

Due for the exercise session: February 9., 2021

- (1) Try to find good time bounds for computing
 - a) The Ferrer's shape $\text{Fer}(P)$ of a poset P .
 - b) A maximum k -chain in P .
 - c) A maximum ℓ antichain in P .

- (2) Think about analogs of the Greene-Kleitman theory for directed acyclic graphs.
- (3) Given a 2-dimensional poset P , find good time bounds for computing
 - a) The skeleton of P
 - b) The Ferrer's shape $\text{Fer}(P)$ of P .
 - c) A maximum k -chain in P .

- (4) This refers to lecture 28, so P is 2-dimensional. Show that the k chain \mathcal{C} of X obtained from a maximum $k - 1$ chain \mathcal{C}' of $S(X)$ is maximum without referring to ℓ antichains and orthogonal pairs.

Hint: Define a mapping from k chains in X to $k - 1$ chains in $S(X)$.
- (5) Visit the pages of Xavier Viennot (www.viennot.org and www.xavierviennot.org), stray around on these pages and tell us about findings you enjoyed.
- (6) Find out about a result obtained by Logan–Shepp and Vershik–Kerov. What does it say about posets?
- (7) Find an up-growing on-line width 2 poset P such that the First-Fit algorithm (greedy) uses an unbounded number of chains.
- (8) Within Lecture 29 we have seen that Algorithm has a strategy in on-line partition game on up-growing orders of width at most w to use at most $\binom{w+1}{2}$ chains. Show that this is best possible (so devise a strategy for Presenter).
- (9) Find the value of the on-line antichain partition game on interval orders of height at most h presented with representation.
- (10) Find the value of the on-line chain partition game on semi-orders of width at most w (presented without representation)