(1) Let T(n) be the number of spanning trees of the graph, which has 2n vertices and is shown below.



Prove  $T(n) = 4 \cdot T(n-1) - T(n-2)$  for all  $n \ge 2$ . How fast does T(n) grow for  $n \to \infty$ ?

- (2) Let G = (V, E) be a simple graph. The *line graph*  $\mathcal{L}(G)$  has vertex set E and two vertices  $e, e' \in E$  of  $\mathcal{L}(G)$  are connected if and only if they share an endpoint, i.e. if  $e \cap e' \neq \emptyset$ . A *Hamilton cycle* of a graph G is a cycle, visiting every vertex exactly once. Specify an algorithm which decides if  $\mathcal{L}(G)$  has a Hamilton cycle in  $\mathcal{O}(|V| + |E|)$  steps, i.e. the algorithm does only a number of steps which is (at most) linear in |V| + |E|.
- (3)
- (a) How many spanning trees does  $K_{2,l}$  have? How many classes of isomorphic spanning trees does it have?
- (b) How many spanning trees and classes of isomorphic spanning trees does  $K_{3,l}$  have?
- (c) Let  $m \leq n$ . How many spanning trees does  $K_{m,n}$  have?
- (4) Let G be a graph and v a vector in the cycle space of G. Let M be the set of edges of G, whose corresponding entries in v are equal to 1. Characterize M.
- (5) Please hand in your solution of this exercise: A matroid M = (E, U) consist of a set E and a family  $U \subseteq Pot(E), U \neq \emptyset$  of subsets of E, which fullfill the following properties:
  - $\emptyset \in \mathcal{U}$ ,
  - $A \subseteq B$  and  $B \in \mathcal{U}$  imply  $A \in \mathcal{U}$ , and
  - for all  $A, B \in \mathcal{U}$  with |A| < |B| exists  $x \in B$ , such that  $A \cup \{x\} \in \mathcal{U}$  holds.

Let G = (V, E) be a graph and  $U \subseteq Pot(E)$  be the set of cycle free sets of edges of G. Prove that (E, U) is a matroid.