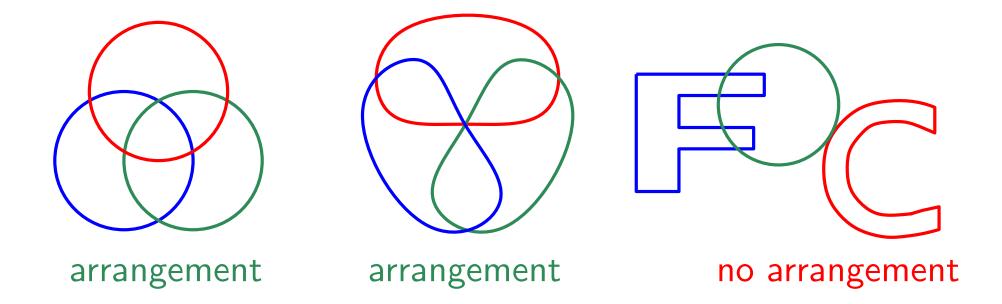


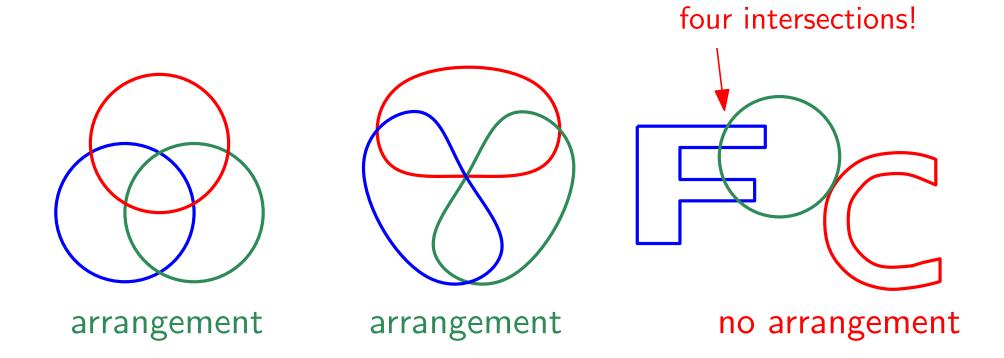
Arrangements of Pseudocircles

Stefan Felsner and Manfred Scheucher

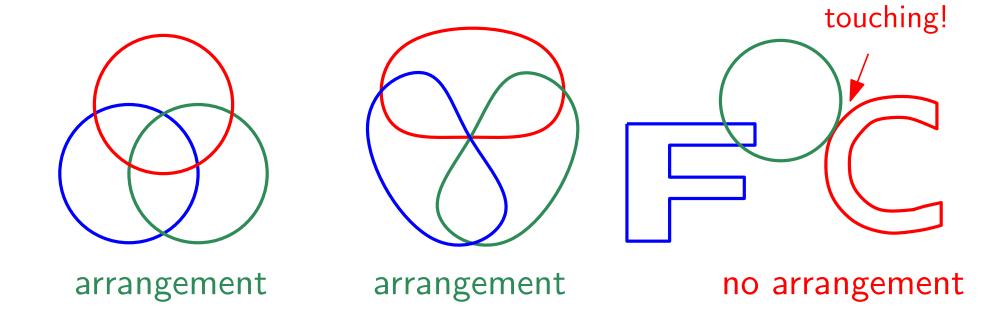
pseudocircle ...simple closed curve



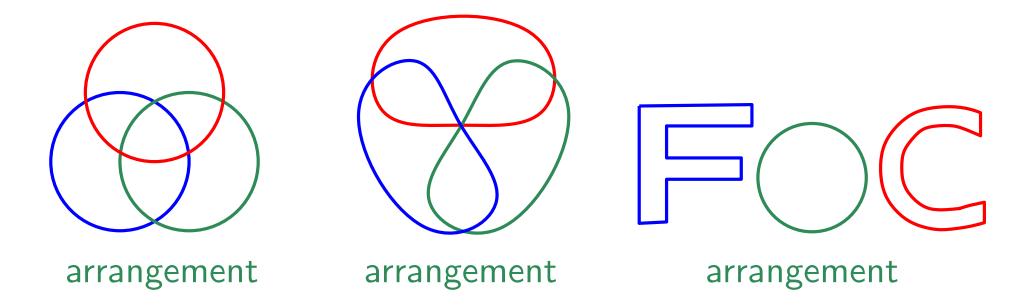
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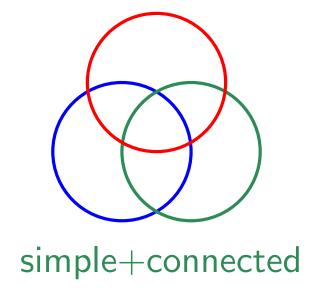
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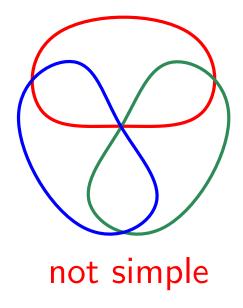


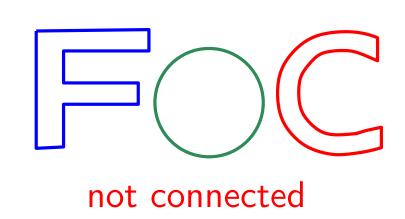
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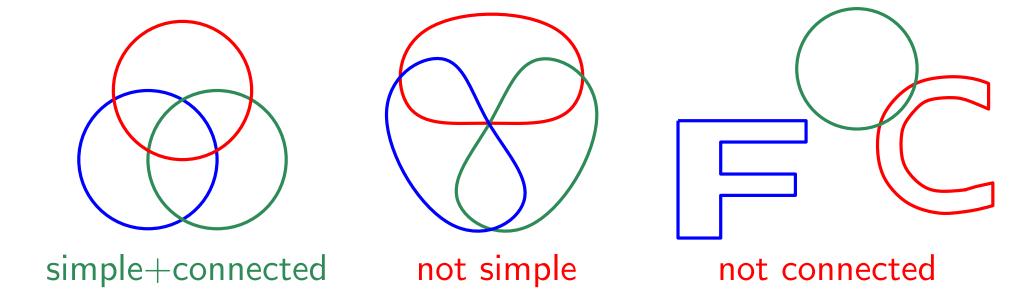
simple ... no 3 pcs. intersect in common point
connected ... intersection graph is connected





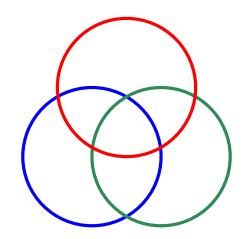


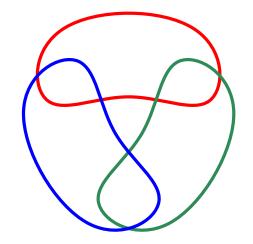
simple ... no 3 pcs. intersect in common point *connected* ... intersection graph is connected

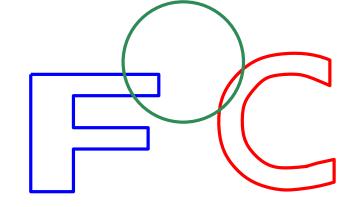


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assumptions throughout presentation

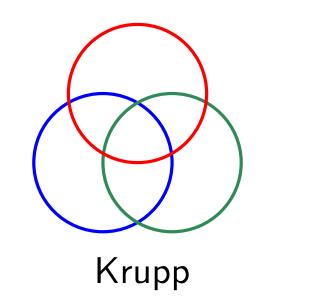




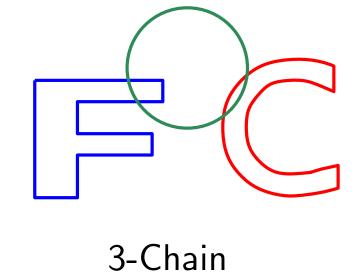


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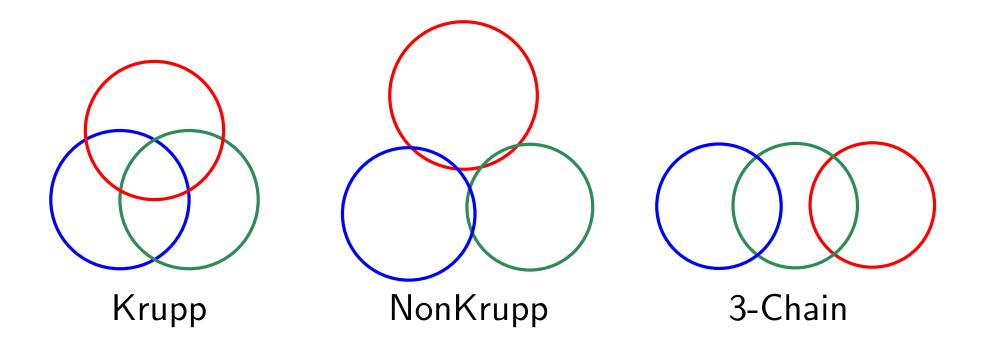




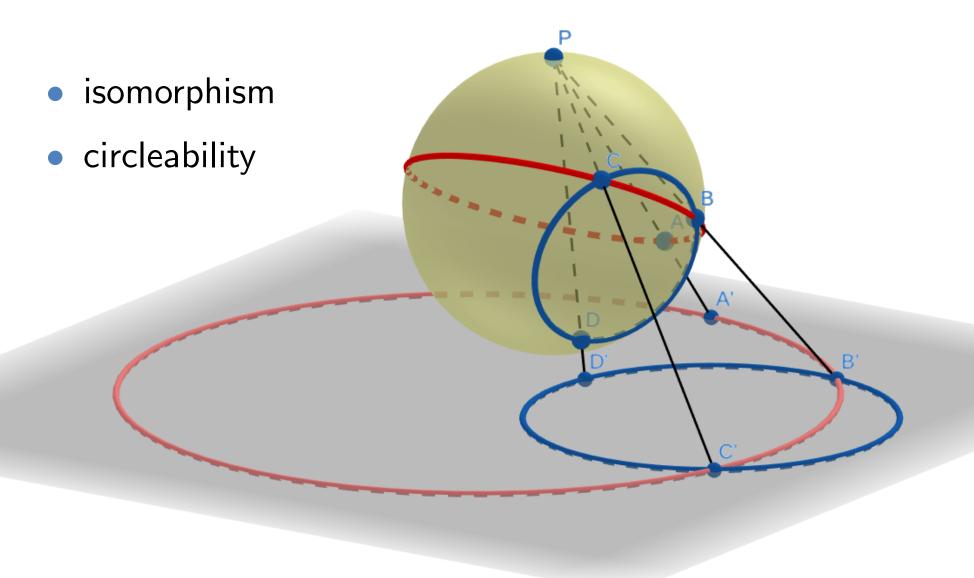
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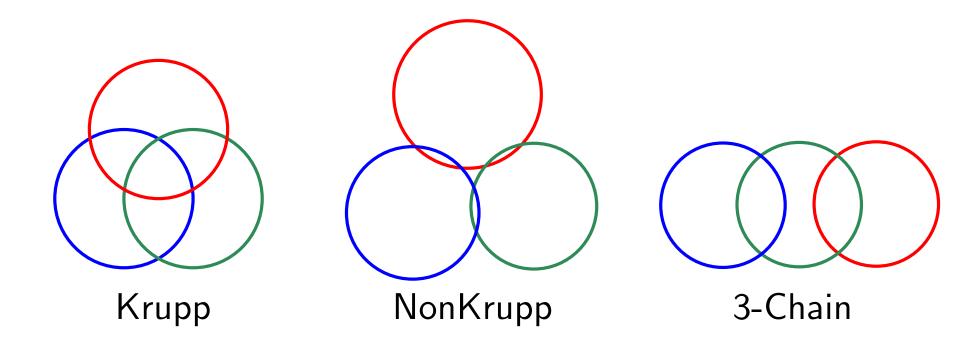
circleable ... \exists isomorphic arrangement of circles



Plane VS Sphere

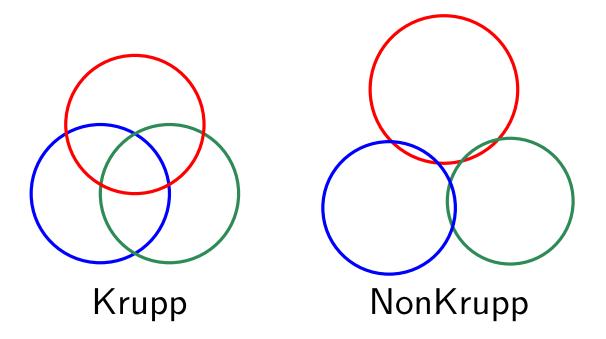


connected ... graph of arrangement is connected



connected ... graph of arrangement is connected

intersecting ...any 2 pseudocircles cross twice



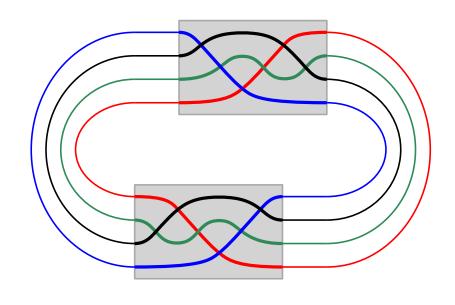
connected ... graph of arrangement is connected

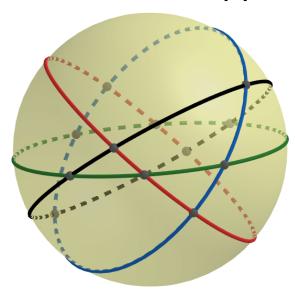


intersecting ... any 2 pseudocircles cross twice



arr. of great-pseudocircles ...any 3 pcs. form a Krupp





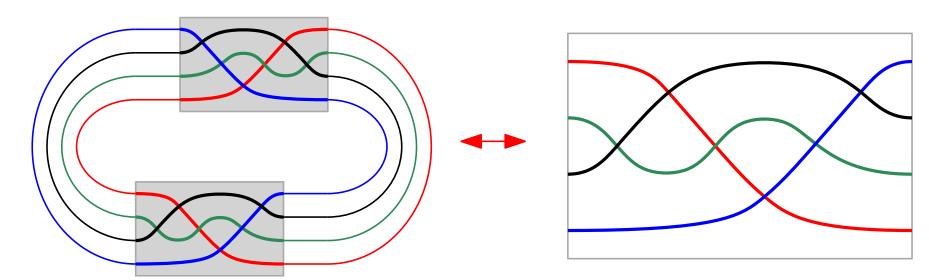
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connected ... graph of arrangement is connected intersecting ... any 2 pseudocircles cross twice arr. of great-pseudocircles ...any 3 pcs. form a Krupp digon-free ... no cell bounded by two pcs.

connected ... graph of arrangement is connected intersecting ... any 2 pseudocircles cross twice arr. of great-pseudocircles ...any 3 pcs. form a Krupp digon-free ... no cell bounded by two pcs. *cylindrical* ... \exists two cells separated by each of the pcs.

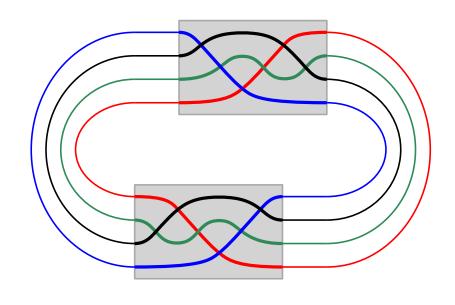
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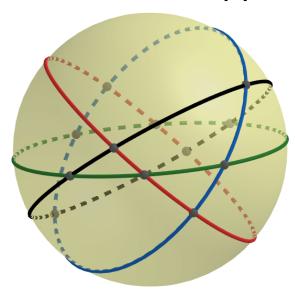


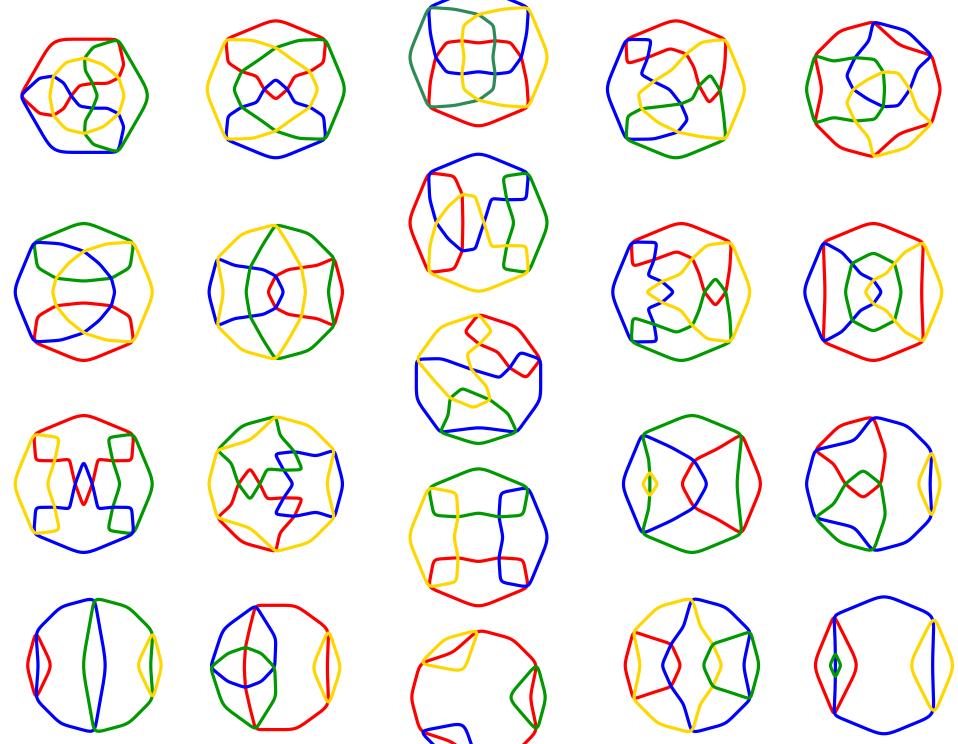
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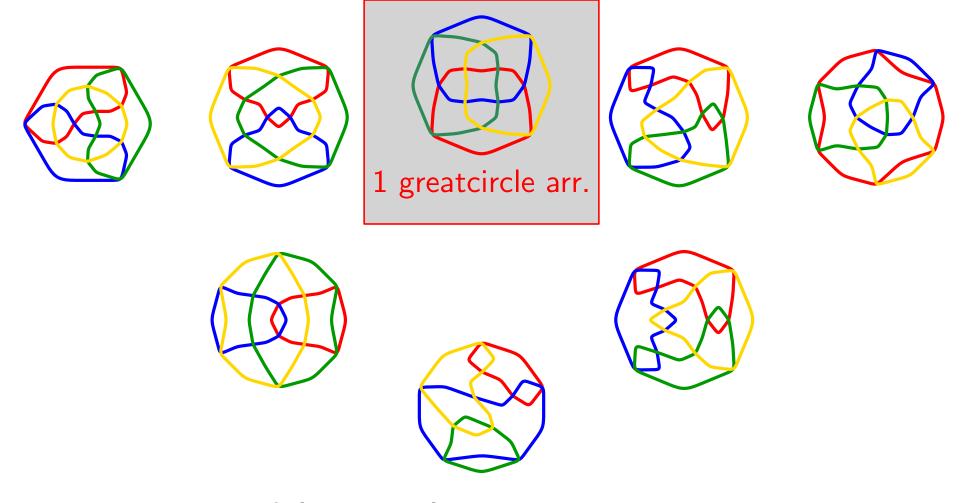


arr. of great-pseudocircles ...any 3 pcs. form a Krupp









8 intersecting arrangements

21 connected arrangements

Enumeration of Arrangements

n	3	4	5	6	7
connected	3	21	984	609 423	?
+digon-free	1	3	30	4 509	?
intersecting	2	8	278	145 058	447 905 202
+digon-free	1	2	14	2 131	3 012 972
great-p.c.s	1	1	1	4	11

Table: # of combinatorially different arragements of n pcs.

Enumeration of Arrangements

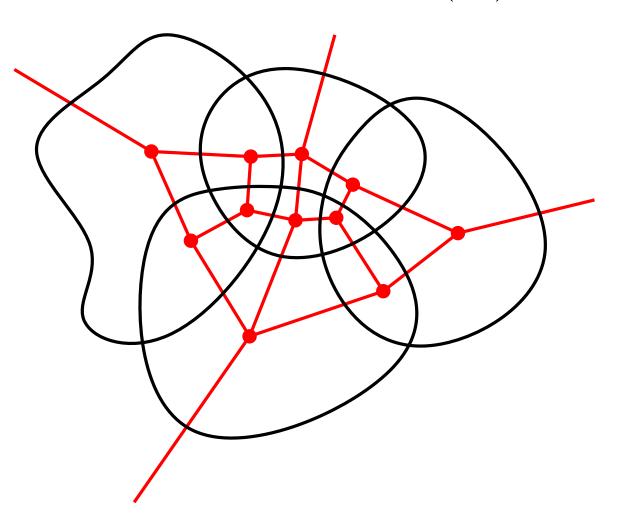
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Table: # of combinatorially different arragements of n pcs.

arrangements of pcs: $2^{\Theta(n^2)}$

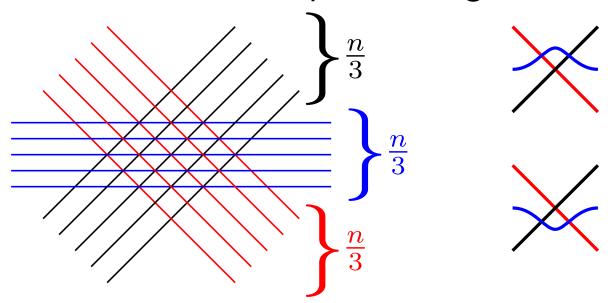
arrangements of circles: $2^{\Theta(n \log n)}$

ullet dual graph is quadrangulation on $O(n^2)$ vertices



- dual graph is quadrangulation on ${\cal O}(n^2)$ vertices
- Tutte'62: $2^{\Theta(m)}$ triangulations on m vertices
- \Rightarrow Upper bound: $2^{O(n^2)}$ non-isomorphic arrangements

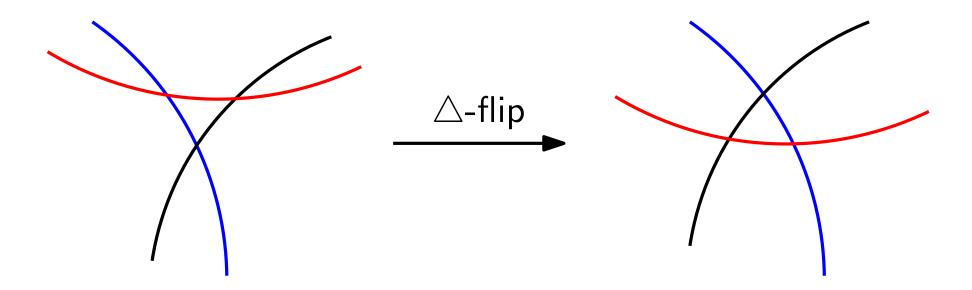
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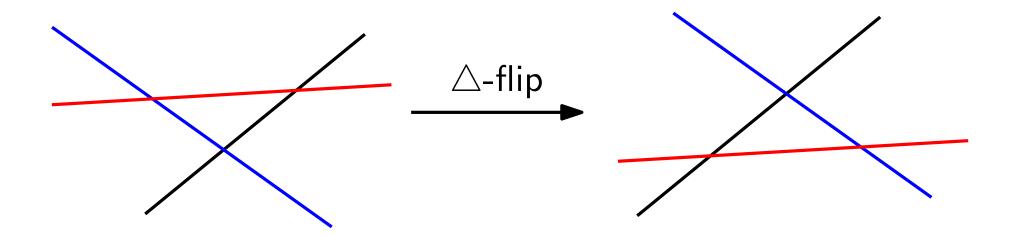
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Theorem: There are $2^{\Theta(n^2)}$ arrangements on n pcs.

Upper bound: arrangement changes if a triangle "flips"



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- we sketch the proof for *line*-arrangements

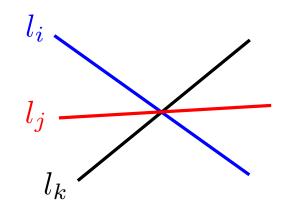


• lines l_1, \ldots, l_n given by $l_i : y_i = a_i x + b_i$

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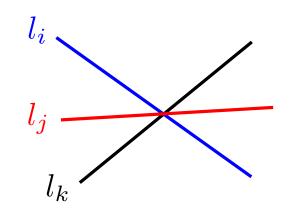
• l_i , l_j , and l_k meet in a common point

$$\iff \det \begin{pmatrix} 1 & 1 & 1 \\ a_i & a_j & a_k \\ b_i & b_j & b_k \end{pmatrix} = 0 \qquad l_j - l_j$$



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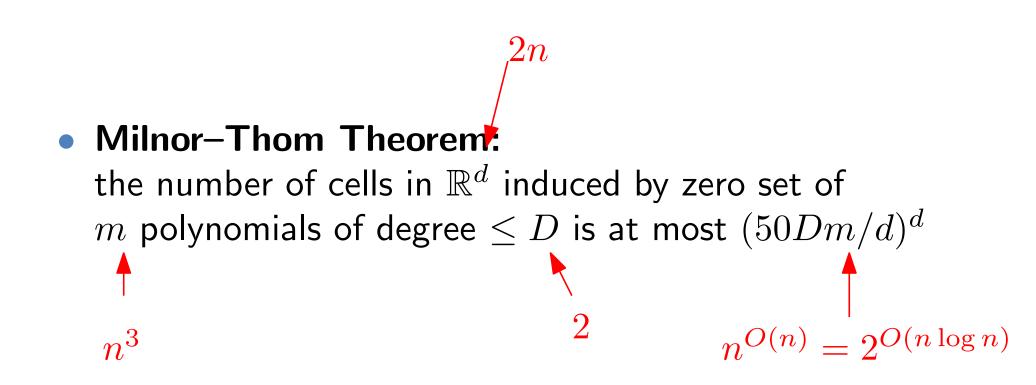
$$\iff \det \begin{pmatrix} \mathbf{1} & \mathbf{1} & \mathbf{1} \\ a_i & a_j & a_k \\ b_i & b_j & b_k \end{pmatrix} = 0 \qquad \qquad \mathbf{l}_j - \mathbf{l}_j$$



- system of $\binom{n}{3}$ quadratic polynomials in 2n variables
- simple arr. ⇔ all polynomials non-zero

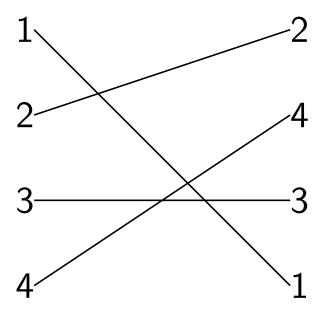
• Milnor–Thom Theorem:

the number of cells in \mathbb{R}^d induced by zero set of m polynomials of degree $\leq D$ is at most $(50Dm/d)^d$

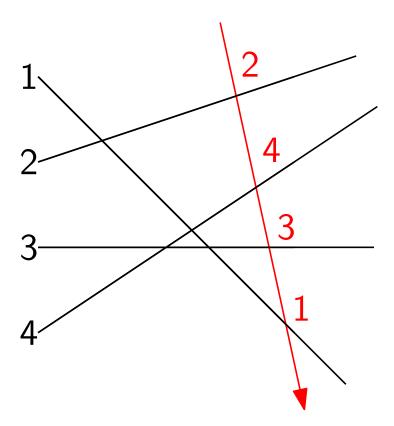


• Upper bound: $2^{O(n \log n)}$

Lower bound: # of permutations



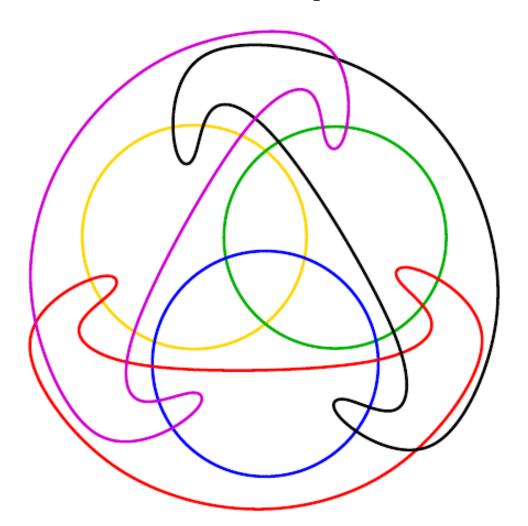
Lower bound: # of permutations



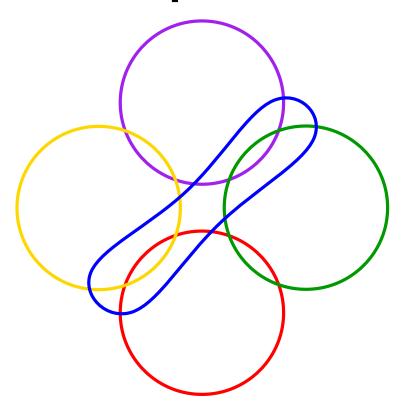
Theorem: There are $2^{\Theta(n \log n)}$ arrangements on n circles.

Part I: Circleability

• non-circleability of intersecting n=6 arrangement [Edelsbrunner and Ramos '97]



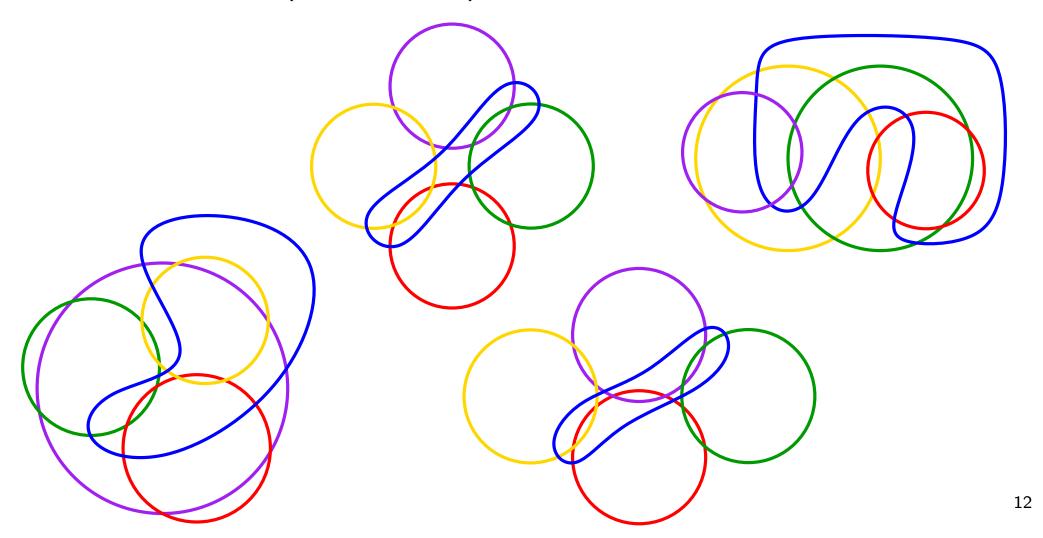
- non-circleability of intersecting n=6 arrangement [Edelsbrunner and Ramos '97]
- non-circleability of n=5 arrangement [Linhart and Ortner '05]



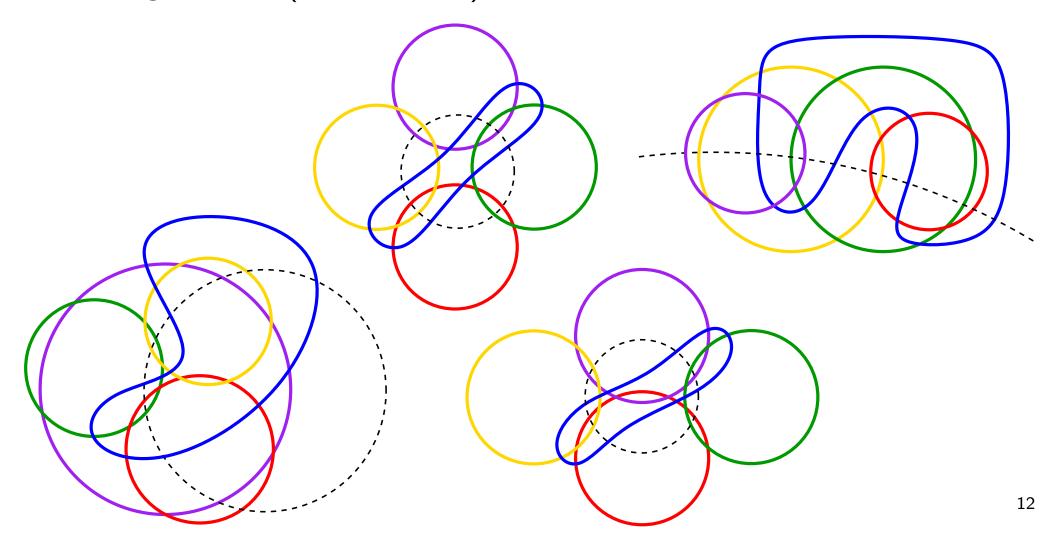
- non-circleability of intersecting n=6 arrangement [Edelsbrunner and Ramos '97]
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- circleability of all n=4 arrangements [Kang and Müller '14]

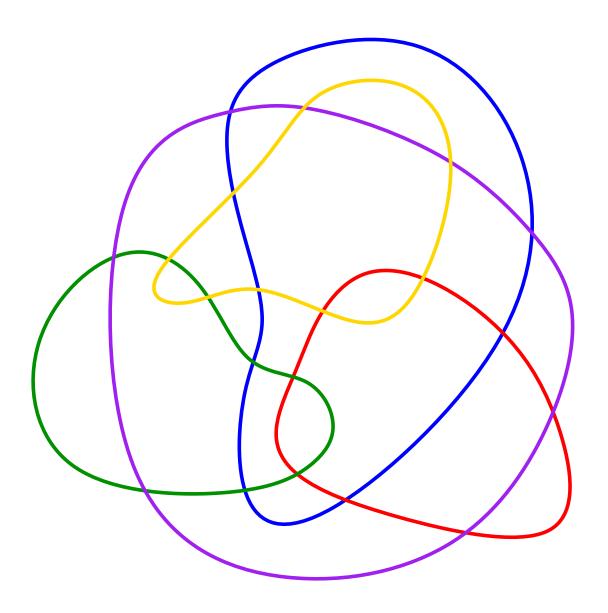
- non-circleability of intersecting n=6 arrangement [Edelsbrunner and Ramos '97]
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- NP-hardness of circleability [Kang and Müller '14]

Theorem. There are exactly 4 non-circleable n=5 arrangements (984 classes).

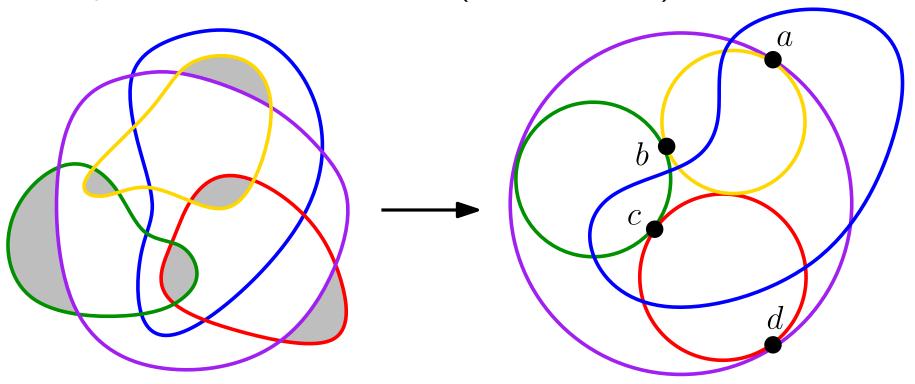


Theorem. There are exactly 4 non-circleable n=5 arrangements (984 classes).

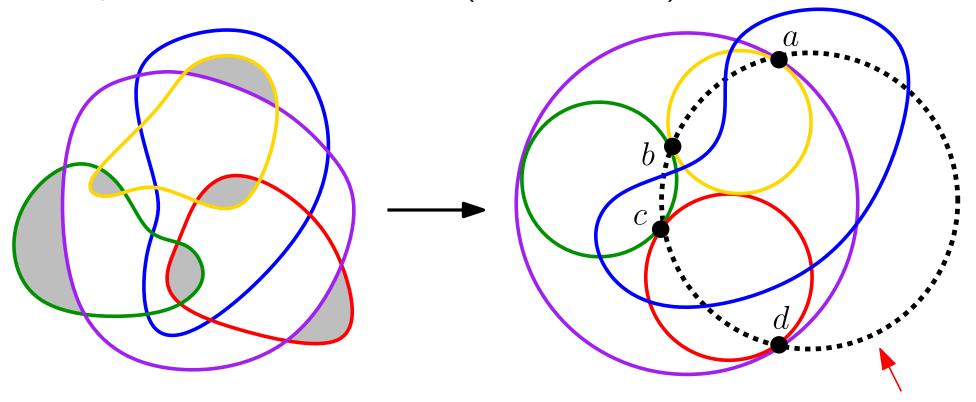




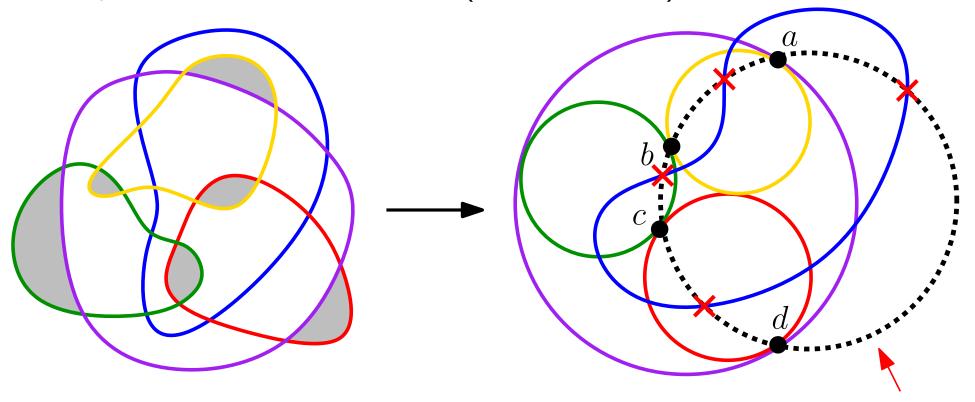
- ullet assume there is a circle representation of \mathcal{N}_5^1
- shrink the yellow, green, and red circle
- cyclic order is preserved (also for blue)



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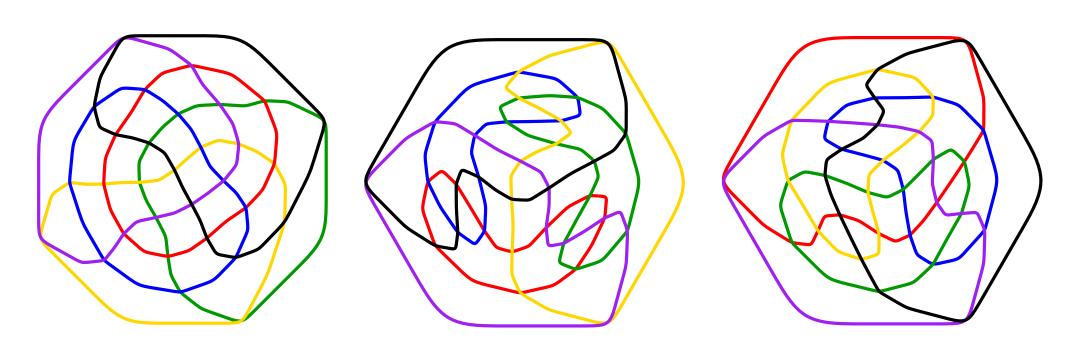


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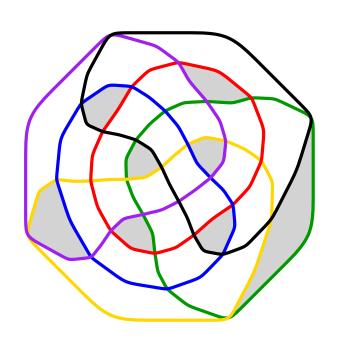


contradiction: 4 crossings

Theorem. There are exactly 3 non-circleable digon-free intersecting n=6 arrangements (2131 classes).

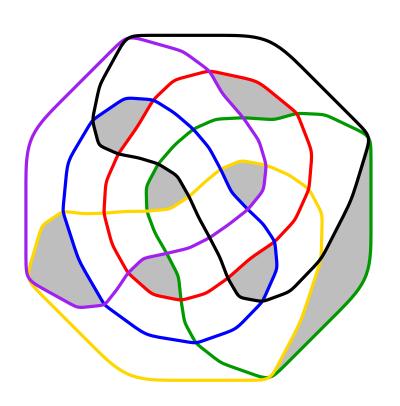


Theorem. There are exactly 3 non-circleable digon-free intersecting n=6 arrangements (2131 classes).



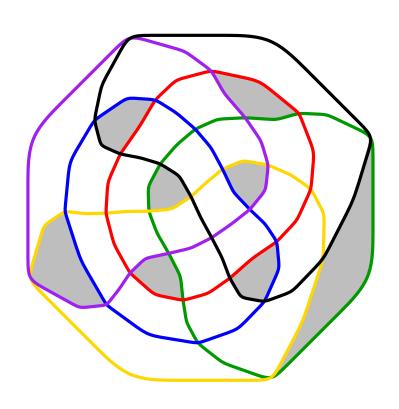
 $\mathcal{N}_6^{\triangle}$ is unique digon-free intersecting with 8 triangular cells

Grünbaum Conjecture: $p_3 \ge 2n-4$



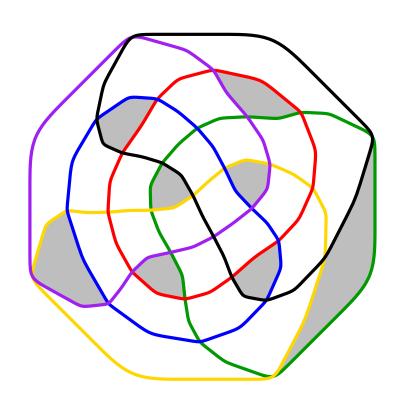
Proof.

based on sweeping argument in 3-D



$$C_1,\ldots,C_6$$
 ...circles (on \mathbb{S}^2)

$$E_1,\ldots,E_6\ldots$$
 planes (in \mathbb{R}^3)



Proof.

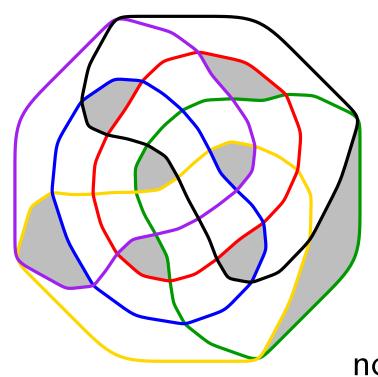
 C_1,\ldots,C_6 ...circles (on \mathbb{S}^2)

 $E_1,\ldots,E_6\ldots$ planes (in \mathbb{R}^3)

move planes away from the origin



 E_i moves to $t \cdot E_i$ as $t \to \infty$



Proof.

 C_1,\ldots,C_6 ...circles (on \mathbb{S}^2)

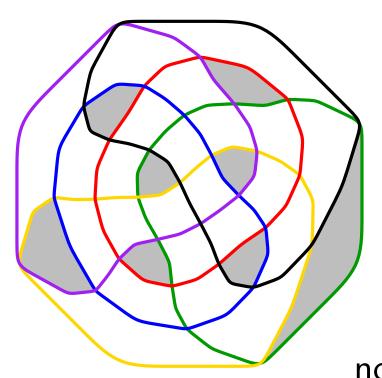
 $E_1,\ldots,E_6\ldots$ planes (in \mathbb{R}^3)

move planes away from the origin

no great-circle arr. \Rightarrow events occur



not all planes contain the origin



Proof.

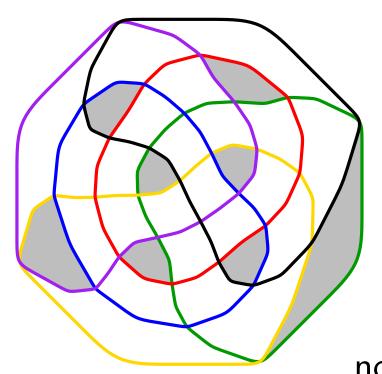
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first event is triangle flip (∄ digons)



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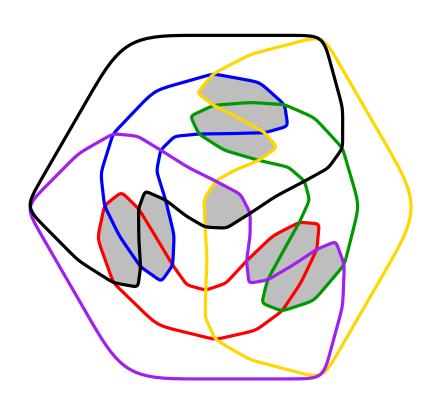
 $E_1,\ldots,E_6\ldots$ planes (in \mathbb{R}^3)

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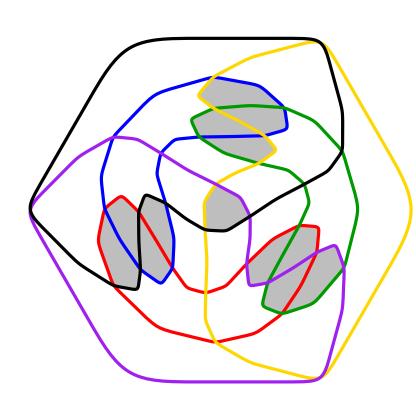
but triangle flip not possible because all triangles in NonKrupp. Contradiction.



Proof. (similar)

 $C_1, \ldots, C_6 \ldots$ circles

 $E_1, \ldots, E_6 \ldots$ planes

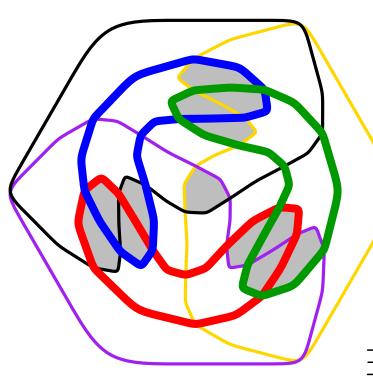


Proof. (similar)

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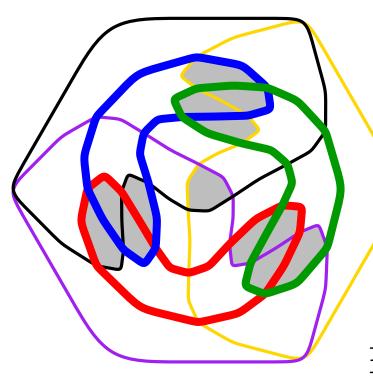
 $E_1, \ldots, E_6 \ldots$ planes

move planes towards the origin

 \exists NonKrupp subarr. \Rightarrow events occur



∃ point of intersection outside the unit-sphere (will move inside)



Proof. (similar)

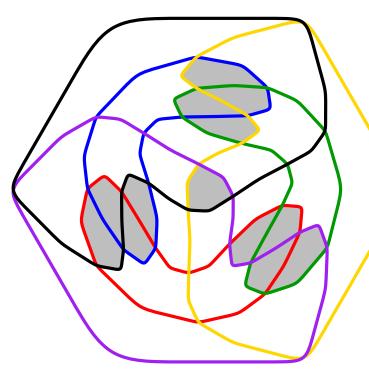
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Proof. (similar)

 $C_1, \ldots, C_6 \ldots$ circles

 $E_1, \ldots, E_6 \ldots$ planes

move planes towards the origin

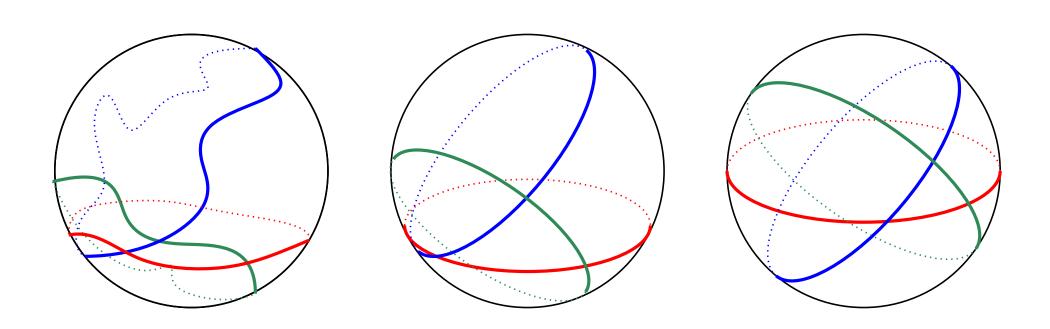
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Great-Circle Theorem:

An arrangement of great-pcs. is circleable (i.e., has a circle representation) if and only if it has a great-circle repr.



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- C_1, \ldots, C_n ... circles E_1, \ldots, E_n ... planes
- move planes towards the origin
- all triples Krupp
 - ⇒ all intersections remain inside
 - \Rightarrow no events

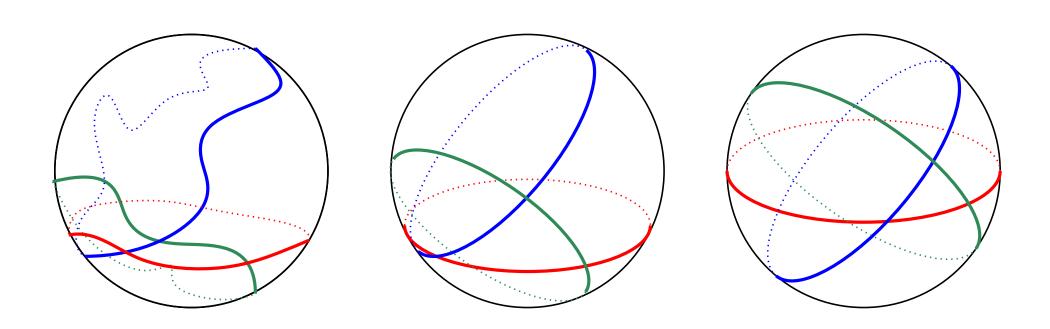
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- move planes towards the origin
- all triples Krupp
 - ⇒ all intersections remain inside
 - \Rightarrow no events
- we obtain a great-circle arrangement

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Corollaries:

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 - ∃ corresponding non-circleable arr. of pseudocircles

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- deciding circleability is $\exists \mathbb{R}$ -complete

```
( NP \subseteq \exists \mathbb{R} \subseteq PSPACE )
```

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An arrangement of great-pcs. is circleable (i.e., has a circle representation) if and only if it has a great-circle repr.

Corollaries:

- ∀ non-stretchable arr. of pseudolines
 ∃ corresponding non-circleable arr. of pseudocircles
- deciding circleability is $\exists \mathbb{R}$ -complete
- ∃ infinite families of minimal non-circ. arrangements
- ∃ arr with a disconnected realization space

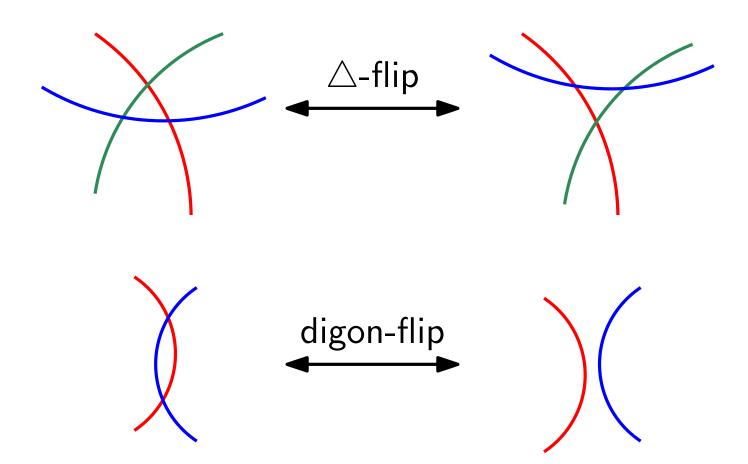
• . . .

Computational Part

- find circle representations heuristically
- hard instances by hand

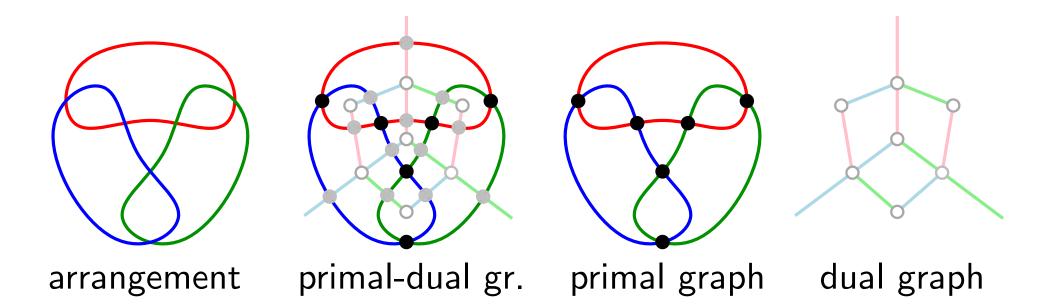
Computational Part

enumeration via recursive search on flip graph



Computational Part

- connected arrangements encoded via primal-dual graph
- intersecting arrangements encoded via dual graph



Part II: Triangles in Arrangements

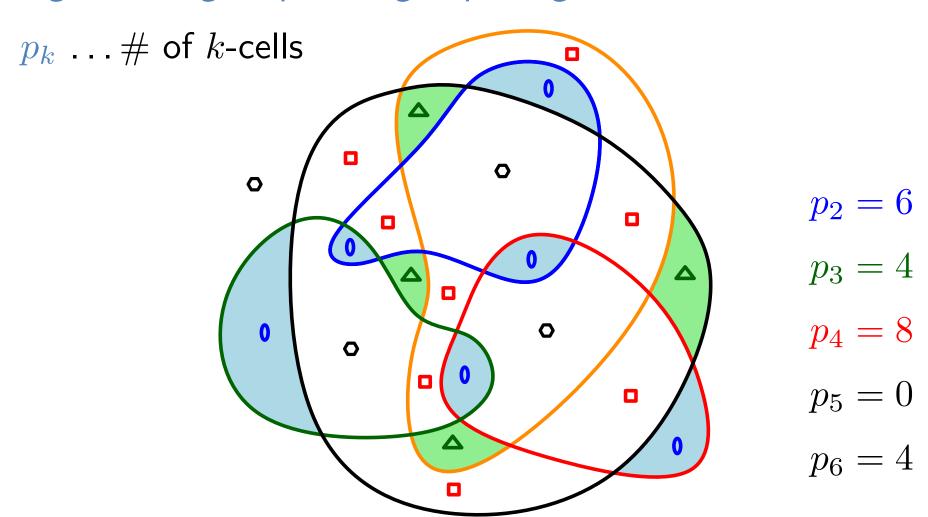
Part II: Triangles in Arrangements

assumption throughout part II:

intersecting ... any 2 pseudocircles cross twice

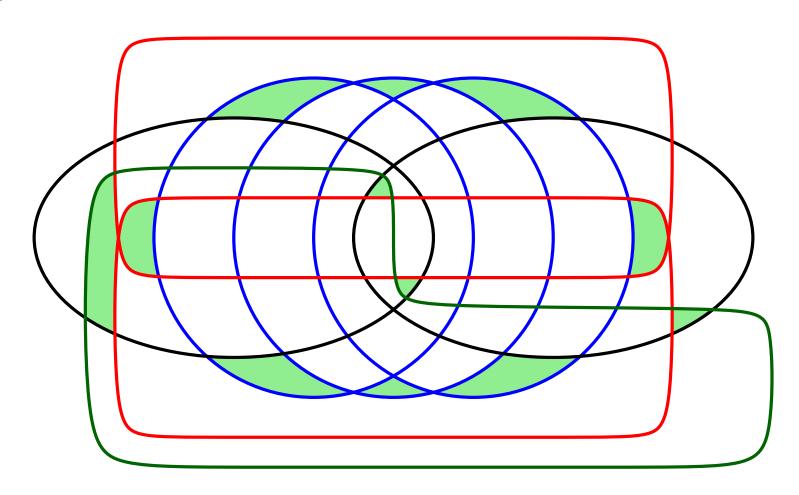
Cells in Arrangements

digon, triangle, quadrangle, pentagon, . . . , k-cell



Grünbaum's Conjecture ('72):

• $p_3 \ge 2n - 4$?



Grünbaum's Conjecture ('72):

• $p_3 \ge 2n - 4$?

Known:

- $p_3 \ge 4n/3$ [Hershberger and Snoeyink '91]
- $p_3 \ge 4n/3$ for non-simple arrangements, tight for infinite family [Felsner and Kriegel '98]

Grünbaum's Conjecture ('72):

• $p_3 \ge 2n - 4$?

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Our Contribution:

- disprove Grünbaum's Conjecture
- $p_3 < 1.\overline{45}n$
- New Conjecture: 4n/3 is tight

Theorem. The minimum number of triangles in digon-free arrangements of n pseudocircles is

- (i) 8 for $3 \le n \le 6$.
- (ii) $\lceil \frac{4}{3}n \rceil$ for $6 \le n \le 14$.
- (iii) $< 1.\overline{45}n$ for all n = 11k + 1 with $k \in \mathbb{N}$.

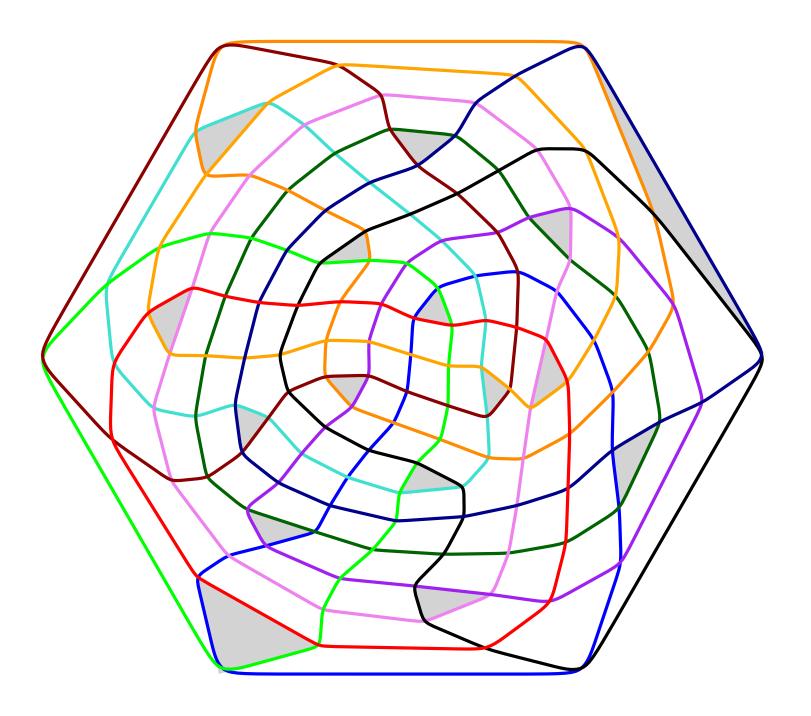


Figure: Arrangement of n=12 pcs with $p_3=16$ triangles.

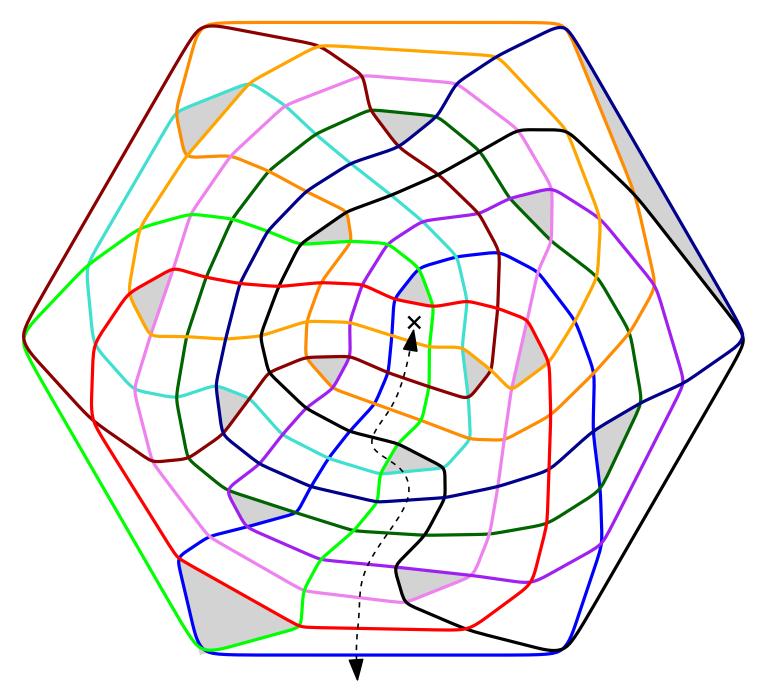
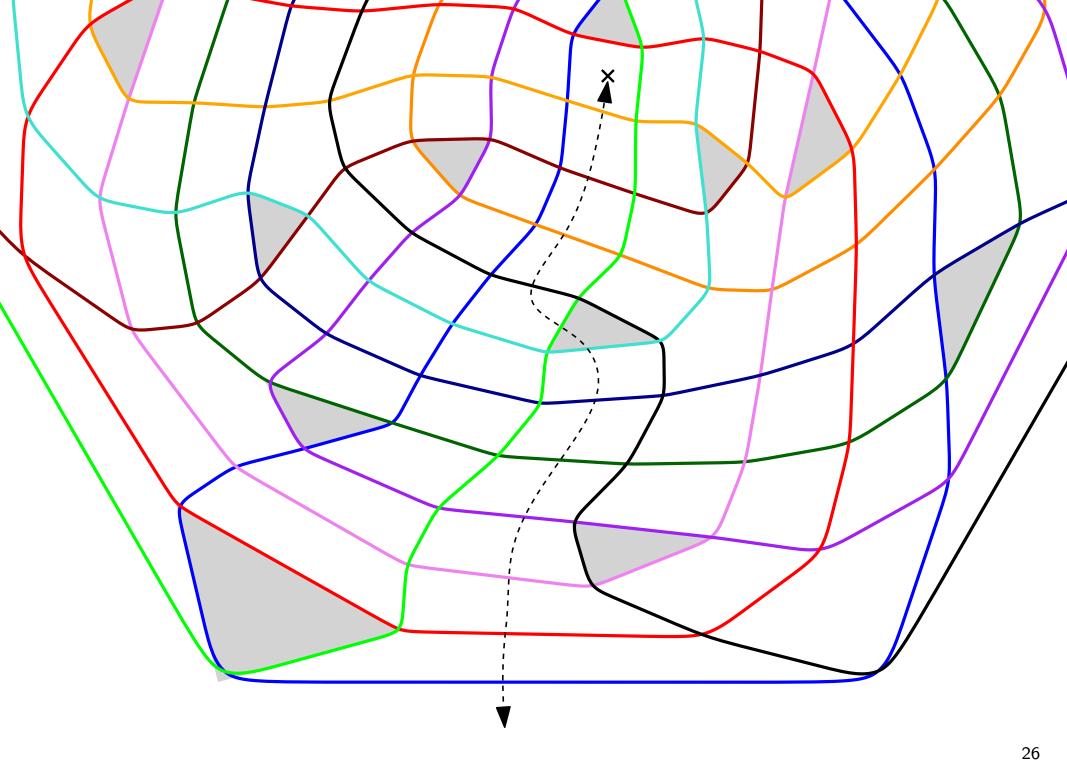
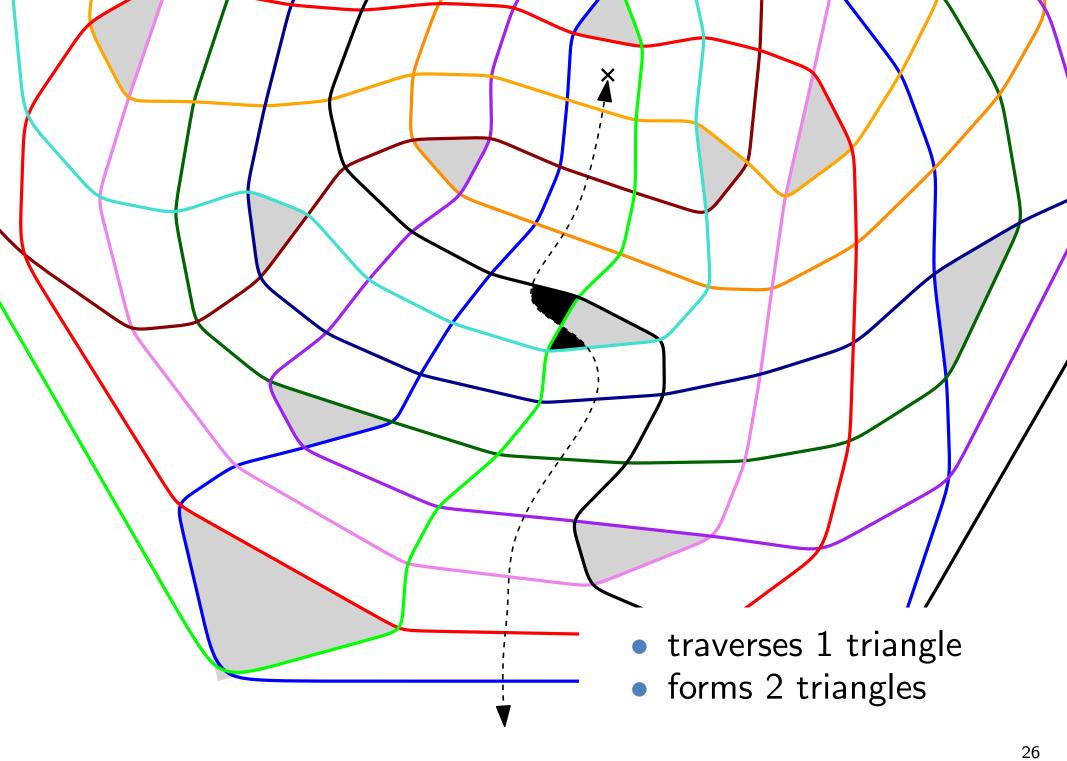
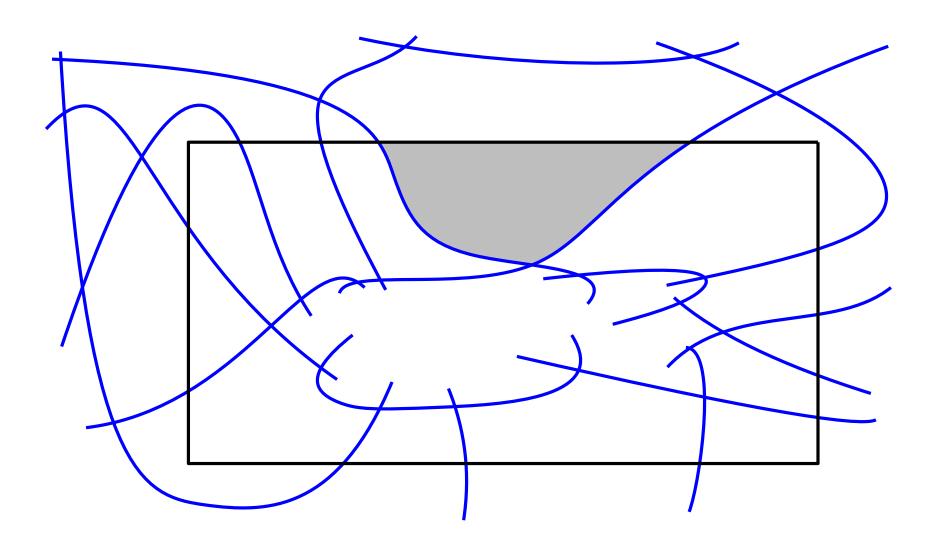
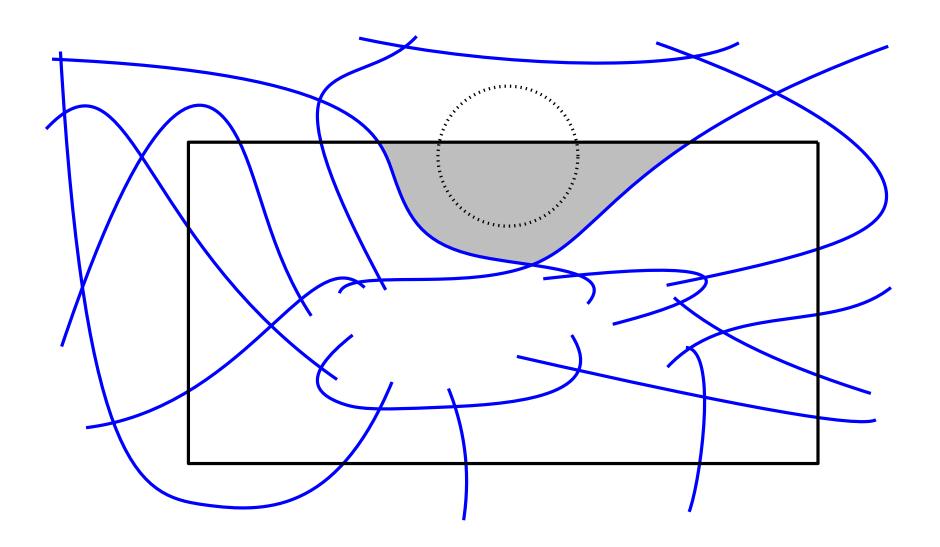


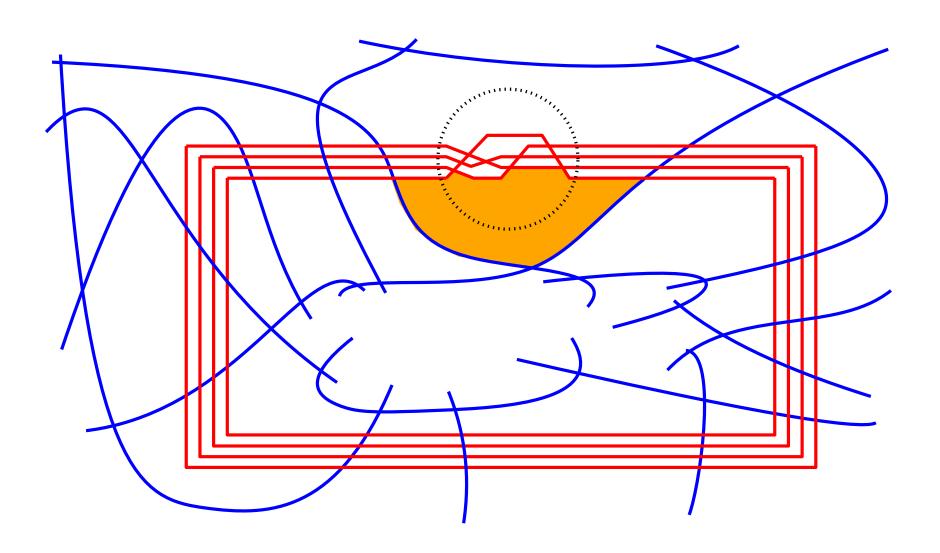
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- start with $\mathscr{C}_1 := \mathscr{A}_{12}$
- ullet merge \mathscr{C}_k and $\mathscr{A}_{12} \longrightarrow \mathscr{C}_{k+1}$
- $n(\mathscr{C}_k) = 11k + 1$, $p_3(\mathscr{C}_k) = 16k$
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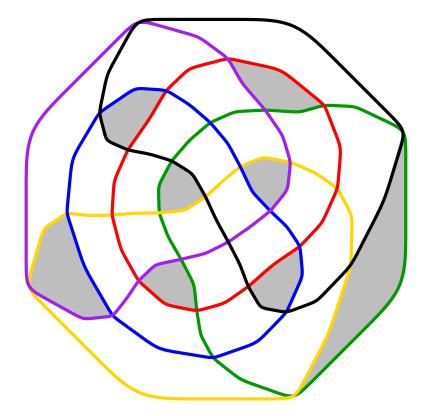
Conjecture. $\lceil 4n/3 \rceil$ is tight for infinitely many n.

• \exists unique arrangement $\mathcal{N}_6^{\triangle}$ with $n=6, p_3=8$

• $\mathcal{N}_6^{\triangle}$ appears as a subarrangement of every arr. with

$$p_3 < 2n - 4$$
 for $n = 7, 8, 9$

• $\mathcal{N}_6^{\triangle}$ is non-circularizable

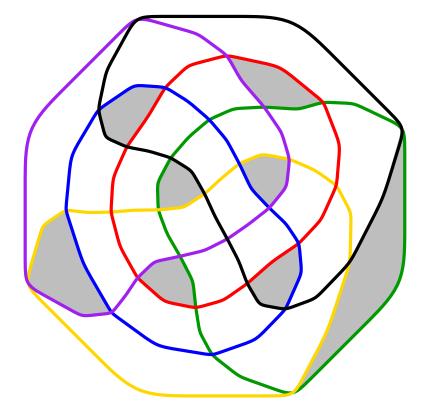


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- ⇒ Grünbaum's Conjecture might still be true for arrangements of circles!



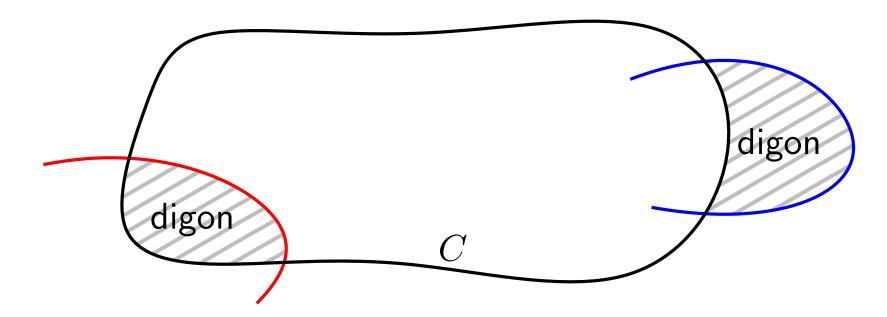
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Proof.

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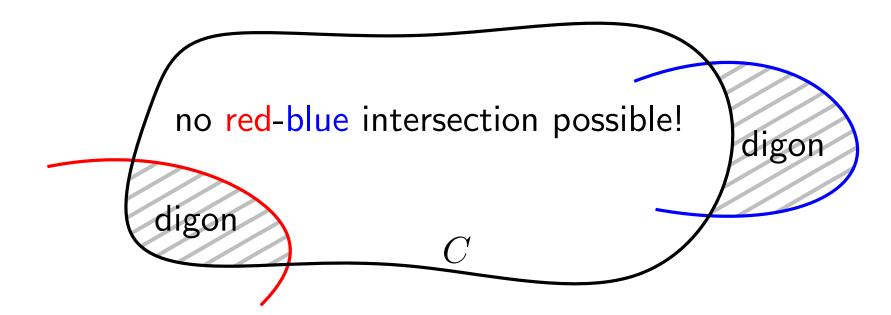
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- All incident digons lie on the same side of C.



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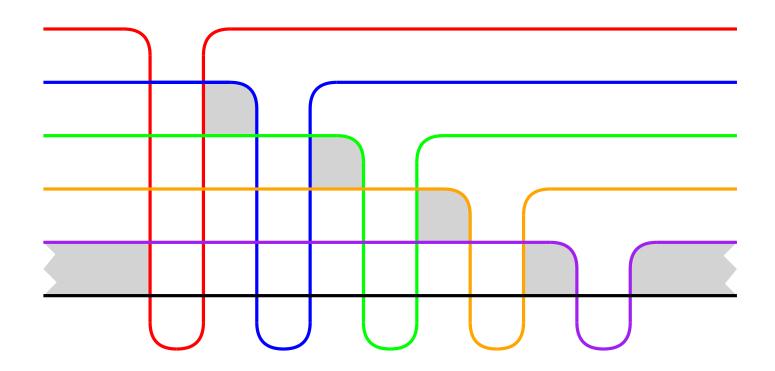
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- C ... pseudocircle in \mathscr{A}
- ullet All incident digons lie on the same side of C.
- ullet \exists two digons or triangles on each side of C [Hershberger and Snoeyink '91] .

Theorem. $p_3 \geq 2n/3$

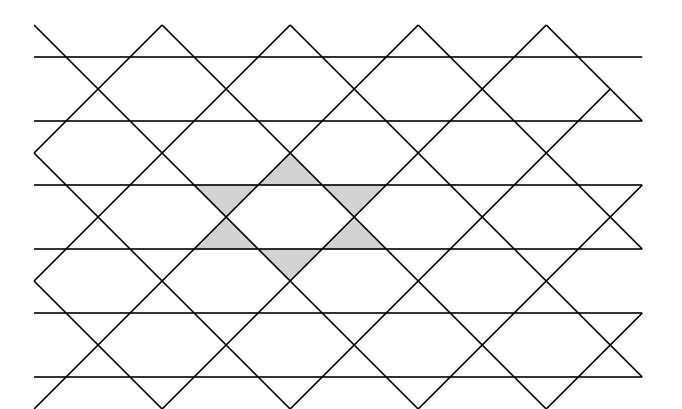
Conjecture. $p_3 \ge n-1$



Theorem.
$$p_3 \leq \frac{2}{3}n^2 + O(n)$$

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n	2	3	4	5	6	7	8	9	10
simple	0	8	8	13	20	29	≥ 37	≥ 48	≥ 60
$+ {\sf digon-free}$	_	8	8	12	20	29	≥ 37	≥ 48	≥ 60
$\lfloor \frac{4}{3} \binom{n}{2} \rfloor$	1	4	8	13	20	28	37	48	60

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 \mathscr{A} ... arrangement of $n \geq 4$ pseudocircles

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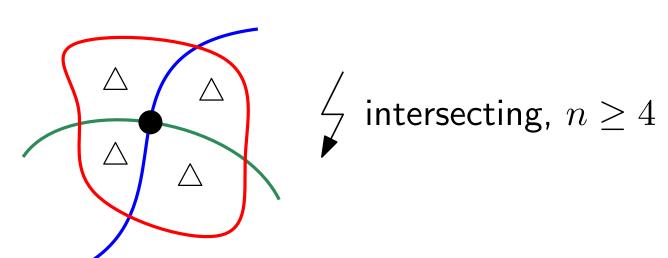
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X'...crossings incident to *precisely* 3 triangles

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- ullet each remaining triangle is incident to $oldsymbol{3}$ crossings of Y



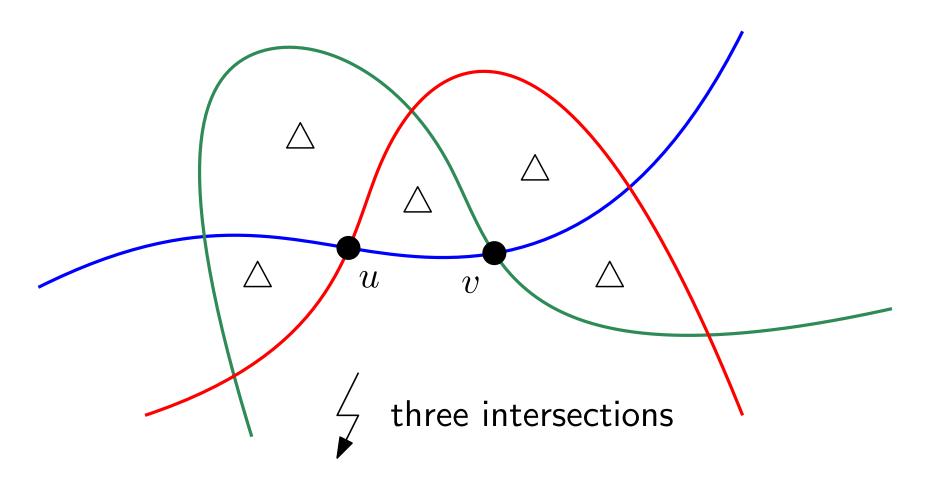
not incident to any vertex from X'

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- ullet each remaining triangle is incident to $oldsymbol{3}$ crossings of Y
- Since $|Y| \le |X| = n(n-1)$, we count

$$p_3 \le \frac{2}{3}|Y| + O(n) \le \frac{2}{3}n^2 + O(n)$$

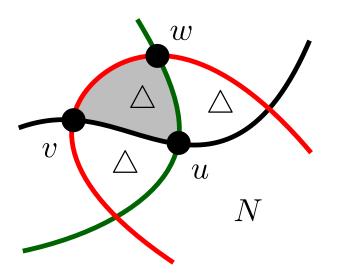
Claim B: Two adjacent crossings u, v in X' share two triangles.



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Claim C: Let u, v, w be three distinct crossings in X'. If u is adjacent to both v and w, then v is adjacent to w.

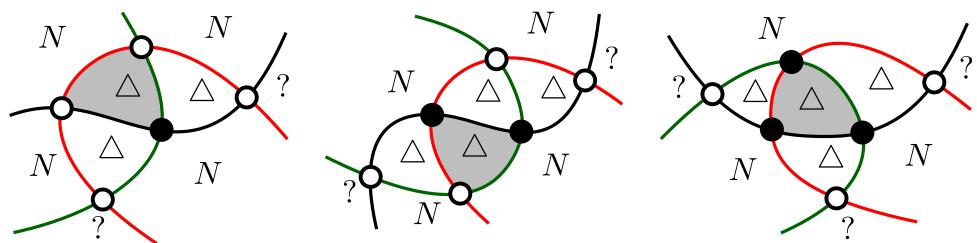
(If both edges uv and uw are incident to two triangles, then uvw form a triangle)



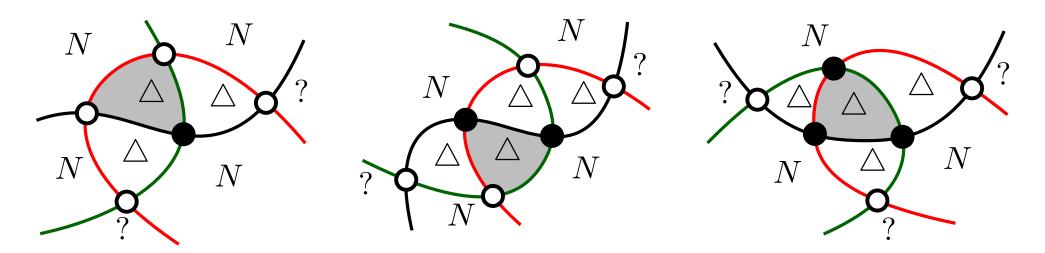
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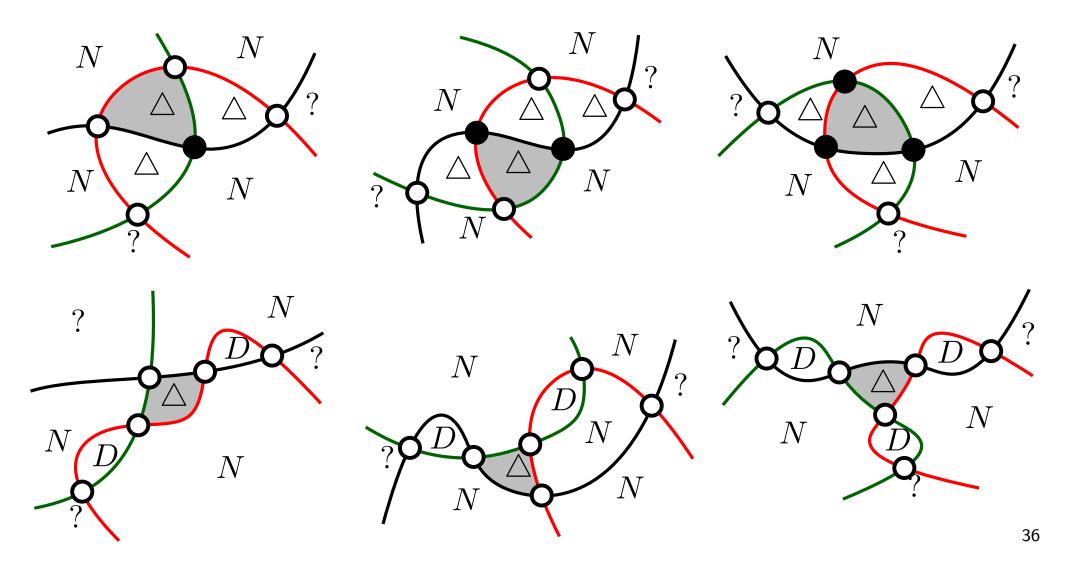
 \Rightarrow each connected comp. of the graph induced by X' is either singleton, edge, or triangle.



We can convert crossings of X' into digons using \triangle -flips!



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There are at most O(n) digons [Agarwal, Nevo, Pach, Pinchasi, Sharir, Smorodinsky 2004]

- \Rightarrow at most O(n) flips
- $\Rightarrow |X'|$ at most O(n)

$$\Rightarrow p_3 \le \frac{2}{3}n^2 + O(n)$$

