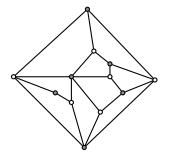
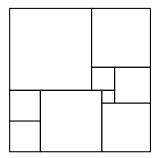
Add some conditions.

- Find a representing rectangulation R in a square such that all rectangles intersect the diagonal. (We saw a solution)
- Find a representing rectangulation *R* such that all inner rectangles are squares.

*The dissection of rectangles into squares* Brooks, Smith, Stone and Tutte 1940.

# Segment Contact Squarings

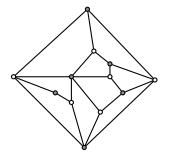


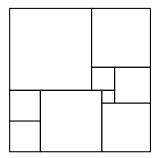


View the bipolar graph as electrical network with edges of resistance 1 Ohm. Consider an  $s \rightarrow t$  flow in this network. The distribution of current in edges corresponds to a squaring. Based on this theory they give explicit solutions:

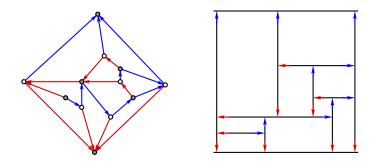
size(i, j) =# spanning trees T with (i, j) on the  $s \to t$  path in T- # spanning trees T with (j, i) on the  $s \to t$  path in T.

# Segment Contact Squarings



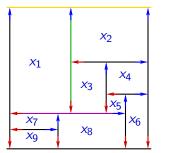


Computing a Segment Contact Squaring



Step I: Compute a separating decomposition on Q. This yields a rectangular dissection.

## Computing a Segment Contact Squaring



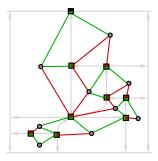
• 
$$x_1 = x_2 + x_3$$
  
•  $x_1 + x_3 + x_5 = x_7 + x_8$   
•  $x_1 + x_2 = 1$ 

Step II: Set up a linear system of equations:  $A_S \cdot x = e_1$ 

• A<sub>5</sub> is a square matrix.

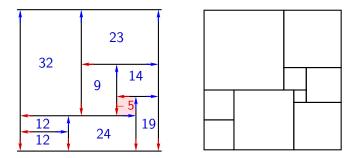
**Theorem.** det( $A_S$ )  $\neq$  0.

## Computing a Segment Contact Squaring



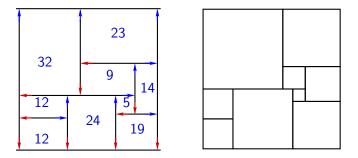
- $A_S$  is a bipartite adjacency matrix of a bipartite graph H.
- $det(A_S) = \sum_{\pi} sign(\pi) \prod a_{i,\pi(i)}$
- $det(A_S) = \sum_M sign(M)$ , with M perfect matching of H.
- sign(M) = sign(M') for all M and M' perfect matching.
- *H* has a perfect matching.

#### Squarings with Segment Contacts



Step III: Flip negative faces to get the good separating decomposition and the squaring.

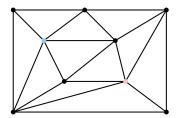
#### Squarings with Segment Contacts

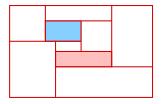


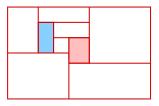
Step III: Flip negative faces to get the good separating decomposition and the squaring.

# Rectangular and square duals

#### Rectangular Duals





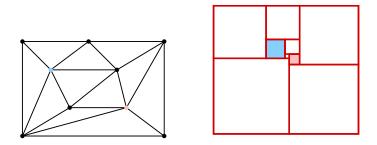


Rectangular duals are not unique.

They exist if the triangulation is 4-connected.

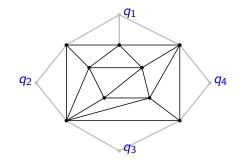
## Squarings for Inner Triangulations

The squaring is unique.



O. Schramm Square Tilings with prescribed Combinatorics. 1993

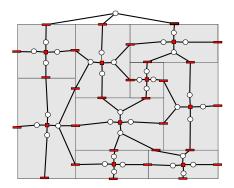
#### A Polyhedral view on Squarings



 $P = \{x \in \mathbb{R}^{V}_{\geq 0} : \sum_{i \in \gamma} x_{i} \geq 1 \text{ for all } q_{1} \to q_{3} \text{ paths } \gamma \}$ solution vector for min( $\sum_{i} x_{i}^{2} : x \in P$ ) yields squaring. L. Lovász Geometric Representations of Graphs, 2009, Sec. 6.3.2.

## Computing Squarings with Systems of Equations

transversal structure  $\implies$  rectangular dissection.

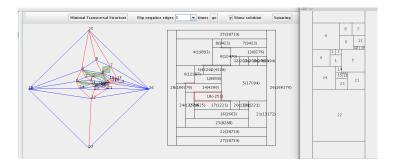


- red rectangles = variables
- white circles = equations

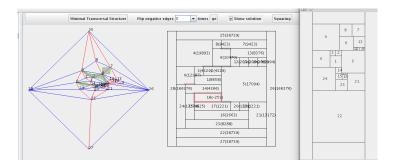
#### The algorithm

- 1. T transversal structure on a 4-triangulation G.
- 2. Calculate the solution  $x_T$  of the equation system Ax = b.
- 3. If  $x_T \ge 0$ 
  - Construct a square contact representation from  $x_T$ .
- 4. Else
  - T' obtained by 'reverting bad-directed cycle' in T.
  - Set T := T' and go to 2.

#### The program



#### The program



In the lecture we saw a demo.