## Contact Representations of Planar Graphs: Combinatorial Structure and Algorithm $\mathcal{X}$

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## Outline

## Segment Contacts - Quadrangulations <br> Algorithm $\mathcal{X}$ does its job

Algorithm $\mathcal{X}$
The abstract view
Homothetic Triangle Contact Representations
Monster packing and Algorithm $\mathcal{X}$
Square Contact Representations Including an existence proof

More Contact Representations - Pentagons
5-color trees and the coin theorem

## Segment Contact Representations



A quadrangulation with a segment contact representation.

## Induced separating Decompositions



The contact representation induces a separating decomposition on the quadrangulation.

## Separating Decompositions

$G=(V, E)$ a plane quadrangulation with outer face
$F=\{s, x, t, y\}$.
A coloring and orientation of the edges of $G$ with colors red and blue is a separating decomposition of $G$ iff


## Separating Decompositions and 2-Orientations

Theorem.
Separating decompositions and 2-orientations are in bijection.
Proof.

- Define the path of an edge:

- The path is simple (Euler), hence, ends in a sink: in the red $s$ or in the blue $t$.


## Sketch: Compute Segment Contact Representations

- Compute a separating decomposition.
- Separate the two trees - 2-page book embedding.



## Alternating and Full Binary Trees

Proposition. There is bijection between on-sided and binary trees that preserves types (left/right) of nodes.


## Sketch: Compute Segment Contact Representations

- The two binary trees obtained from the separating decomposition fit together.



## More Problems for Segment Contact Representations

Add some conditions.

- Find a representation in a square such that all intersect the diagonal.
(We just saw a solution)
- Find a representation such that all inner rectangles are squares.

The dissection of rectangles into squares
Brooks, Smith, Stone and Tutte 1940.

- Find a representation such that each inner rectangle has its own prescribed aspect ratio.
(aspect ratio universality)


## Segment Contact Squarings



## Computing a Segment Contact Squaring



Step I: Compute a separating decomposition on $Q$.
This describes the contacts of the segments.
(combinatorial or abstract contact representation)

## Equations



$$
\begin{aligned}
& x_{1}+x_{2}=1 \\
& \quad x_{1}=x_{2}+x_{3} \\
& x_{1}+x_{3}+x_{5}=x_{7}+x_{8} \\
& \quad x_{2}=x_{3}+x_{4} \\
& \quad \vdots
\end{aligned}
$$

Step II: Set up a linear system of equations: $A_{S} \cdot x=e_{1}$

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$$

Step II: Set up a linear system of equations: $A_{S} \cdot x=e_{1}$

- $A_{S}$ is a square matrix.

Theorem. $\operatorname{det}\left(A_{S}\right) \neq 0$.

## The Determinant



- $A_{S}$ is a bipartite adjacency matrix of a bipartite graph $H$.
- $\operatorname{det}\left(A_{S}\right)=\sum_{\pi} \operatorname{sign}(\pi) \prod a_{i, \pi(i)}$
- $\operatorname{det}\left(A_{S}\right)=\sum_{M} \operatorname{sign}(M)$, with $M$ perfect matching of $H$.
- $\operatorname{sign}(M)=\operatorname{sign}\left(M^{\prime}\right)$ for all $M$ and $M^{\prime}$ perfect matching. Face cycles: length 4, odd number of minus signs.
- $H$ has a perfect matching.


## Solve the System



## Flip the Negative



- The boundary of the negative contains no complete segment.
- Hence, the boundary is a directed cycle in the separating decomposition
- Flip boundaries of negative areas to get the good separating decomposition and the squaring.


## The Squaring



- The method also works for given aspect ratios.
- Alternative algorithms exist.

Example: electric flow in bipolar networks.

## Algorithm $\mathcal{X}$

## Algorithm $\mathcal{X}$ a Coarse Description

Aim for a contact representation of $G$.

- Compute a combinatorial representation - directed graph $D$.
- Extract linear equations: $A_{D}=e_{1}$.
- Show that $A_{D}$ is square and $\operatorname{det} A_{D} \neq 0$.
- Solve the system.

Solution positive - Done.

- Show that negative variables in the solution correspond to directed cycles in $D$ which can be flipped ( $D \rightarrow D^{\prime}$ ).
- Try again with $D^{\prime}$.


## Homothetic Triangle Contact Representations



## Triangle Contact Representation

Theorem [ de Fraysseix, Ossona de Mendez and Rosenstiehl ]. Triangulations have triangle contact representations.


## Homothetic Triangle Contact Representations

Theorem [ Gonçalves, Lévêque, Pinlou 2010].
Every 4-connected triangulation has a triangle contact representation with homothetic triangles.


## Triangle Contact Representations

Proof uses Schramm's "Monster Packing Theorem".
Theorem. Let $T$ be a planar triangulation with outer face $\{a, b, c\}$ and let $C$ be a simple closed curve partitioned into arcs $\left\{C_{a}, C_{b}, C_{c}\right\}$. For each interior vertex $v$ of $T$ prescribe a convex set $P_{v}$ containing more than one point. Then there is a contact representation of $T$ with homothetic copies.

## Remark.

In general homothetic copies of the $P_{v}$ can degenerate to a point.
This is impossible if $T$ is 4-connected and all the $P_{v}$ are homothetic triangles.

## Schnyder Woods

$G=(V, E)$ a plane triangulation, $F=\left\{a_{1}, a_{2}, a_{3}\right\}$ the outer triangle.
A coloring and orientation of the interior edges of $G$ with colors $1,2,3$ is a Schnyder wood of $G$ iff

- Inner vertex condition:

- Edges $\left\{v, a_{i}\right\}$ are oriented $v \rightarrow a_{i}$ in color $i$.


## Schnyder Woods and 3-Orientations

Theorem. Schnyder wood and 3-orientation are in bijection.

## Proof.

- All edges incident to $a_{i}$ are oriented $\rightarrow a_{i}$.
$G$ has $3 n-9$ interior edges and $n-3$ interior vertices.
- Define the path of an edge:

- The path is simple (Euler), hence, ends at some $a_{i}$.
- Two path starting at a vertex do not meet again (Euler).


## Schnyder Woods as Abstract Triangle Contact

 Representations

## Triangle Contacts and Equations



The abstract triangle contact representation implies equations for the sidelength:
$x_{i}+x_{j}+x_{k}+x_{\ell}=1$ and $x_{a}+x_{b}+x_{c}=x_{v}$ and $x_{d}=x_{v}$ and $x_{e}=x_{v}$ and $x_{d}+x_{e}=x_{w}$ and $\ldots$

## The System of Equations

- The matrix $A_{S}$ is square.
- $A_{S}$ corresponds to a bipartite graph


Every face has length 6 and two negative signs.

- The bipartite graph has a perfect matching.

Theorem. $\operatorname{det} A_{S} \neq 0$.

- The system of equations has a unique solution.


## Negatve Variables

In the solution some variables may be negative.

- The boundary of the negative variables induces a directed cycle in the Schnyder wood.



## Flipping Cycles

From the bijection
Schnyder woods $\Longleftrightarrow$ 3-orientations it follows that cycles can be reverted (flipped).


## The Status of Algorithm $\mathcal{X}$ in this Case

- Homothetic triangle contact representations are an instance for Algorithm $\mathcal{X}$.
- We have not been able to prove that the algorithm stops. In practice, however, it does!


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Program written by Julia Rucker.

## The Status of Algorithm $\mathcal{X}$ in this Case

Theoretical support:
Theorem. A negative triangle becomes positive by flipping.


## Squarings of Inner Triangulations.



## Squarings for Inner Triangulations

Theorem [ O. Schramm 1993]. The squaring of a 5-connected inner triangulation exists and it is is unique.

O. Schramm Square Tilings with prescribed Combinatorics. 1993 Schramm uses extremal lengths,
Lovász gave a proof using convex corners.

## Transversal structures

Proposition. 4-connected inner triangulations of a quadrangle admit transversal structures (a.k.a. regular edge labeling).


## A Systems of Equations

 transversal structure $\Longrightarrow$ rectangular dissection (abstract squaring).

- red rectangles and red circles $=$ variables
- white circles $=$ equations


## Properties of the System

- The matrix $A_{T}$ is a square matrix.
- The bipartite graph has facial cycles of length 10 with four negative signs $\Longrightarrow$ perfect matchings have the same sign.
- The graph has a perfect matching



## Flips on Transversal Structures

- It is possible to associate a digraph $D$ to a inner triangulation with a transversal structure such that negative variables induce a directed boundary cycle which can be reverted.
- We describe the elementary flips directly:



## A Program



## Written by Thomas Picchetti

## Existence Reproved (Hendrik Schrezenmaier)

## Theorem.

Every inner triangulation $G$ of a 4-gon admits a squaring.
Proof (Sketch)

- Let $R$ be a rectangulation of $G$ with aspect ratio vector $\alpha_{0}$ and transversal structure $T_{0}$.
- Let $\alpha_{1}=\mathbf{1}$ be the aspect ratio vector of a squaring and let $\ell=\left\{\alpha_{t}: t \in[0,1]\right\}$ be the line from $\alpha_{0}$ to $\alpha_{1}$.
- The set $A_{0}$ of all aspect ratio vectors $\beta$ representable by $T_{0}$ is a subset of $\mathbf{R}^{n}$ containing $\alpha_{0}$. The set is defined by polynomial inequalities (positivity) with polynomials of bounded degree (determinant, Cramer's rule).
- When $\ell$ leaves $A_{0}$ some variables change their sign in $T_{0}$ this set corresponds to a flippable set, this defines $T_{1}$ and $A_{1}$.
- Continue until $\alpha_{1}$ is reached.

Pentagon contact representations


## Pentagon Contact Representations


$G$ an inner triangulation of the 5 -gon $a_{1}, \ldots, a_{5}$

- Existence
- Uniqueness
- Combinatorial structure
- Computation


## Homothetic Pentagon Contact Representations

Theorem.
Every triangulation of a 5-gon has a contact representation with homothetic pentagons.

Proof: Use Schramm's "Monster Packing Theorem".


With pentagons there are no degeneracies.

## The combinatorial structure: five color forests

Definition (Five color forest) Orientation and coloring of inner edges of inner triangulation of 5-gon $a_{1}, \ldots, a_{5}$, s.t.


- no incoming edge of color $i$
$\Rightarrow$ outgoing edge of color $i-2$ or $i+2$ exists

Theorem
Regular pentagon contact representation induces five color forest on its contact graph.


Five color forests $\leftrightarrow(5,2)$-orientations


- outdeg $(\bullet)=5$
- outdeg(○) $=2$

Five color forests $\leftrightarrow(5,2)$-orientations


Theorem
There is a bijection between the five color forests and
$(5,2)$-orientations of a graph $G$.

## Abstract contact representations

- Compute a five color forest.
- This yields an abstract contact representation.



## A system of linear equations



Variables:

- one side length for each vertex: $x_{v}$
- four side lengths for each face: $x_{f}^{(1)}, \ldots, x_{f}^{(4)}$

Equations:

- five for each vertex: $x_{v}=$ sum of touching face side lengths
- two for each face: $x_{f}^{(3)}=x_{f}^{(1)}+\phi x_{f}^{(2)}, \quad x_{f}^{(4)}=\phi x_{f}^{(1)}+x_{f}^{(2)}$
- one inhomogeneous: length of upper segment $=1$


## System of linear equations

Computing a regular pentagon contact representation induced by a fixed five color forest

Lemma
The system $A_{F} X=\mathbf{e}_{1}$ is uniquely solvable.
Lemma
$x \geq 0 \Leftrightarrow$ there is a regular pentagon contact repr. inducing $F$

## Algorithm $\mathcal{X}$ for Pentagons

- Guess a five color forest $F$
- Case 1: solution of $A_{F} X=\mathbf{e}_{1}$ is nonnegative
- construct contact repr. from solution
- Case 2: solution contains negative and nonnegative variables
- Lemma: neg. and nonneg. variables are separated by oriented cycles in the $(5,2)$-orientation

- change orientation of these cycles
- restart with new (5, 2)-orientation, resp. five color forest


## The Status

It works!


But, we have no proof that the process always ends.

## 5-gons



13-gons


27-gons


Before you ask:
Yes circle contact representations are accumulation points.

## The Efficiency of Algorithm $\mathcal{X}$



- For each pair $(n, K)$ the dot is the average over 100 random graphs.


## The End

## Visit and enjoy:

www3.math.tu-berlin.de/diskremath/research/kgon-representations/index.html


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