

# Erdős–Szekeres–type Problems on Planar Point Sets

Manfred Scheucher

#### **General Position**

a finite point set S in the plane is in general position if  $\nexists$  collinear points in S



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throughout this presentation, every set is in general position

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**Theorem (Erdős & Szekeres 1935).**  $\forall k \in \mathbb{N}, \exists a \text{ smallest integer } g(k) \text{ such that}$ every set of g(k) points contains a k-gon.

# **Theorem (Erdős & Szekeres '35)** $2^{k-2} + 1 \le g(k) \le \binom{2k-4}{k-2} + 1$

equality conjectured by Szekeres, Erdős offered 500\$ for a proof

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several improvements of order  $4^{k-o(k)}$ 

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$$g(k) \le 2^{k+O(k^{2/3}\log k)}$$
 [Suk '16]

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Known: 
$$g(4) = 5$$
,  $g(5) = 9$ ,  $g(6) = 17$ 

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computer assisted proof, 1500 CPU hours [Szekeres–Peters '06]

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**Theorem.**  $g(k) \le 2^{k+o(k)}$ . [Suk '16] < 1 hour using SAT solvers [S.'18, Marić '19] Known: g(4) = 5, g(5) = 9,  $g(6) \stackrel{\checkmark}{=} 17$ 

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using a = b = k, we then get  $g(k) \leq \binom{2k-4}{k-2} + 1$ .

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Base Case: If a = 2 or b = 2, we have  $\phi(a, b) + 1 = {\geq 0 \choose 0} + 1 = 2$  points  $\Rightarrow$  2-cup / 2-cap

Let S be a set of  $\phi(a, b) + 1 = {a+b-4 \choose a-2} + 1$  points and suppose it does not contain a *a*-cup. We show that there is a *b*-cap.

Step: Let E be the set of rightmost (end)points of all (a-1)-cups

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Step: Let E be the set of rightmost (end)points of all (a-1)-cups  $\Rightarrow \exists b$ -cap! Q.E.D.

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this bound is actually tight because there exist sets  $S_{a,b}$ with  $\phi(a,b) = {a+b-4 \choose a-2}$  points without *a*-cups and *b*-caps.

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- this will give

$$|S_{a,b}| = \phi(a,b) = \phi(a,b-1) + \phi(a-1,b) = \binom{a+b-4}{a-2,b-2}$$

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• no *a*-cup or *b*-cap

• Q.E.D.

- $=S_{a-1,b}$

 $=S_{a,b-1}$
• use  $S_{2,k}$ ,  $S_{3,k-1}$ , ...,  $S_{k,2}$  as gadgets to construct a large set without k-gons

• place  $S_{a,b}$ 's very flat in very small bubbles











• place  $S_{a,b}$ 's very flat in very small bubbles (high above/deep below)



 $S_{k,2}$ 



- place  $S_{a,b}$ 's very flat in very small bubbles
- bubbles can have arbitrary relative positions







 $S_{k,2}$ 



$$S_{k,2}$$

$$S_{k-1,3}$$

$$S_{k-2,4}$$

$$S_{k-3,5}$$

$$S_{2,k}$$

• each  $\ell$ -gon has:

 $(\leq t)$ -cap in top-layer t,  $(\leq k - b)$ -cup in bottom-layer b,  $\leq 1$  point per interm. layer  $\Rightarrow \ell < k$ 

$$S_{k,2}$$

$$S_{k-1,3}$$

$$t = 2$$

$$S_{k-2,4}$$

$$b = 4$$

$$S_{2,k}$$

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• number of points:  $\phi(k,2) + \phi(k-1,3) + \ldots = \sum_{j=0}^{k-2} {k-2 \choose j} = 2^{k-2}$ 

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- number of points:  $\phi(k, 2) + \phi(k - 1, 3) + \ldots = \sum_{j=0}^{k-2} {k-2 \choose j} = 2^{k-2}$
- therefore  $g(k) \ge 2^{k-2} + 1$

$$S_{k,2}$$
 $S_{k-1,3}$ 
 $S_{k-2,4}$ 
 $S_{k-3,5}$ 
 $b = 4$ 



**Lemma** (Fractional EST, Pór & Valtr '02) Let S be a point set with  $|S| \ge 2^{32k}$  points. Then there exists k-cup/cap  $X \subset S$  satisfying  $|T_i \cap S| \ge \frac{|S|}{2^{32k}}$  for every i.



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linear in S!



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- auxiliary poset on each  $T_i$



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• ... which combine to k-gons

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n points in convex position

any 
$$k$$
-subset is  $k$ -gon

$$\Rightarrow \max = \binom{n}{k}$$

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• 
$$g_k(n) = \Theta(n^k)$$
 each  $g(k)$ -subset gives  $k$ -gon  
because  $g_k(n) \ge \frac{\binom{n}{g(k)}}{\binom{n-k}{g(k)-k}} \ge c \cdot \frac{n^{g(k)}}{n^{g(k)-k}} = c \cdot n^k$ 

each k-gon counted by at most that many g(k)-subsets

 $g_k(n) :=$  minimum # of k-gons among all sets of n points

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• k = 4: rectilinear crossing number of  $K_n$ :  $g_4(n) = \overline{cr}(K_n) \sim c_4 \cdot {n \choose 4}$  with  $0.3799 < c_4 < 0.3805$ [Ábrego et al. '08, Aichholzer et al. '20]



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- various notions of crossing numbers have been studied intensively (not necessarily straight-line drawings, not necessarily complete graphs)

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$$g_4(S) + \sum_{k=0}^{\lfloor n/2 \rfloor - 1} k(n - 2 - k) e_k(S) = 3\binom{n}{4}$$

$$e_0 = 6$$
  
 $e_1 = 6$   
 $e_2 = 9$ 

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**Corollary:** Given a set S of n points, the value  $g_4(S)$ (which is of order  $\Theta(n^4)$ ) can be computed in  $O(n^2)$  time

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Remark:  $g_k(S)$  can be computed in  $O(k \cdot n^3)$  time [Mitchell Rote Sundaram Woeginger '95]

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Proof Idea: 
$$\sum_{e} k_e \cdot (n-2-k_e) = 2 \boxtimes + 3 \bigtriangleup$$



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 QED



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Remark: this crossings – k-edges relation generalizes to simple topological drawings of  $K_n$ [Ábrego, Fernández-Merchant, Ramos, Salazar '11]

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- 10 points  $\Rightarrow \exists$  5-hole [Harborth '78]

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- ∃ arbitrarily large point sets with no 7-hole [Horton '83]
- Sufficiently large point sets ⇒ ∃ 6-hole
  [Gerken '08 and Nicolás '07, independently]







4 points, no 4-hole (h(4) = 5)





29 points, no 6-hole [Overmars '02] ( $30 \le h(6) \le 463$ )



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# Horton's construction for $n=2^1$ points, no 7-holes



Horton's construction for  $n=2^2$  points, no 7-holes



Horton's construction for  $n=2^3$  points, no 7-holes





Horton's construction for  $n=2^4\,$  points, no 7-holes



Horton's construction for  $n=2^4$  points, no 7-holes


Horton's construction:  $n = 2^k$  points, no 7-holes  $(h(7) = \infty)$ 

- $S_1 = \text{single point}$ , and  $S_n$  recursively: two copies L, U of  $S_{\frac{n}{2}}$ 
  - from left to right: points alternatingly in L and U
  - U high above L



 $S_1 = \text{single point, and } S_n$  recursively: two copies L, U of  $S_{\frac{n}{2}}$ 

• from left to right: points alternatingly in L and U



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• we show  $h(5) \leq g(6)$ :

consider 6-gon G with minimum number of interior points I



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if I is empty, G is a 6-hole



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if I is empty, G is a 6-hole if I contains 1 point, we find a 5-hole if I contains  $\geq 2$  points, we choose ab as bounding edge of  $\operatorname{conv}(I)$ ,



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Q1:  $\exists$  shorter (computed assisted?) proof for existence of 6-holes? Q2: is h(6) bounded in terms of max{const, g(7)}?

Thm (affine 1st moment, Edelman and Jamison '85):  $\underbrace{h_0(P)}_{1} - \underbrace{h_1(P)}_{n} + \underbrace{h_2(P)}_{\binom{n}{2}} - h_3(P) \pm \ldots = \sum_{k \ge 0} (-1)^k h_k(P) = 0$ 

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Proof Idea: holds for n points in convex position,



any k-subset is k-hole

$$\sum_{k\geq 0} (-1)^k \binom{n}{k} = (1+(-1))^n = 0$$
  
**Binomial theorem**

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and invariant to mutations! (i.e., when a point moves over a line)



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Proof Idea: holds for n points in convex position,

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(i.e., when a point moves over a line)

for each k-hole which we get/destroy with ab we also get/destroy a (k + 1)-hole with abc

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Idea:  $k \cdot h_k(P) = \sum_e h_k(P;e)$ 



• 
$$h_3(n) \ge \lfloor \frac{1}{3} {n \choose 2} \rfloor = \Omega(n^2)$$
  
•  $\forall \text{ edge } e$   
 $\exists 3\text{-hole with closest point}$ 



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 $h_k(n) :=$  minimum # of k-holes among all sets of n points

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 $\forall$  crossed edge e

 $\exists$  4-hole with diagonal e



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O(n) uncrossed edges (planar graph)

• 
$$h_5(n) \ge \lfloor \frac{1}{10}n \rfloor = \Omega(n)$$



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 • same idea:  $h_6(n) \ge \Omega(n)$ 





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Bárány and Füredi '87]
h<sub>3</sub>, h<sub>4</sub> both in Θ(n<sup>2</sup>)
h<sub>5</sub> in Ω(n log<sup>4/5</sup> n) and O(n<sup>2</sup>)
[Aichholzer, Balko, Hackl, Kynčl, Parada, S., Valtr, and Vogtenhuber '17] (computer assisted proof, 20 pages)







Horton set  $S_n$  defined recursively: two copies L, U of  $S_{\frac{n}{2}}$ 

• from left to right: points alternatingly in L and U



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- each 6-hole is one of two types:
- entirely from U or L points from both, U and L



• each 6-hole is one of two types:

entirely from U or L

points from both, U and L

no 4 caps  $\Rightarrow$  3-cup plus 3-cap



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• similar recurrences for 3-, 4-, and 5-holes

## Horton Lattice

• Horton lattice: perturbation of  $\sqrt{n} \times \sqrt{n}$  grid



gives currently best bounds for  $h_3, h_4, h_5, h_6$  and no 7-holes [Bárány & Valtr '04]

## What about Higher Dimensions??

# **Higher Dimensions**

a finite point set P in  $\mathbb{R}^d$  is in general position if no d + 1 points lie in a common hyperplane

k-gon = k points in convex position

k-hole = k-gon with no other points of P in its convex hull

dimension reduction (Károlyi '01):



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$$g^{(d)}(k) \le g^{(d-1)}(k-1) + 1 \le \dots \le \underbrace{g^{(2)}(k-d+1) + d - 2}_{\le 2^{k+o(k)} \text{ (Suk'17)}}$$

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Károlyi and Valtr '03: 
$$g^{(d)}(k) = \Omega(c^{d-\sqrt[d-1]{k}})$$

asymptotic behavior remains unknown for  $d \geq 3$ 

central problem: determine the largest value k = H(d) such that every sufficiently large set in d-space contains a k-hole

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• H(2) = 6 because  $h^{(2)}(6) < \infty$  and  $h^{(2)}(7) = \infty$ [Gerken '07, Nicolás '07] [Horton '87]

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$$2d + 1 \le H(d) < d^{d+o(d)}$$
 [Valtr '92]

- (2d+1)-holes exist in sufficiently large sets
- $\exists d$ -dimensional Horton sets without  $d^{d+o(d)}$ -holes

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- in particular,  $7 \le H(3) \le 22$  remains open

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   [Conlon & Lim '21]
- do exponentially large holes exist?

#### Number of Holes in Higher Dimensions

• random sets give  $O(n^d)$  bounds for k-holes

Theorem (Balko S. Valtr '20 + '21). Let  $d \ge 2$  and  $k \ge d+1$ , and let K be a convex body in  $\mathbb{R}^d$ . If S is a set of n points chosen uniformly and independently at random from K, then the expected number of k-holes in S is  $\Theta(n^d)$ .

In particular:

 $\exists$  sets of n points in  $\mathbb{R}^d$  with  $O(n^d)$  many k-holes

#### Number of Holes in Higher Dimensions

• random sets give  $O(n^d)$  bounds for k-holes

• d-dimensional Horton sets with d > 2 contain  $\Omega(n^{\min\{k,2^{d-1}\}})$  many k-holes [Balko S. Valtr '20]

• no explicit construction known with  $O(n^d)$  k-holes



