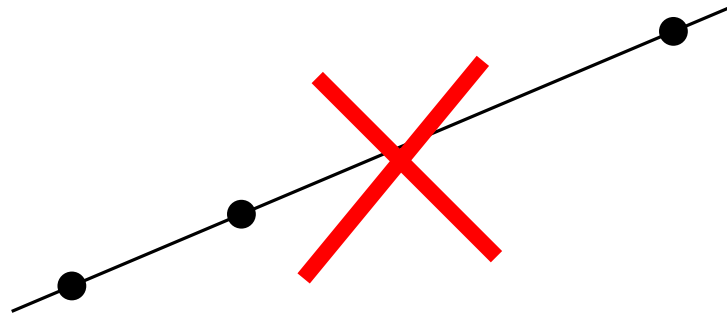


# Erdős–Szekeres-type problems in the real projective plane

Martin Balko, Manfred Scheucher, Pavel Valtr

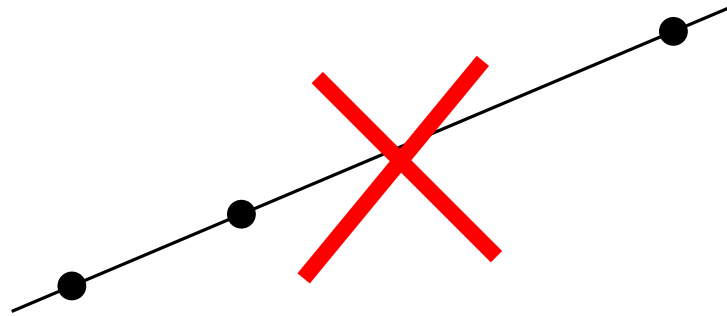
# General Position

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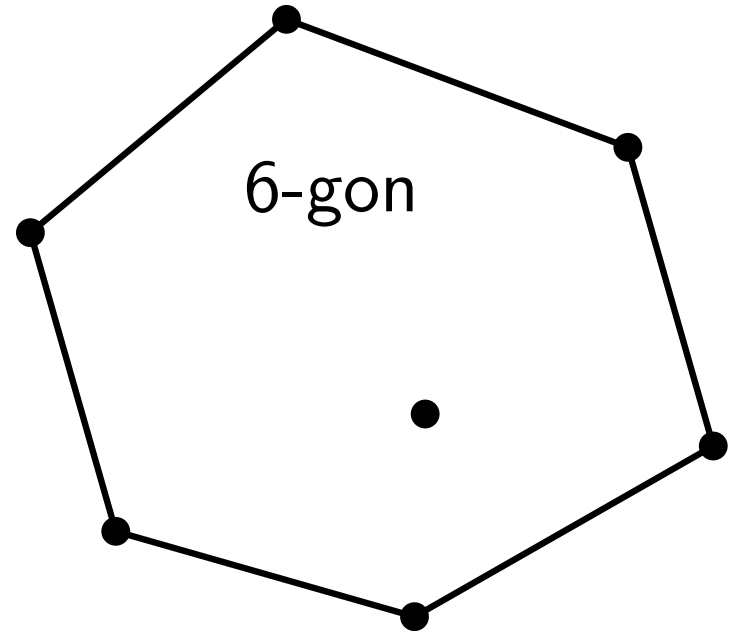
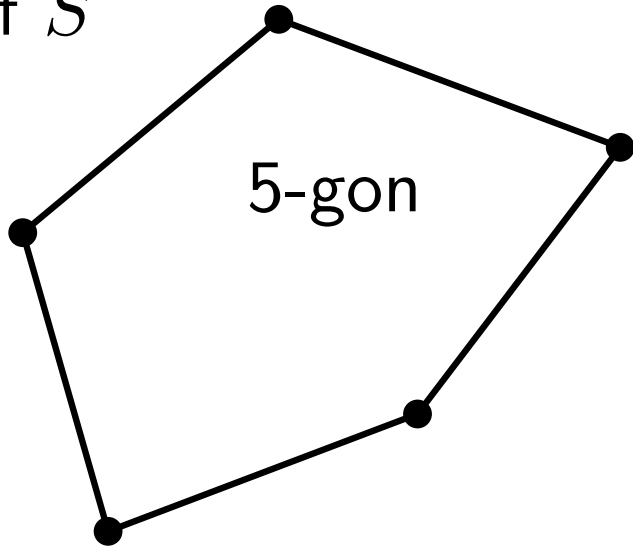


throughout this presentation, every set is in general position

# The Affine World

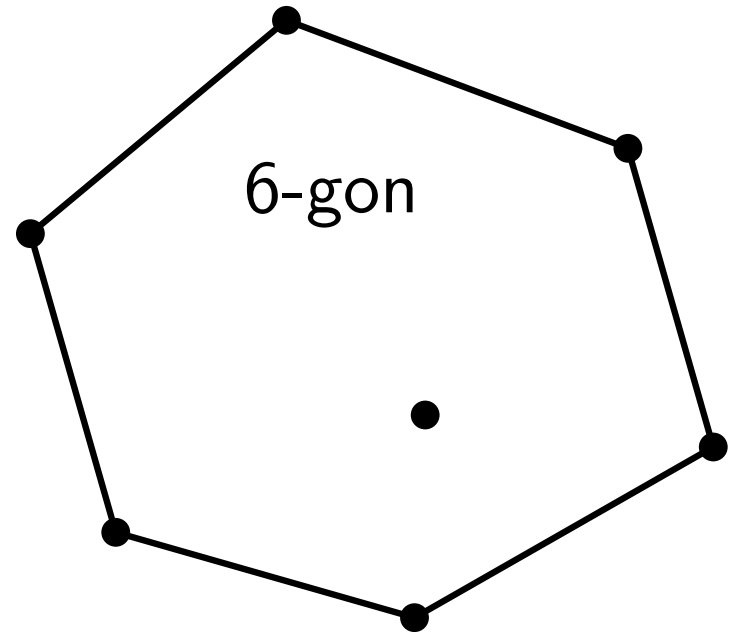
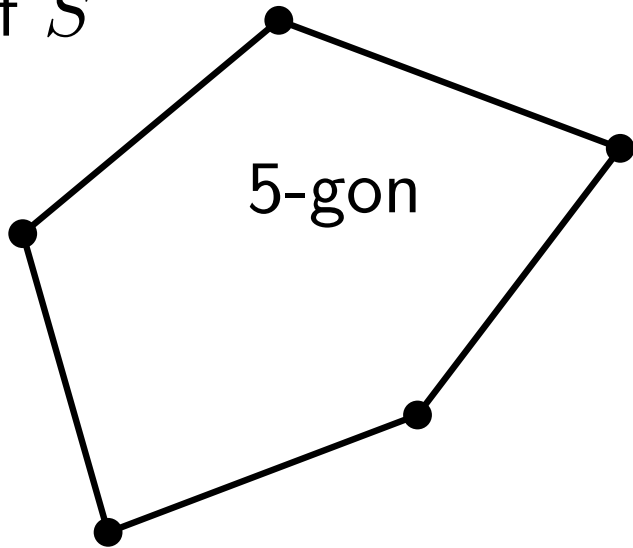
## $k$ -Gons

a  $k$ -gon (in  $S$ ) is a convex polygon spanned by  $k$  points of  $S$



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**Theorem (Erdős & Szekeres 1935).**

$\forall k \in \mathbb{N}$ ,  $\exists$  a smallest integer  $g(k)$  such that every set of  $g(k)$  points determines a  $k$ -gon.

## $k$ -Gons

**Theorem** (Erdős & Szekeres '35)

$$2^{k-2} + 1 \leq g(k) \leq \binom{2k-4}{k-2} + 1$$



equality conjectured by Szekeres, Erdős offered 500\$ for a proof

## $k$ -Gons

**Theorem** (Erdős & Szekeres '35)

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∴ several improvements of order  $4^{k-o(k)}$

**Theorem.**  $g(k) \leq 2^{k+o(k)}$ . [Suk '16]



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**Theorem.**  $g(k) \leq 2^{k+o(k)}$ . [Suk '16]

- $g(k) \leq 2^{k+O(k^{2/3} \log k)}$  [Suk '16]

- $g(k) \leq 2^{k+O(\sqrt{k \log k})}$ ,

also for pseudo-configurations of points

[Holmsen, Mojarrad, Pach and Tardos '17]

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Known:  $g(4) = 5$ ,  $g(5) = 9$ ,  $g(6) = 17$



computer assisted proof, 1500 CPU hours [Szekeres–Peters '06]

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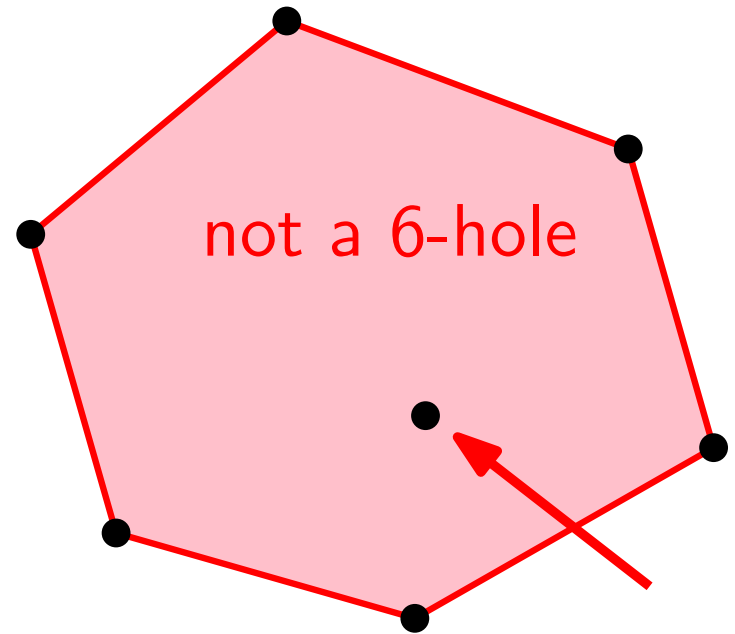
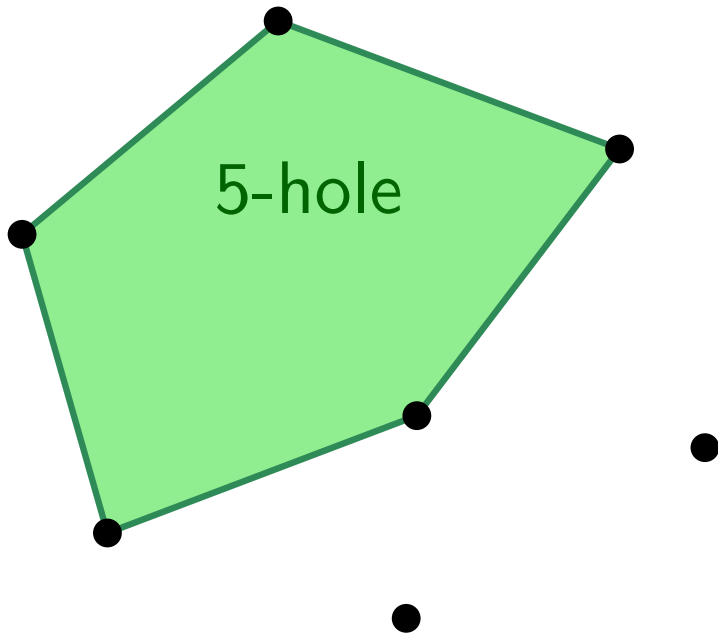
< 1 hour using SAT solvers [S.'18, Marić '19]

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## $k$ -Holes

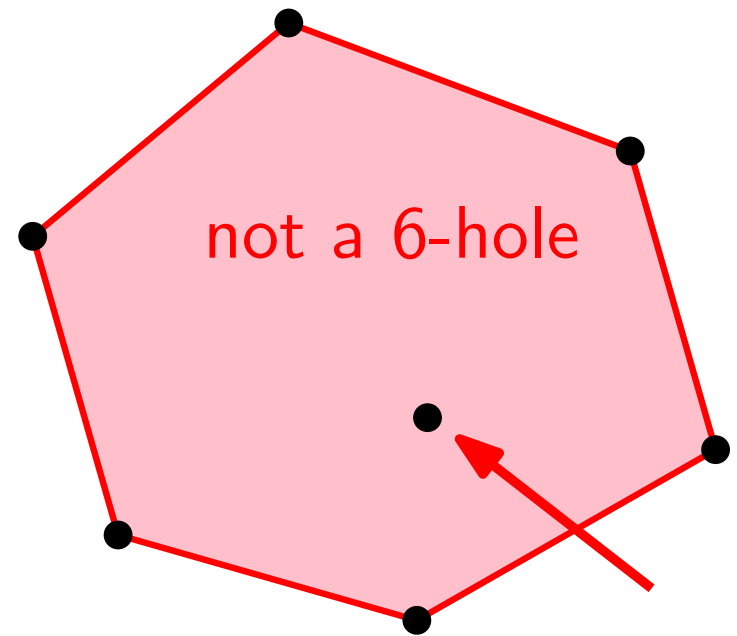
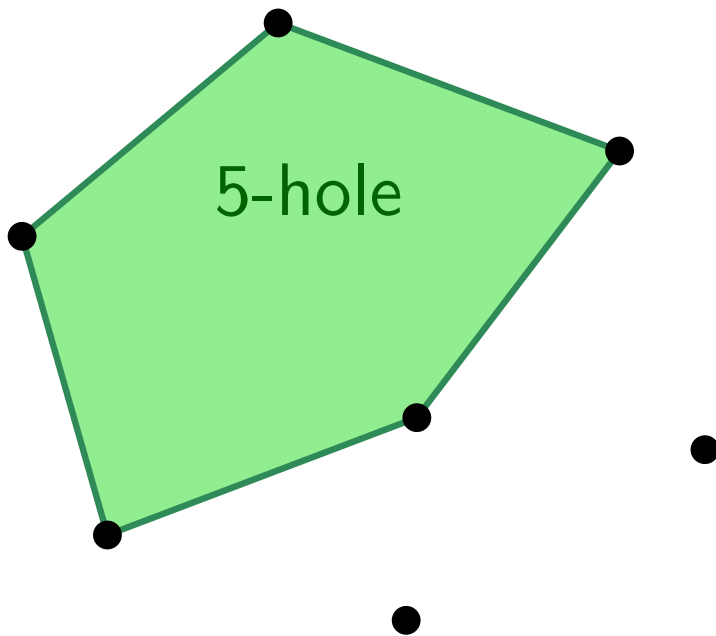
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- 3 points  $\Rightarrow \exists$  3-hole
- 5 points  $\Rightarrow \exists$  4-hole
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about 30 pages



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about 30 pages < 1 cpu hour using SAT (S. '22+)

# $k$ -Holes

- $h(4) = 5, h(5) = 10, 30 \leq h(6) \leq g(9), h(7) = \infty$

Harborth '78

Overmars '02

Gerken '08

Horton '83

# $k$ -Holes

exact value remains unknown

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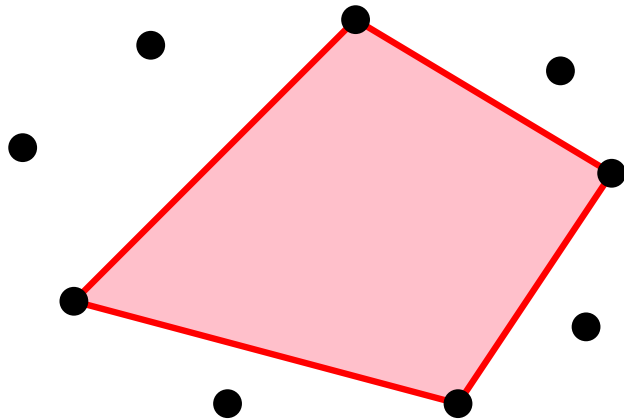
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maximum # of  $k$ -gons among all sets of  $n$  points?

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$n$  points in convex position

any  $k$ -subset is  $k$ -gon

$$\Rightarrow \max = \binom{n}{k}$$

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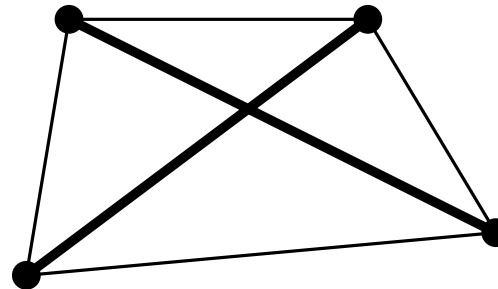
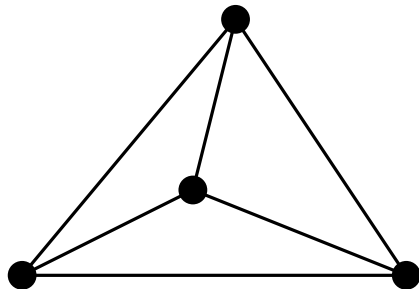
$g_k(n)$  := minimum # of  $k$ -gons among all sets of  $n$  points

- $g_k(n) = \Theta(n^k)$

- $k = 4$  : **rectilinear crossing number of  $K_n$** :

$$g_4(n) = \overline{cr}(K_n) \sim c_4 \cdot \binom{n}{4} \text{ with } 0.3799 < c_4 < 0.3805$$

[Ábrego et al. '08, Aichholzer et al. '20]



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[Ábrego et al. '08, Aichholzer et al. '20]
- various notions of **crossing numbers** have been studied intensively (not necessarily straight-line drawings, not necessarily complete graphs)



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- $h_3, h_4$  both in  $\Theta(n^2)$
- $h_5$  in  $\Omega(n \log^{4/5} n)$  and  $O(n^2)$

[Aichholzer, Balko, Hackl, Kynčl,  
Parada, S., Valtr, and Vogtenhuber '17]  
(computer assisted proof, 20 pages)

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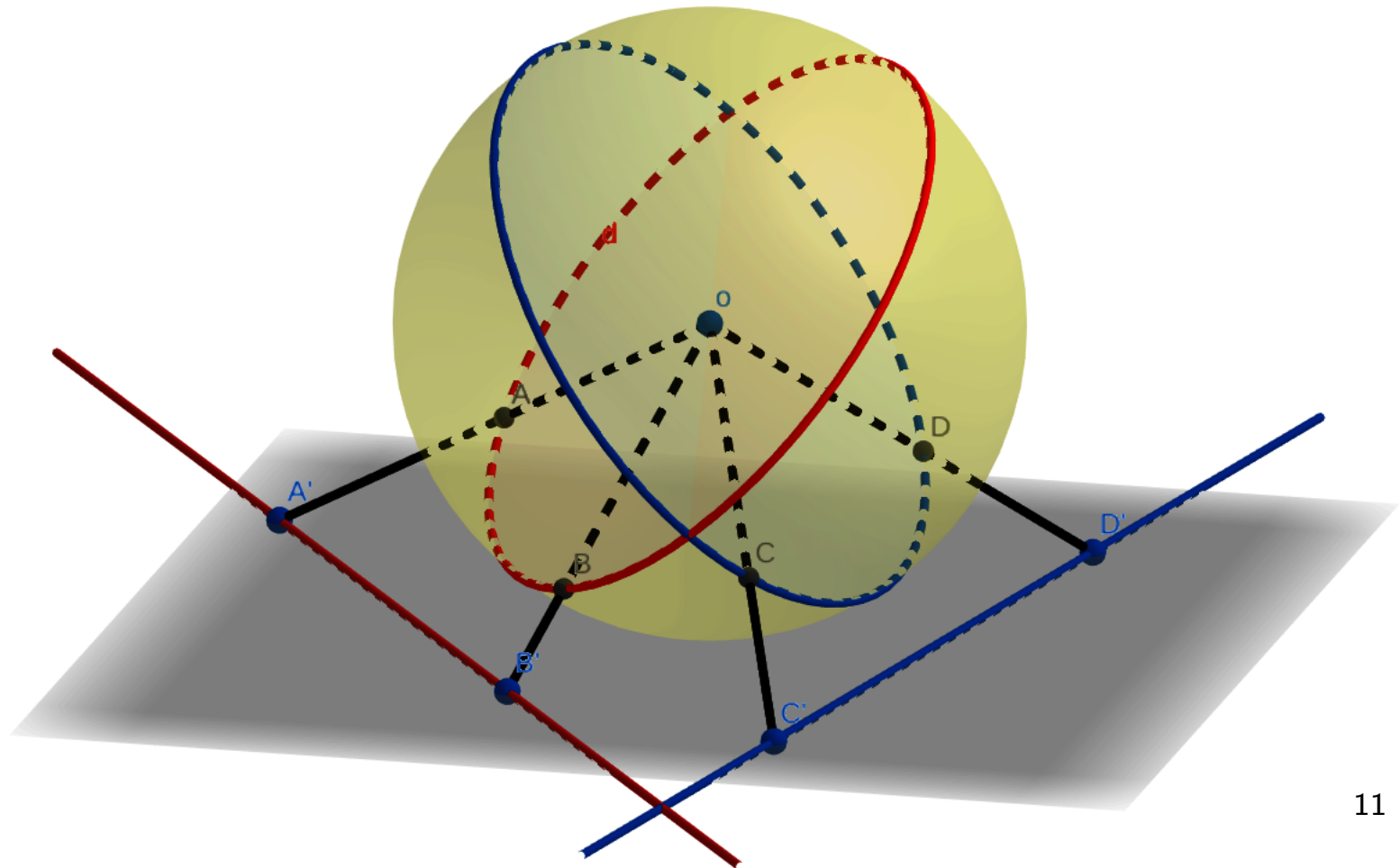
- $h_k(n) = 0$  for  $k \geq 7$  [Horton '83]

**Conjecture:**

$h_5, h_6$  are both in  $\Theta(n^2)$

# The Projective World

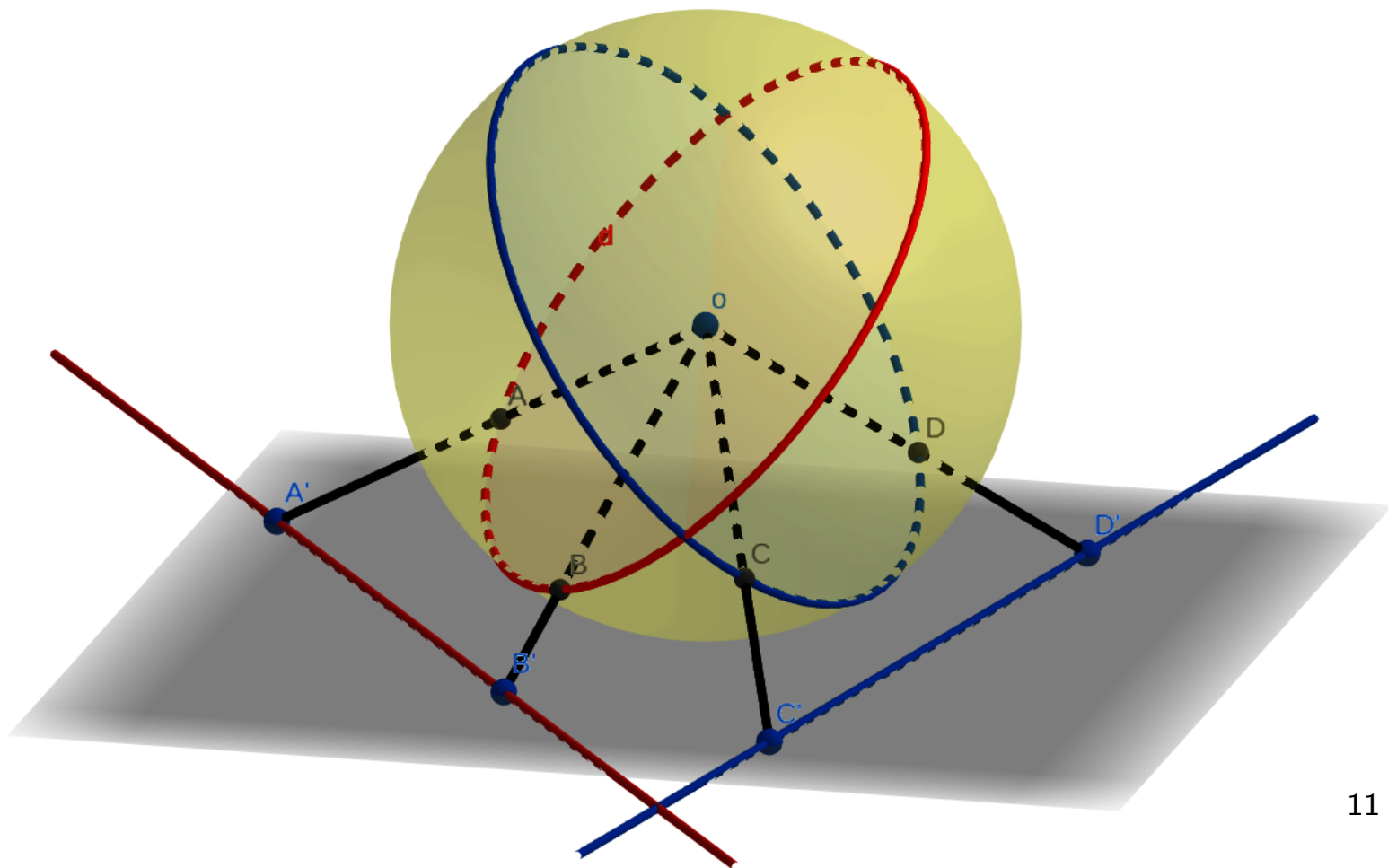
# Convex Sets in Projective Plane (Steinitz 1913)





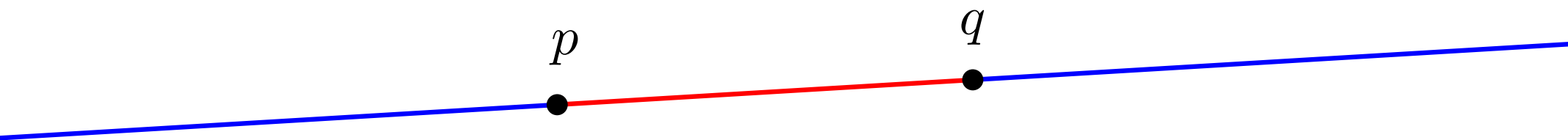
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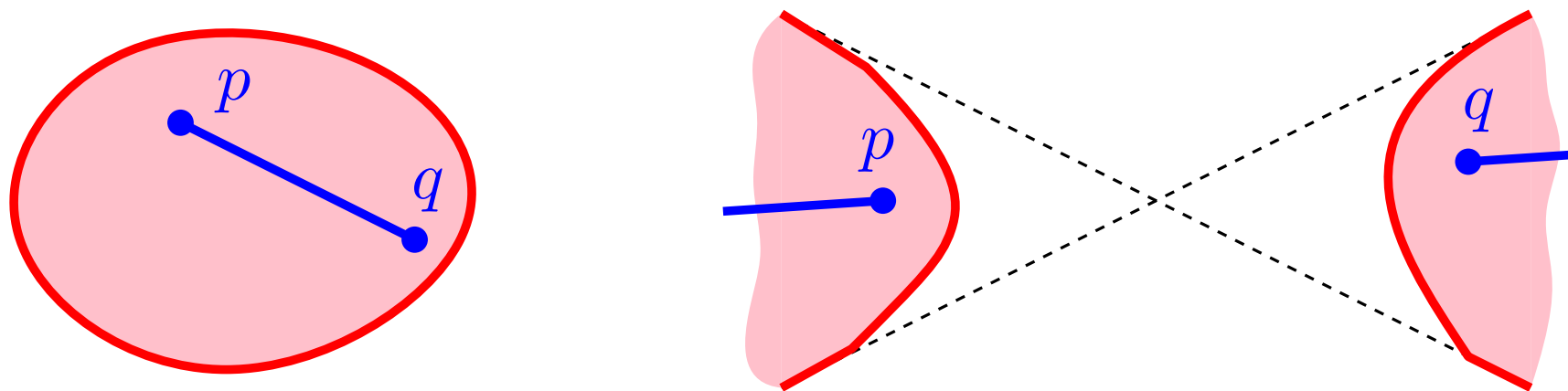
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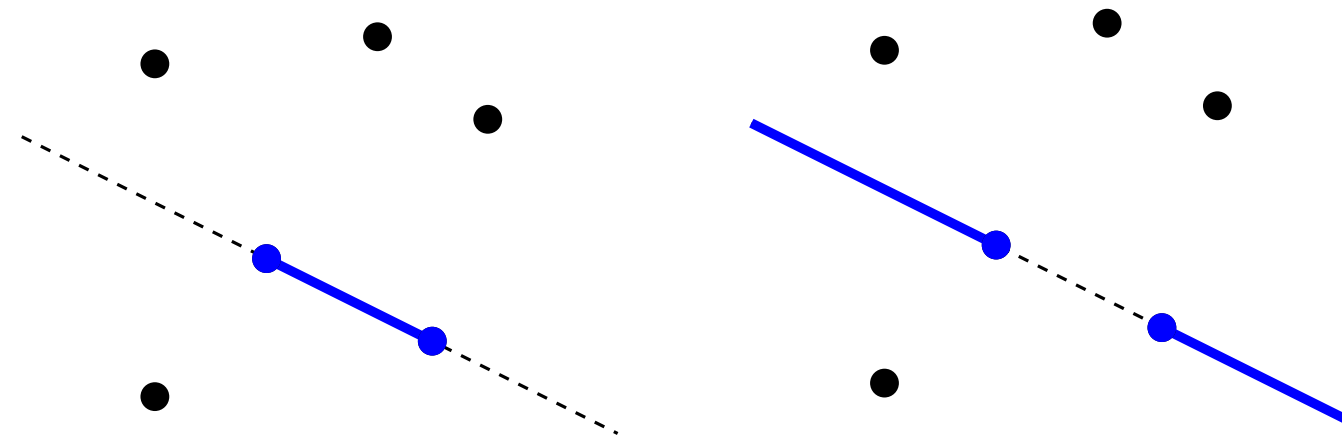
every pair of points  $p, q$  spans two *projective segments*:  
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$C \subseteq \mathbb{RP}^2$  is **projectively convex** if, for every pair  $p, q \in C$ ,  
one of its projective segments is fully contained in  $C$



# Projective Convex Hull, Gons and Holes

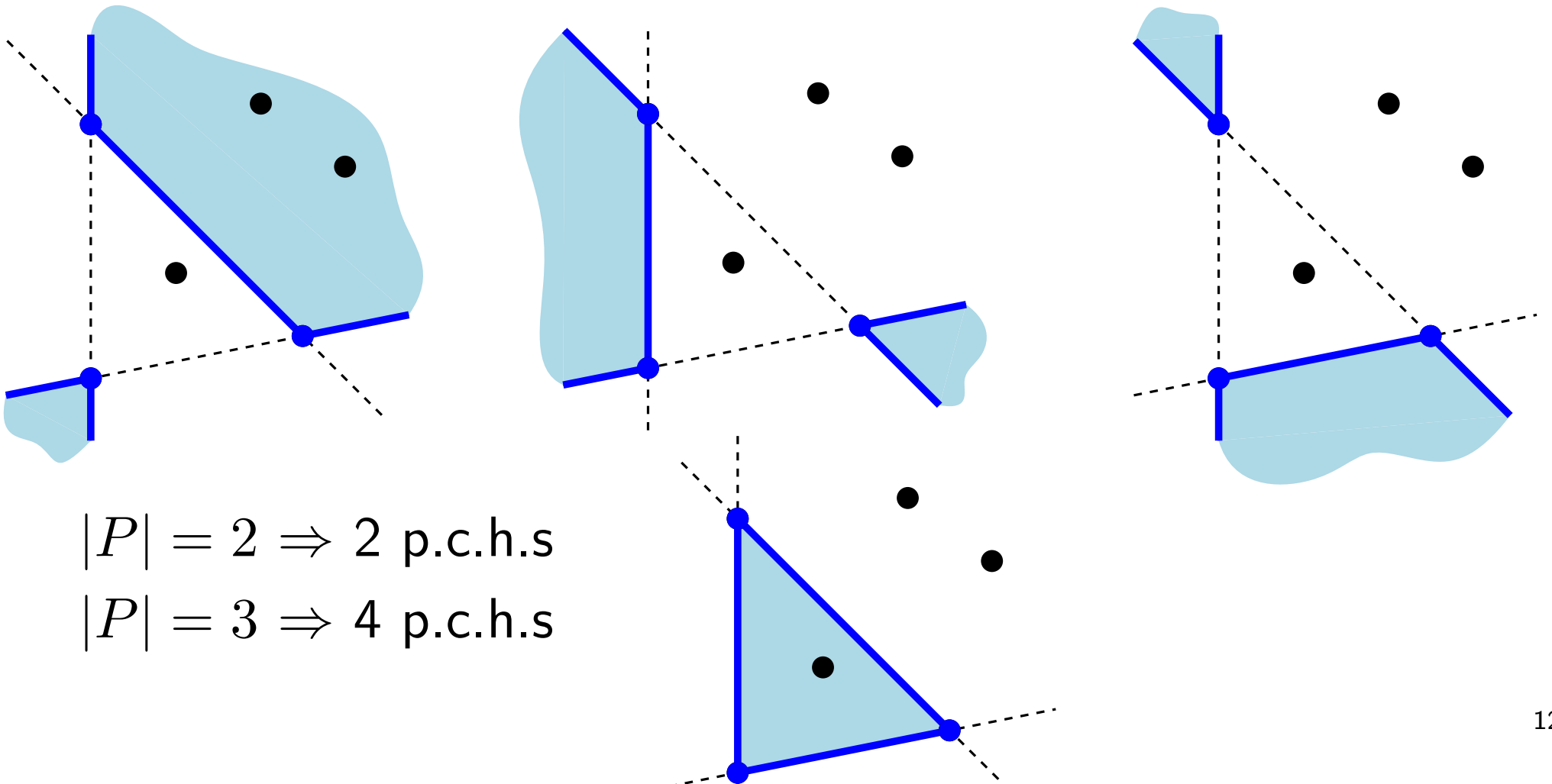
**projective convex hull:** a inclusion-wise minimal convex set containing a given set  $P$  (not unique)



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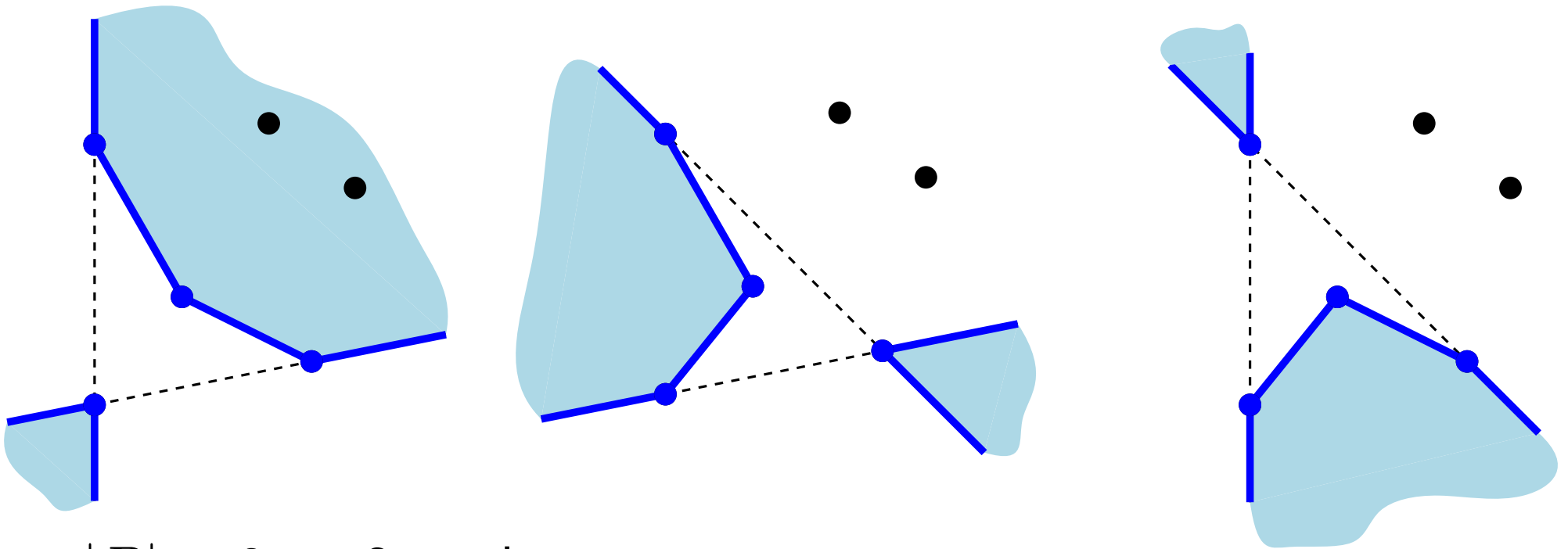


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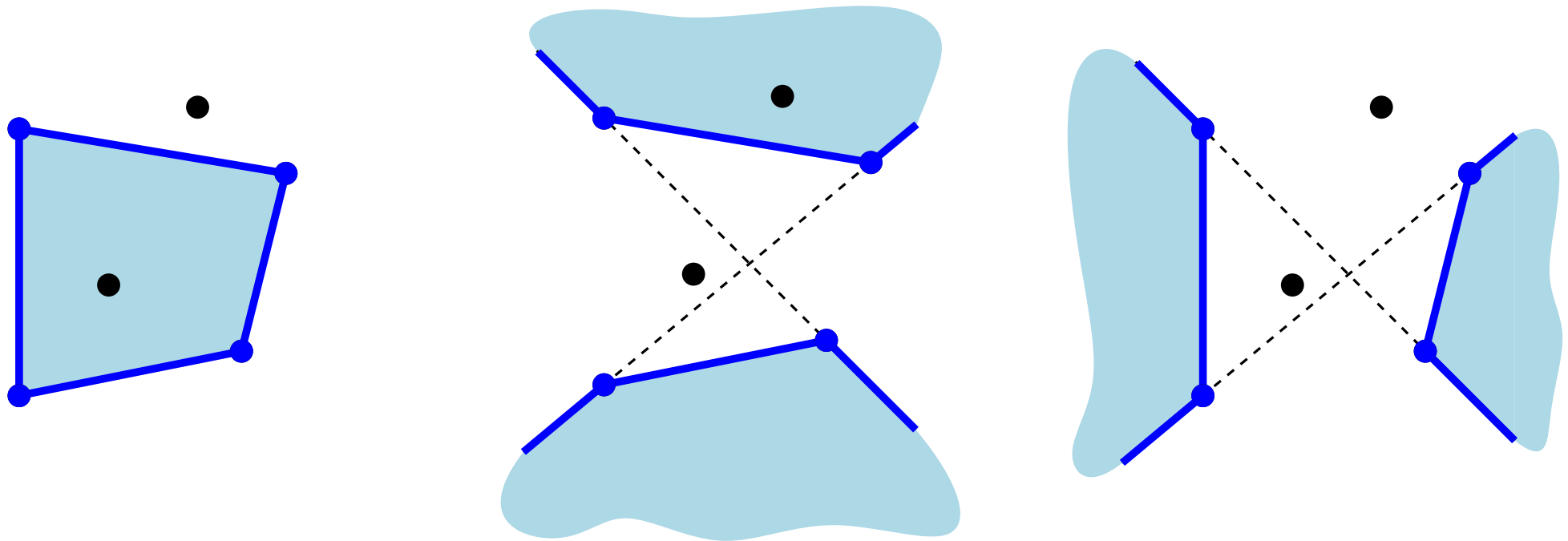
$$|P| = 2 \Rightarrow 2 \text{ p.c.h.s}$$

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$$|P| = 4 \Rightarrow 3 \text{ p.c.h.s (with all points on boundary)}$$

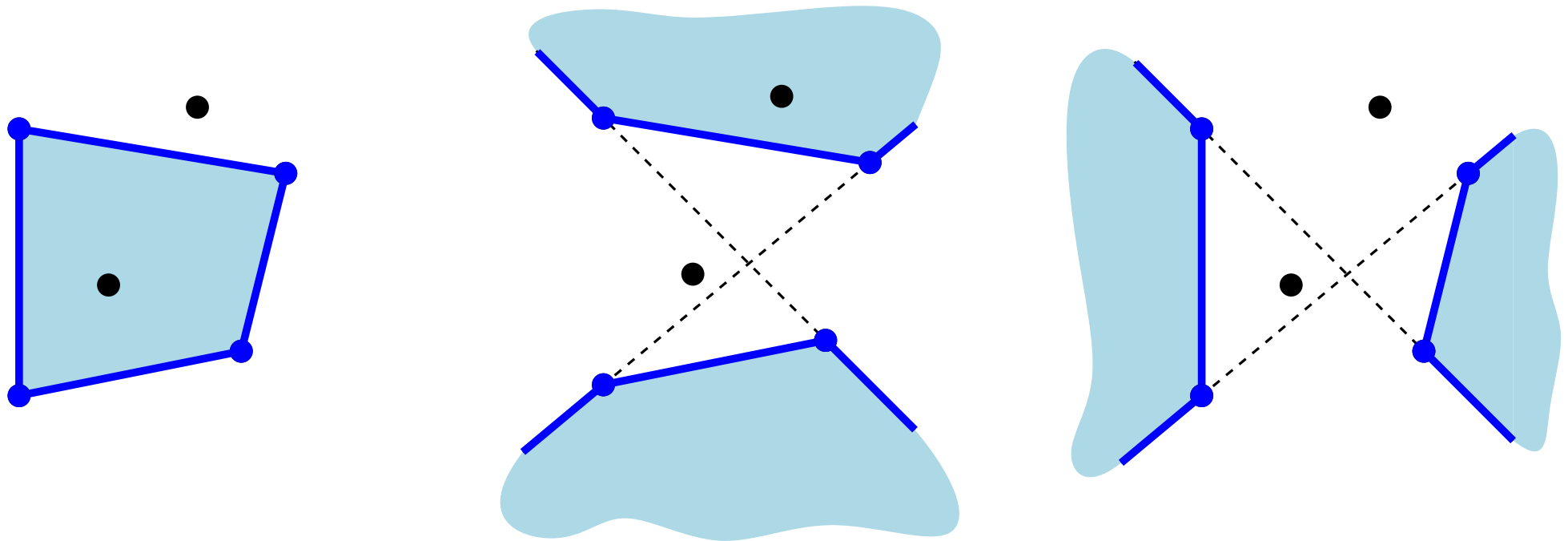
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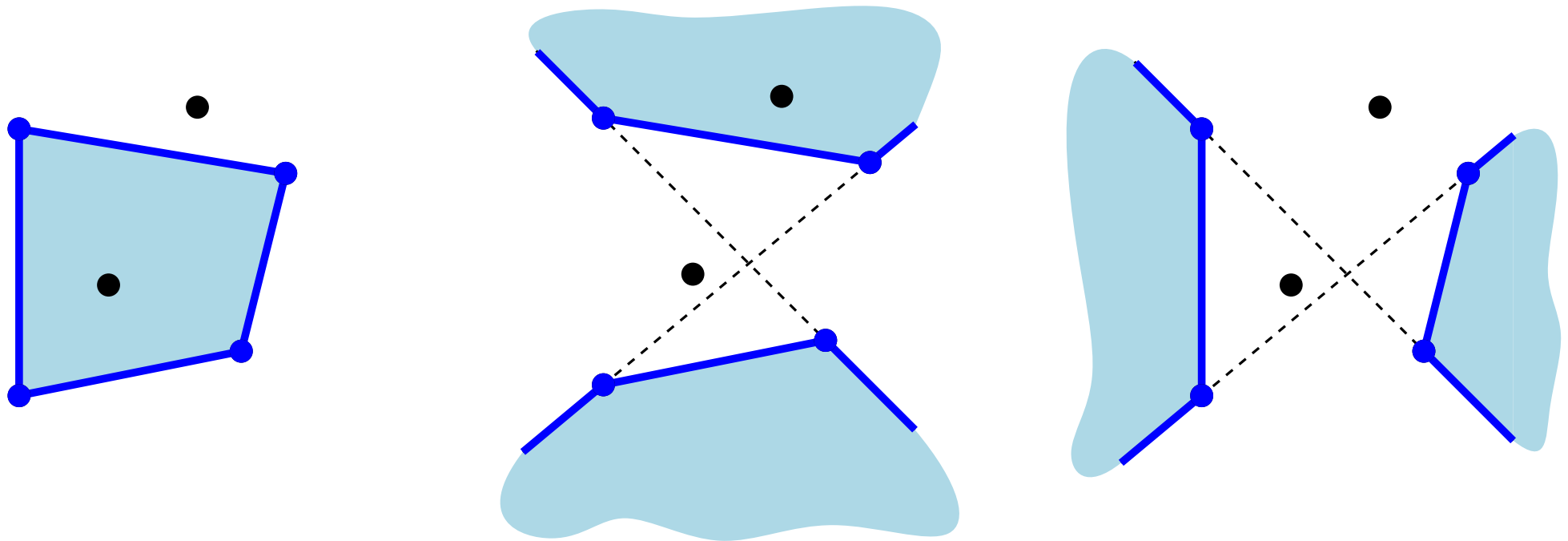


**projective  $k$ -gon:** projective convex set spanned by  $k$  pnts  
(first introduced by Harborth & Möller '93)



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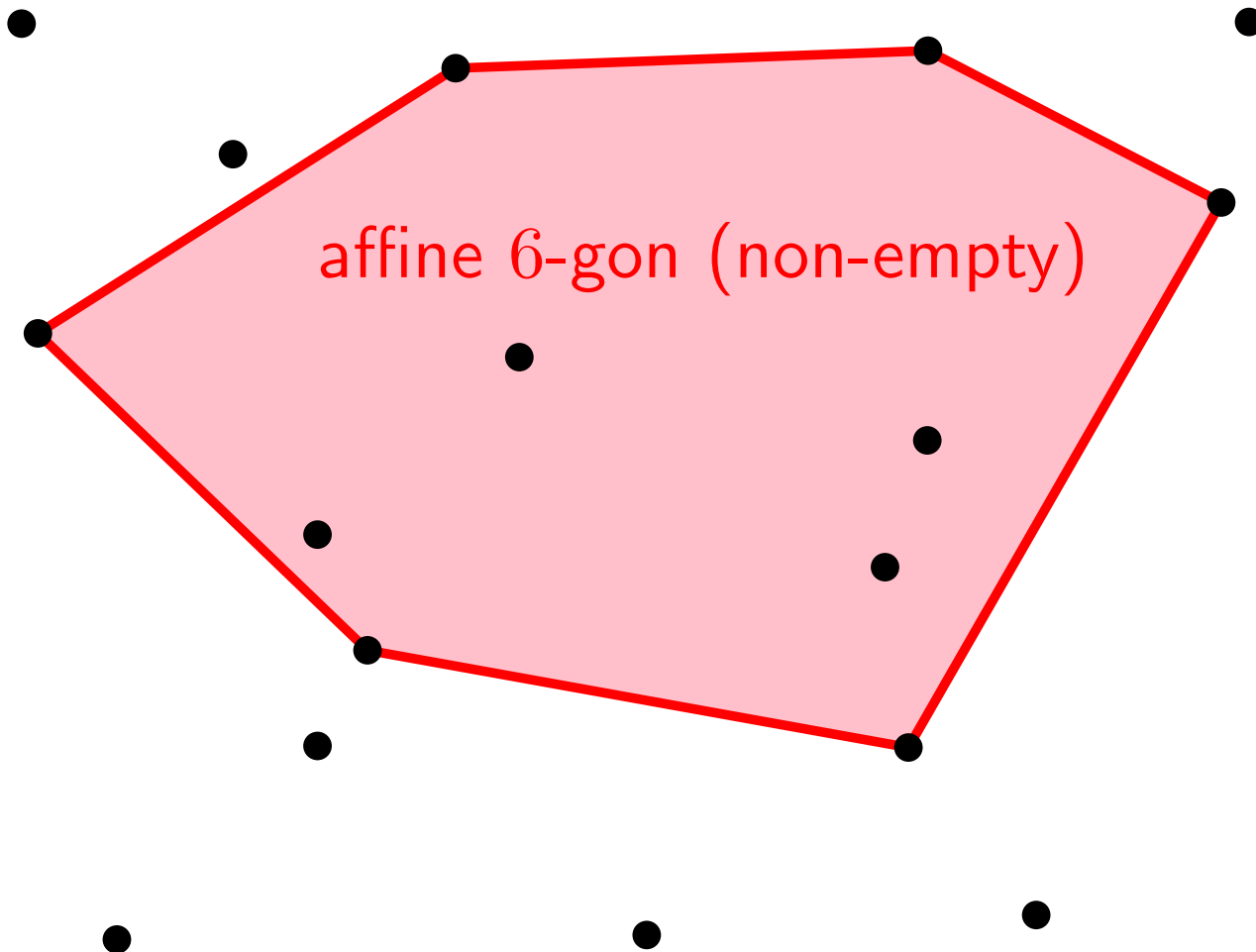


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**projective  $k$ -hole:**  $k$ -gon containing no other point of  $P$

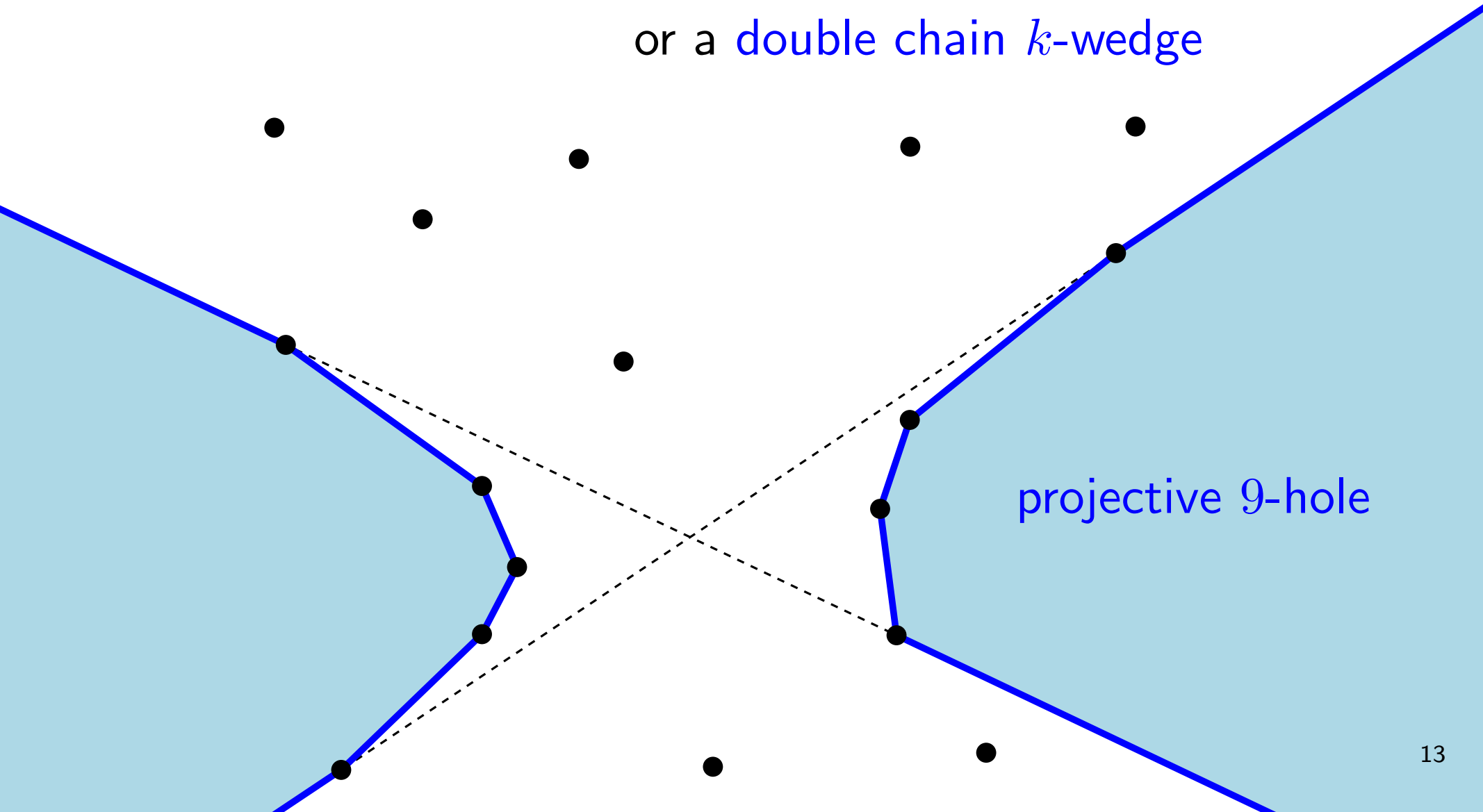
# Projective Gons and Holes

a projective  $k$ -gon is either an **affine  $k$ -gon**



# Projective Gons and Holes

a projective  $k$ -gon is either an **affine  $k$ -gon**  
or a **double chain  $k$ -wedge**





# Erdős–Szekeres Numbers

**Thm** (affine  $k$ -gons; Erdős & Szekeres '35; Suk '16; Holmsen, Mojarrad, Pach & Tardos '17):

$$2^{k-2} + 1 \leq g(k) \leq 2^{k+O(\sqrt{k \log k})}$$

**Thm** (projective  $k$ -gons, BSV '22):

$$2^{k-O(\log k)} \leq g^p(k) \leq 2^{k+O(\sqrt{k \log k})}$$

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## $k$ -Holes in Horton Sets

**Thm** (affine holes, Horton '83, Bárány & Füredi '87):

Let  $S$  be a Horton set of size  $n = 2^t$ . Then

$h_k(S) \leq O(n^2)$  for  $k \leq 6$  and  $h_7(S) = 0$ .

( $h_k \dots \#$  of affine  $k$ -holes)

**Thm** (projective holes, BSV '22):

Let  $S$  be a Horton set of size  $n = 2^t$ . Then

$h_k^p(S) \leq O(n^2)$  for  $k \leq 7$  and  $h_8^p(S) = 0$ .

( $h_k^p \dots \#$  of projective  $k$ -holes)



## $k$ -Holes Summarized

$k = 3$ : always exist

$$k = 4: h_4(\geq 5) \geq 1 \longrightarrow h_4^p(\geq 4) \geq 1$$

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Harborth '78

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weird behavior

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$k = 6$ :

$$h_6(\geq g(9)) \geq 1$$

$$\longrightarrow h_6^p(\geq g^p(9)) \geq 1$$

Gerken '08



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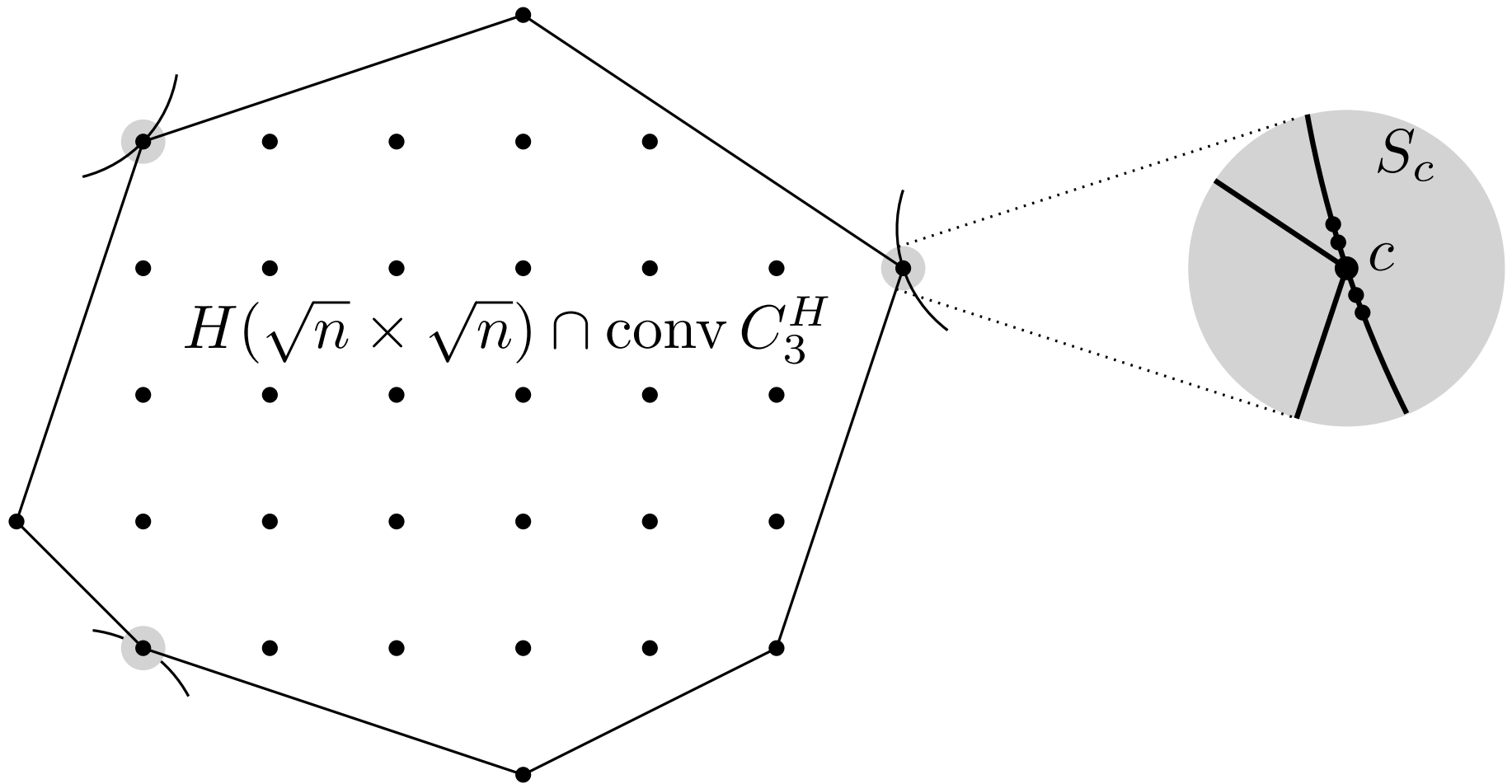
$k \geq 8$ : not exist (Horton sets)

$k = 7$ : no affine, but projective 7-holes in Horton sets.

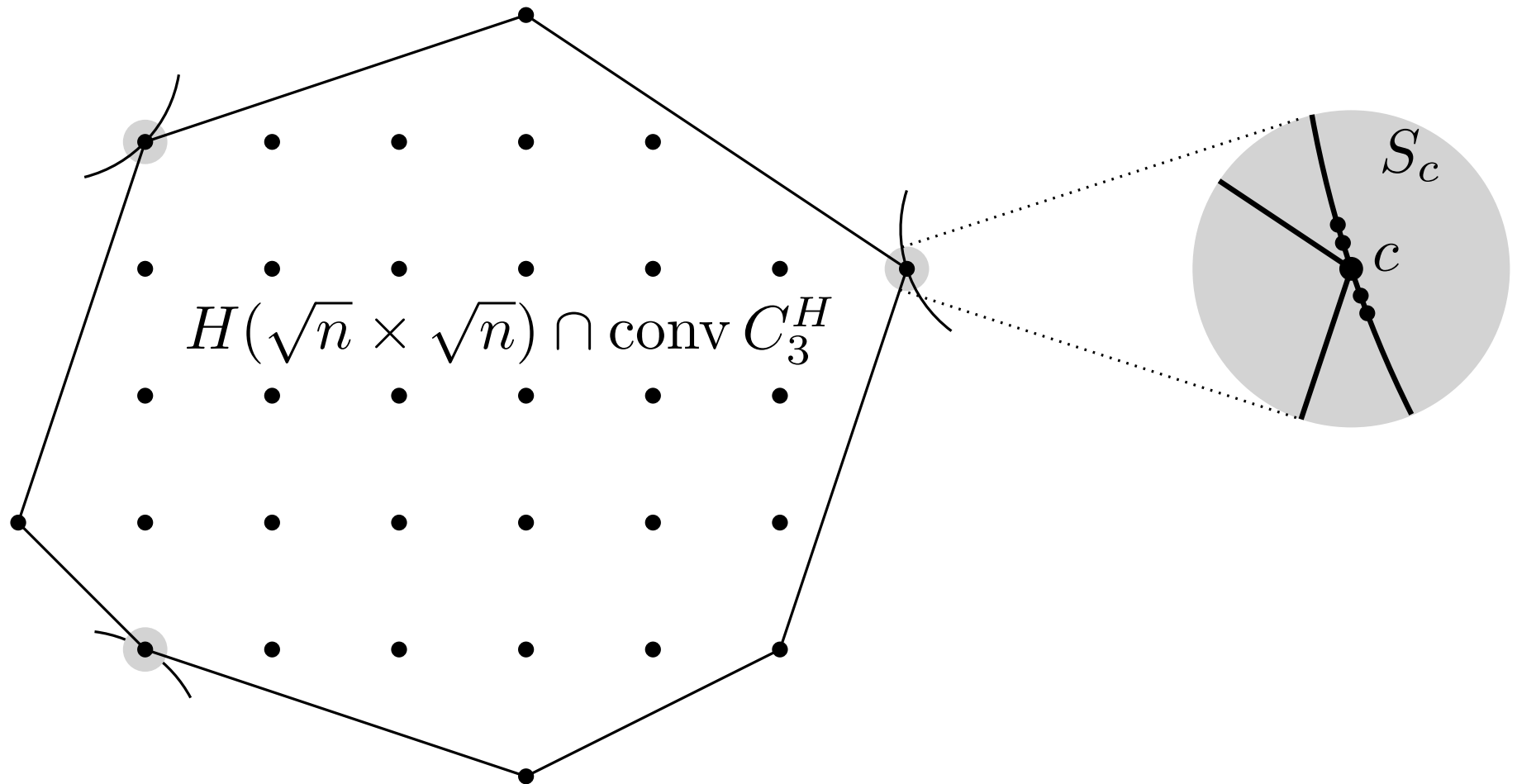
**existence of projective 7-holes remains open!**

Any significant difference?

# Substantially more Projective Holes



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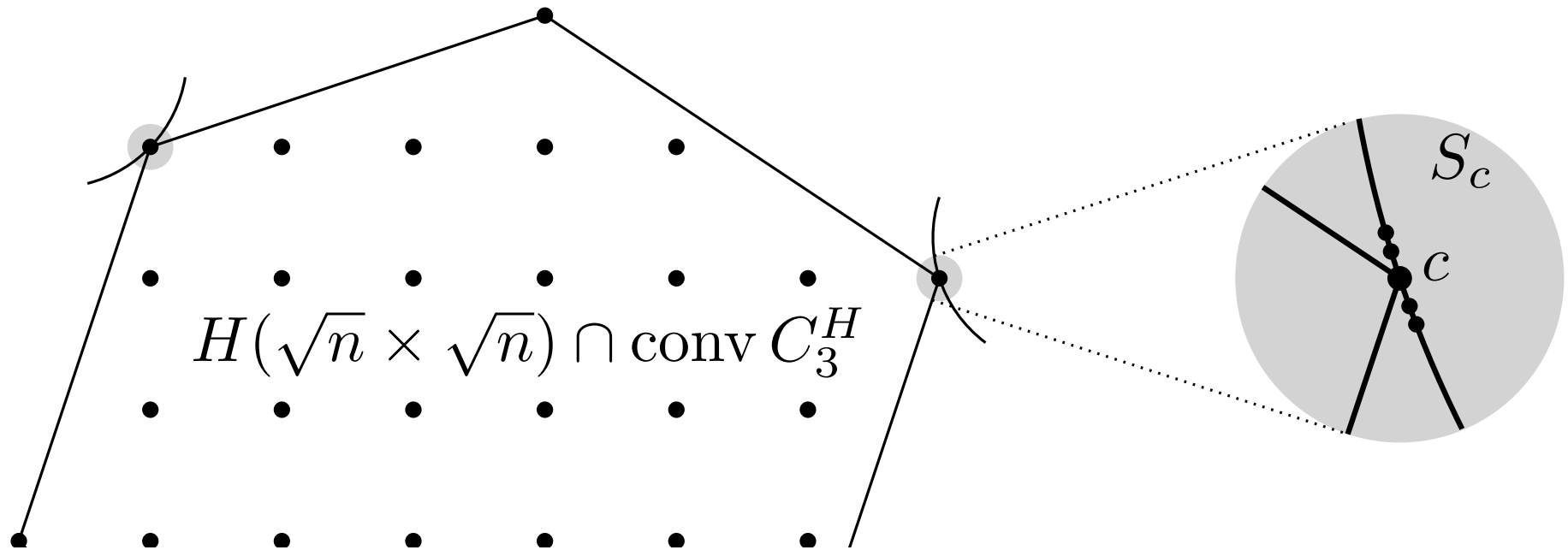


$$H(\sqrt{n} \times \sqrt{n}) \cap \text{conv } C_3^H$$

**Thm** (BSV '22).  $\forall k \in \{3, \dots, 6\}$  and  $n$ ,  $\exists n$ -point set with  $O(n^2)$  affine  $k$ -holes and  $\Omega(n^{3 - \frac{5}{3k}})$  projective  $k$ -holes.



# Substantially more Projective Holes



$$H(\sqrt{n} \times \sqrt{n}) \cap \text{conv } C_3^H$$

**Thm** (BSV '22)  $\forall n$  and  $x \leq 2^{n/2} \exists n$ -point set with  $O(x + n^2)$  affine holes and  $\Omega(x^2)$  projective holes.

**Thm** (BSV '22).  $\forall k \in \{3, \dots, 6\}$  and  $n$ ,  $\exists n$ -point set with  $O(n^2)$  affine  $k$ -holes and  $\Omega(n^{3 - \frac{5}{3k}})$  projective  $k$ -holes.

# Further Results

# Holes in Random Point Sets

Affine:

$$EH_3 \sim 2n^2 \text{ [Valtr '95, Reitzner \& Temesvari '19]}$$

$$EH_k = \Theta(n^2) \text{ [BSV '19+'21]}$$

Projective [BSV '22]:

$$EH_3^p = \Theta(n^2) \text{ with larger multiplicative constant}$$

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no proof for larger holes, but  $\Theta(n^2)$  conjectured!

## Algorithmic Aspects

**Thm** (Mitchell, Rote, Sundaram & Woeginger '95).

The number of affine  $k$ -gons and  $k$ -holes in an  $n$ -point set can be computed in  $O(kn^3)$  time and  $O(kn^2)$  space.

**Thm** (BSV '22).

The number of projective  $k$ -gons and  $k$ -holes in an  $n$ -point set can be computed in  $O(kn^4)$  time and  $O(kn^2)$  space.

