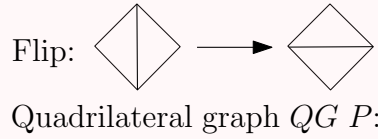


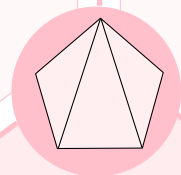
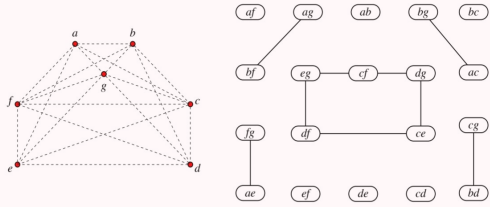
Introduction



Flip graph G

Flip distance: d_G

- Vertices: Diagonals of P
- Edge $ab-cd$:
 ab, cd cross and $\{a,b,c,d\}$ form an empty quadrilateral



Pentagons and partial cubes

Let P be a finite point set. Equivalent:

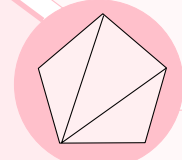
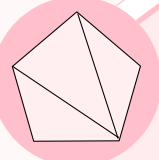
- P has no empty pentagon
- $QG P$ is a forest
- The flipgraph of P is a partial cube

Finding pentagons and flips

- One may test whether any set of n points contains an empty pentagon, in time $\mathcal{O}(n^2)$.
- If a set P of n points has no empty pentagon, one can construct $QG P$ in time $\mathcal{O}(n^2)$.

Main Result

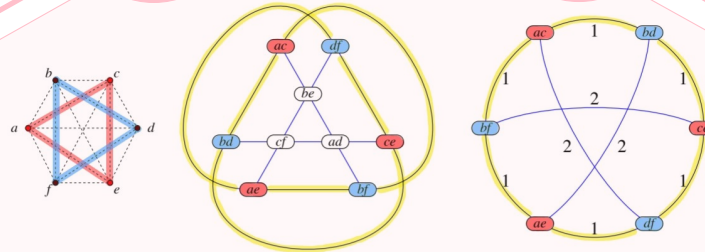
The triangulations of a finite point set from a flip graph that can be embedded isometrically into a hypercube, if and only if the point set has no empty convex pentagon. As a consequence, flip distance in such point sets can be computed efficiently.



Computing and estimating flip distance

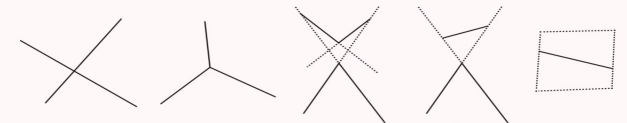
- Let T, T' be any triangulations of P
- Form a complete bipartite graph with the diagonals of T on one side and the diagonals of T' on the other
- Label the edge between any two diagonals by the distance between those diagonals in $QG P$
- Let M be the minimum weight of a perfect matching in this complete bipartite graph

M can be computed in polynomial time, and provides an underestimate of the true flip distance between T and T'

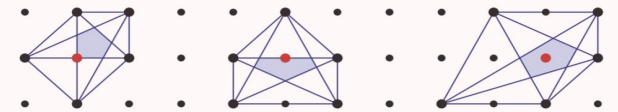


A hexagon (left), its $QG P$ (center) and the complete bipartite graph of distances in $QG P$ from $\{ac, ae, ce\}$ to $\{bd, bf, df\}$ (right)

Examples



Continuous point sets with no pentagons



Any lattice pentagon has another lattice point in it



Point sets with no empty pentagon derived from nested regular polygons

Notice: For instance, if P has no empty pentagons, this estimate ends up calculating the flip distance exactly.