

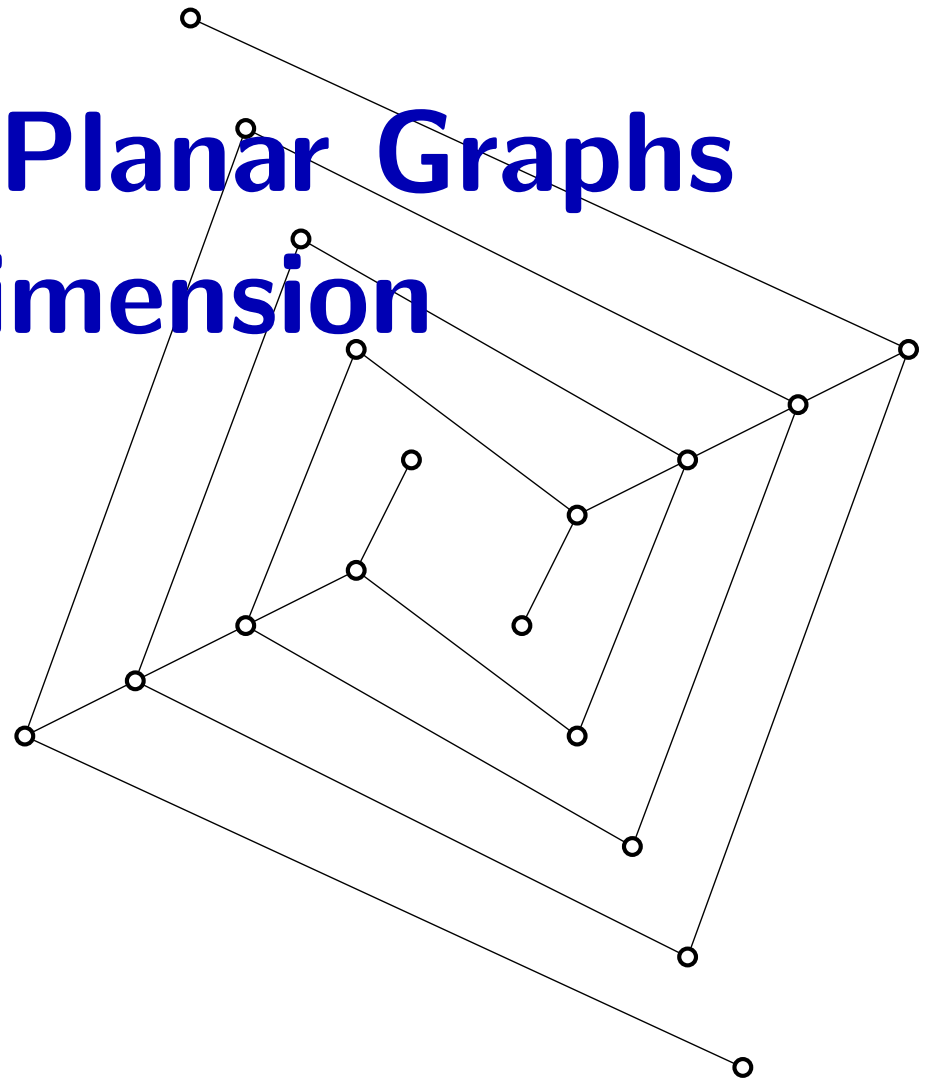
Embeddings of Planar Graphs

Lecture 13 – Dimension

March 2006

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Modules

Dimension

Orthogonal Surfaces and Dimension

Modules

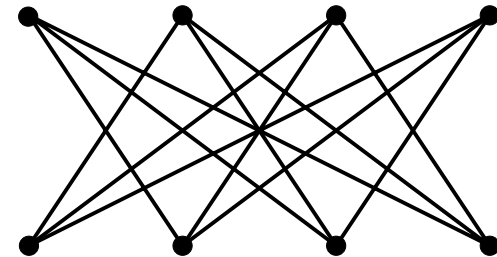
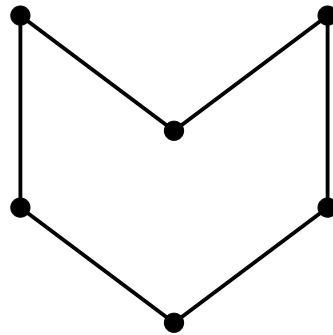
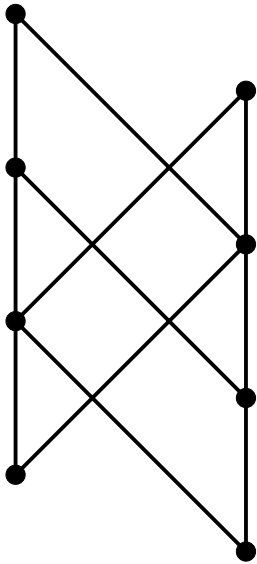
Dimension

Orthogonal Surfaces and Dimension

Finite Orders

$P = (X, <)$ is an **order** iff

- X finite set
- $<$ transitive and irreflexive relation on X .

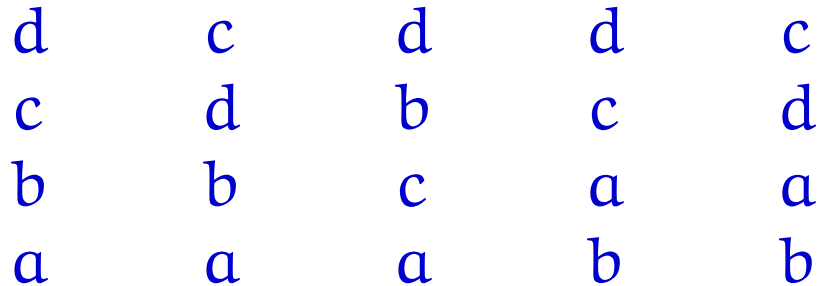
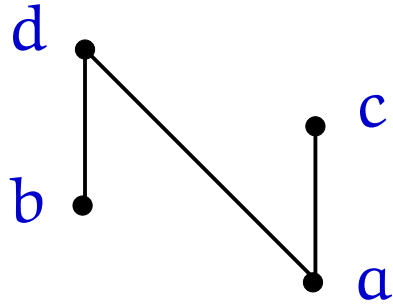


Orders drawn as diagrams.

Linear Extensions

A **linear extension** of $P = (X, <_P)$ is a linear order L , such that

- $x <_P y \implies x <_L y$ ■



Specifying Linear Extensions: Two Recipes

- The generic algorithm
- Acyclic extensions

Dimension of Orders

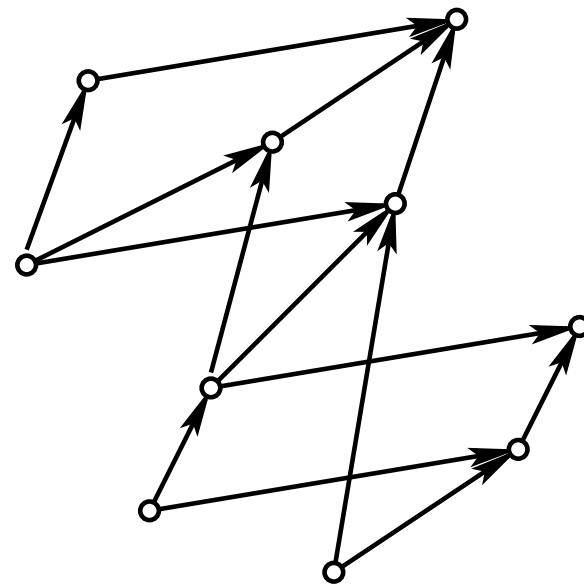
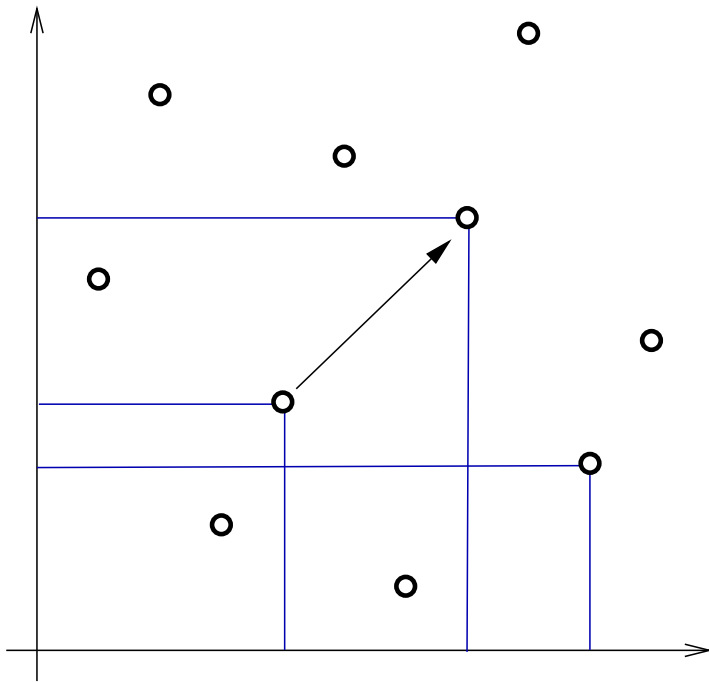
A family \mathcal{L} of linear extensions is a **realizer** for $P = (X, <)$ provided that

- * for every incomparable pair (x, y) there is an $L \in \mathcal{L}$ such that $x < y$ in L .

The **dimension**, $\dim(P)$, of P is the minimum t , such that there is a realizer $\mathcal{L} = \{L_1, L_2, \dots, L_t\}$ for P of size t .

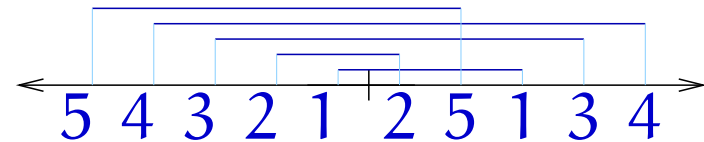
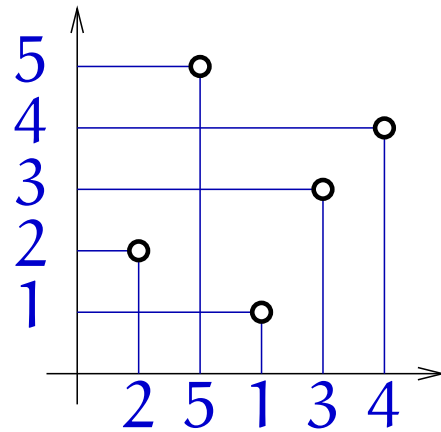
Dimension of Orders

The **dimension** of an order $\mathcal{P} = (X, <)$ is the least t , such that \mathcal{P} is isomorphic to a suborder of \mathbb{R}^t with the product ordering.

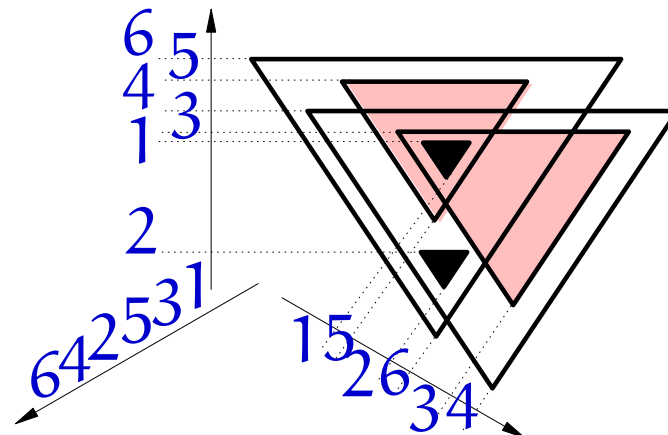
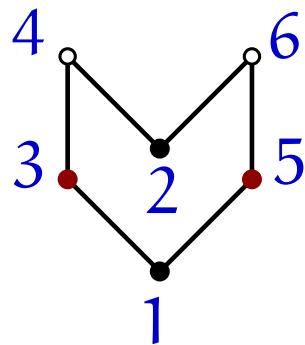


Characterizations

- Dimension 2: Containment orders of intervals.

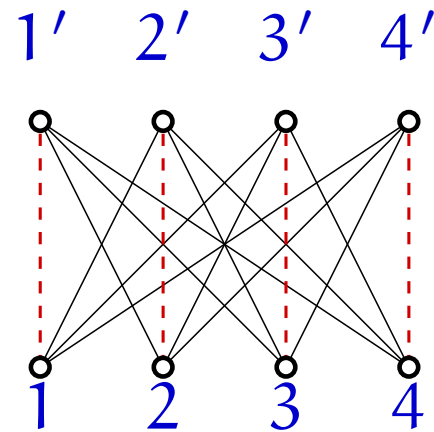
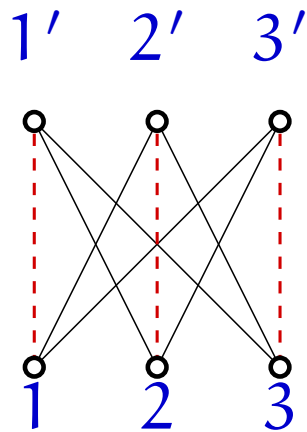


- Dimension 3: Containment orders of triangles.



Standard Examples

- Standard example of an n dimensional order:



Orders and Planarity

- If an order \mathcal{P} has $\mathbf{0}$ and $\mathbf{1}$ and a planar diagram, then $\dim(\mathcal{P}) = 2$.



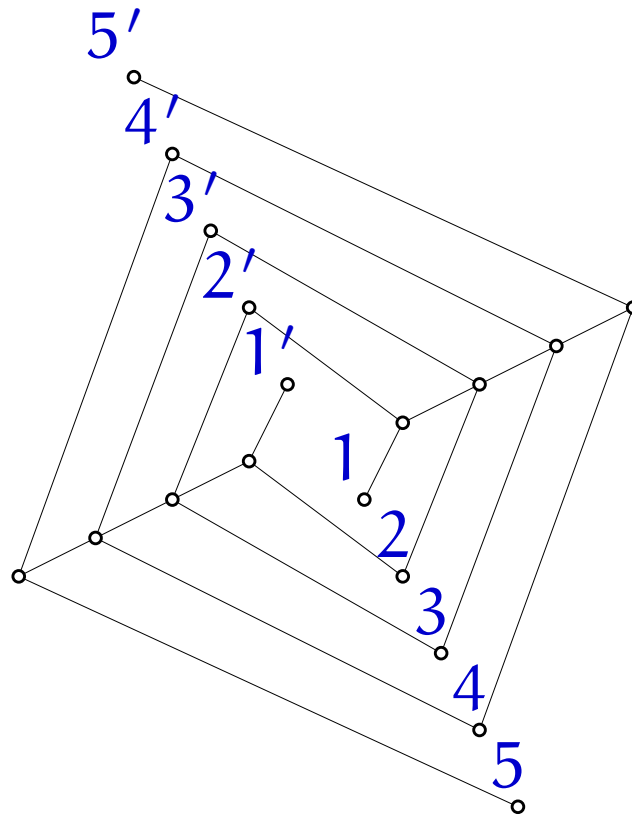
- If an order \mathcal{P} has $\mathbf{0}$ and a planar diagram, then $\dim(\mathcal{P}) \leq 3$.



- The dimension of an order \mathcal{P} with a a planar diagram can be unbounded.

Plane Standard Example

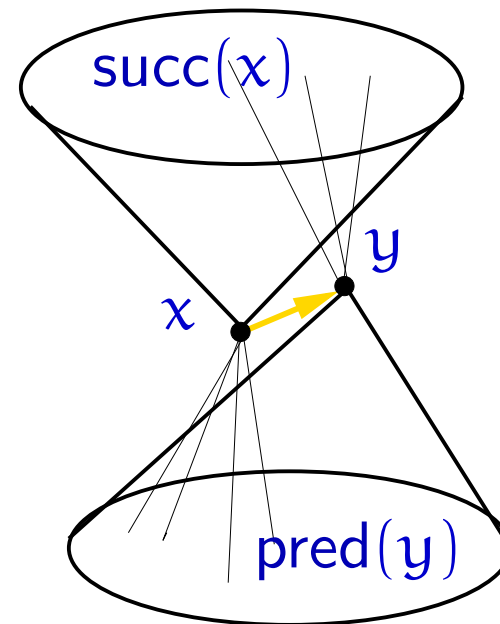
The dimension of an order \mathcal{P} with a planar diagram can be unbounded:



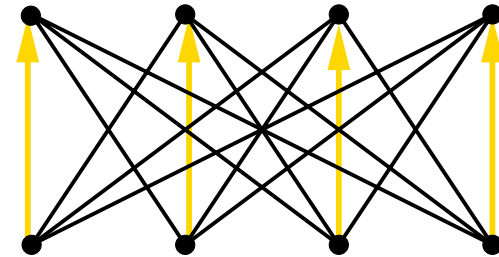
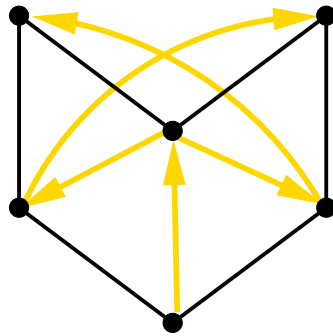
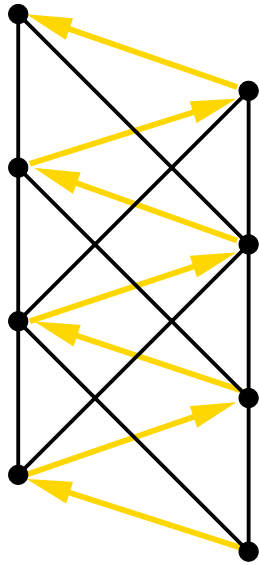
Critical Pairs

Definition. An incomparable pair (x, y) is **critical** if

- $a < x$ implies $a < y$.
- $y < b$ implies $x < b$.



Critical Pairs



Proposition. A family \mathcal{R} of linear extensions of \mathcal{P} is a realizer of \mathcal{P} \iff \mathcal{R} reverses all critical pairs.

The Hypergraph of Critical Pairs

Definition. The **hypergraph of critical pairs** H_P has the set of critical pairs as vertices. A set of critical pairs is an edge if it can't be reversed by a single linear extension and is minimal.

Definition. The **graph of critical pairs** G_P has the set of critical pairs as vertices. A set of two critical pairs is an edge if there is no linear extension which reverses both pairs.

Dimension and Chromatic Number

For every order: $\dim(\mathbf{P}) = \chi(H_{\mathbf{P}}) \geq \chi(G_{\mathbf{P}})$.

Theorem [Doignon, Ducamp and Falmagne 1984, F. and Trotter 2000].

$$\dim(\mathbf{P}) = 2 \iff \chi(G_{\mathbf{P}}) = 2.$$

■

Conjecture. For every $t \geq 3$, there exists an order \mathbf{P} with

$$\chi(G_{\mathbf{P}}) = 3 \quad \text{and} \quad \dim(\mathbf{P}) = t.$$

Important Examples

- Boolean lattice: $\dim(\mathbf{B}_n) = n$



- Standard example of an n dimensional order:
Atoms and coatoms of a Boolean lattice:

$$\dim(\mathbf{B}_n[1, n-1]) = n$$



- First two levels of a Boolean lattice:

$$\dim(\mathbf{B}_n[1, 2]) \sim \log \log n$$

Challenge: Compute $\dim(\mathbf{B}_{100}[1, 2])$.

Important Examples

- Boolean lattice: $\dim(\mathbf{B}_n) = n$
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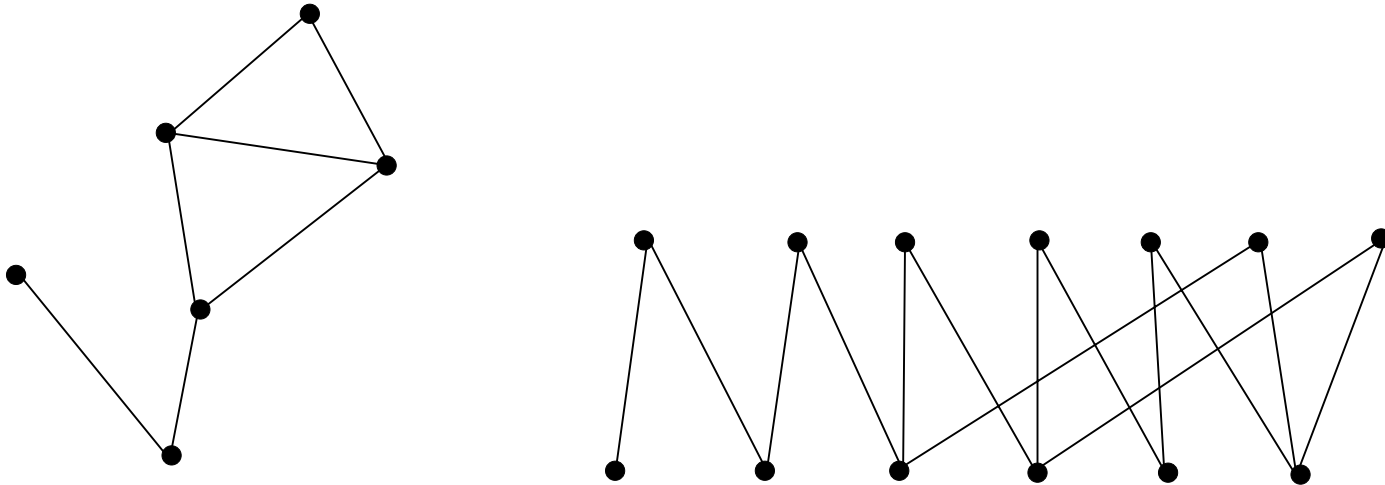
- First two levels of a Boolean lattice:

$$\dim(\mathbf{B}_n[1, 2]) \sim \log \log n$$

$\dim(K_n) \leq$	2	3	4	5	6	7	8
$n \leq$	2	4	12	81	2646	1422564	229809982112

Incidence Orders and Dimension

The incidence order P_G of G



Theorem [Schnyder 1989].

A Graph G is planar $\iff \dim(P_G) \leq 3$.

The Lower Bound

- $\dim(G) \leq 3 \implies G$ planar.

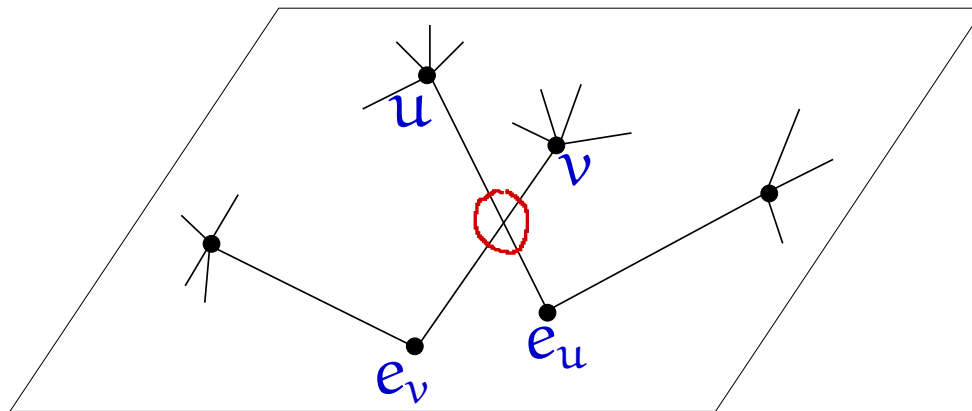
Babai and Duffus 1982

P_G in \mathbb{R}^3



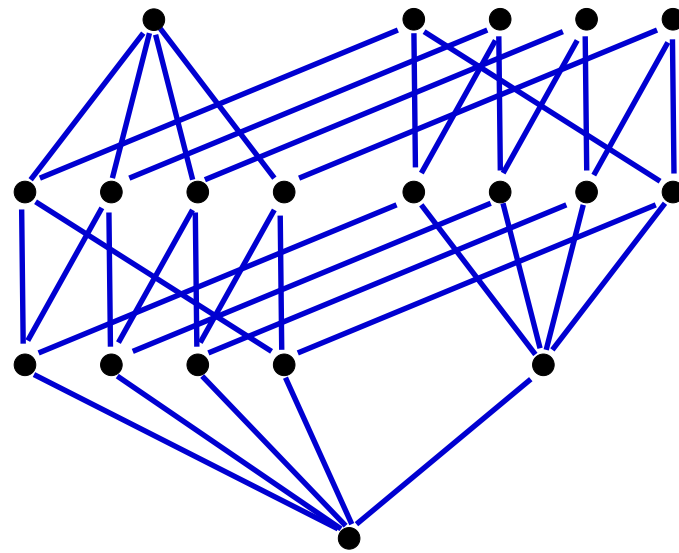
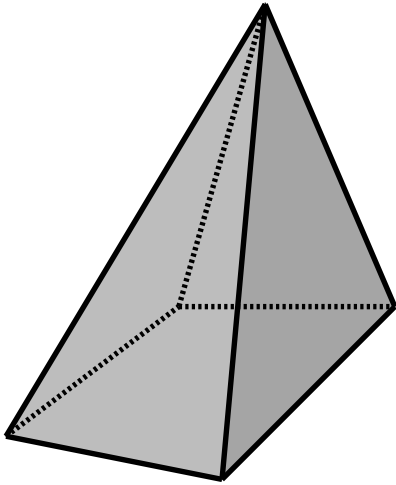
projection

$$x + y + z = 1$$



Dimension of Polytopes

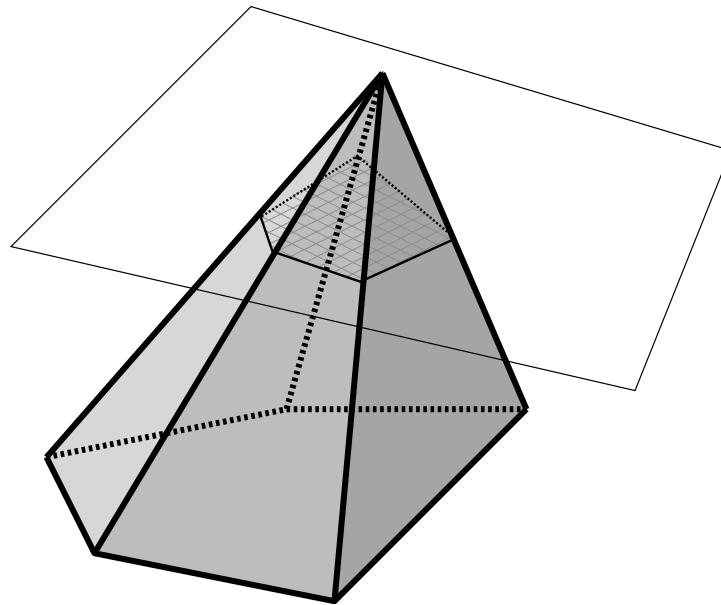
Let \mathcal{F}_P be the face lattice of polytope P .



Dimension of Polytopes: Lower Bound

Theorem [Reuter 1990].

If P is a d -polytope, then $\dim(\mathcal{F}_P) \geq d + 1$.

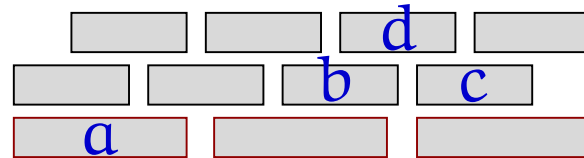
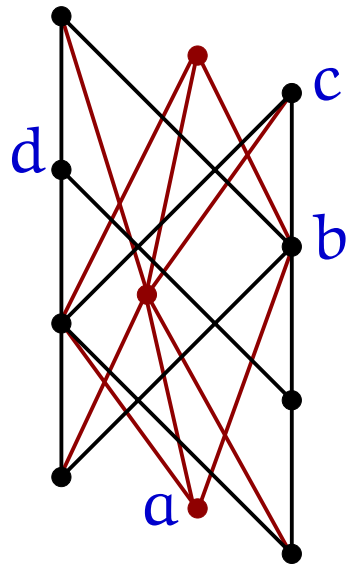


Interval Orders

Definition. An Order $P = (X, <_P)$ is an **interval order** if there is a family $(I_x)_{x \in X}$ of intervals on \mathbb{R} such that

$$x <_P y \iff I_x \text{ is left of } I_y.$$

Example.



Interval Dimension

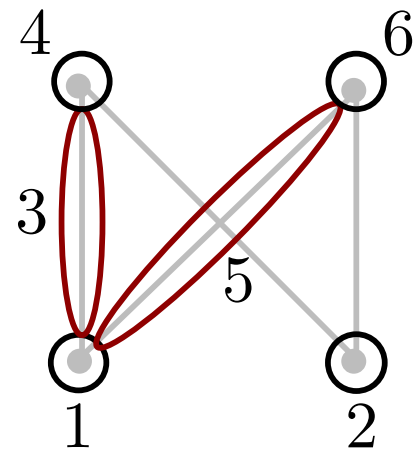
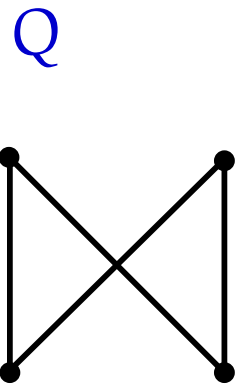
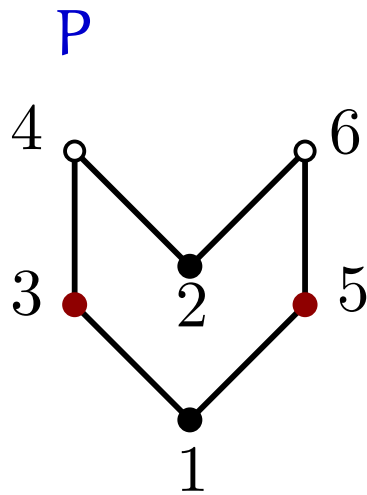
A family \mathcal{I} of interval extensions of \mathbf{P} is an **interval realizer** for $\mathbf{P} = (X, <)$ provided that

- * for every incomparable pair (x, y) there is an $I \in \mathcal{I}$ such that $x \not\prec y$ in I .

The **interval dimension**, $\dim_I(\mathbf{P})$, of \mathbf{P} is the minimum t , such that there is an interval realizer $\mathcal{I} = \{I_1, I_2, \dots, I_t\}$ for \mathbf{P} of size t .

Interval Dimension

- $\dim_I(\mathcal{P}) = \min(\mathfrak{t}$ such that there is an interval realization of \mathcal{P} on some Q with $\dim(Q) = \mathfrak{t}$).



Comparing Notions of Dimension

\mathcal{P} a bipartite order, then

- $\dim_I(\mathcal{P}) \leq \dim(\mathcal{P}) \leq \dim_I(\mathcal{P}) + 1.$

G a graph, then

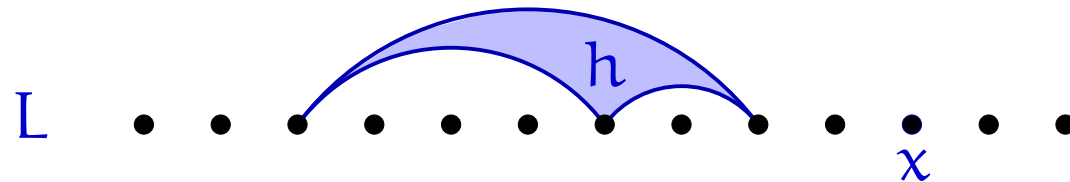
- $\dim(G) = \dim_I(P_G).$

If $\delta(G) > 1$, then $\dim(G) = \dim_I(P_G) = \dim(P_G).$

Dimension of Hypergraphs

A family Γ of permutations of V is a **realizer** for $H = (V, E)$ provided that

- * for every edge h and every $x \in V - h$ there is an $L \in \Gamma$ such that $x > h$ in L .



The **dimension**, $\dim(H)$, of H is the minimum t , such that there is a realizer $\Gamma = \{L_1, L_2, \dots, L_t\}$ for H of size t .

Dimension of Polytopes

P a d -polytope.

H_P the hypergraph of faces and \mathcal{F}_P the face lattice.

- $\dim(H_P) = \dim(\mathcal{F}_P) \geq d + 1$.

Theorem [Schnyder 1989].

If P is a simplicial 3-polytope and F is any face of P , then

- $\dim(\mathcal{F}_P \setminus F) = 3$
- $\dim(\mathcal{F}_P) = 4$

■

Theorem [Brightwell and Trotter 1992].

If P is a 3-polytope and F is any face of P , then

- $\dim(\mathcal{F}_P \setminus F) = 3$
- $\dim(\mathcal{F}_P) = 4$

Dimension and Topological Graphs

- $\dim(G) \leq 2 \iff G$ is a caterpillar.
- $\dim(G) \leq 3 \iff G$ is planar.
- **Challenge:** Find something meaningful for ≤ 4 .

Dimension and Topological Graphs

- $\dim(G) \leq 1\frac{1}{2} \iff G$ is a path.
- $\dim(G) \leq 2 \iff G$ is a caterpillar.
- $\dim(G) \leq 2\frac{1}{2} \iff G$ is outerplanar.
- $\dim(G) \leq 3 \iff G$ is planar.
- **Challenge:** Find something meaningful for $\leq 3\frac{1}{2}$.
- **Challenge:** Find something meaningful for ≤ 4 .

Problems

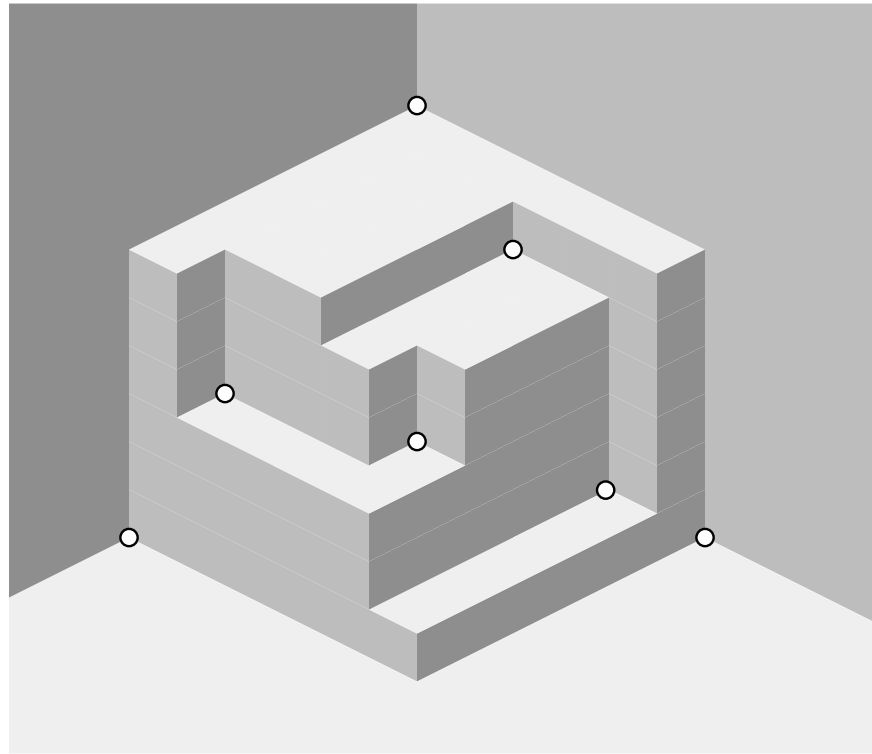
- Graphs of dimension at most 4.
- Dimension of polytopes.
- Given a bipartite order P . Can $\dim(P) \leq 3$ be decided in polynomial time?

Modules

Dimension

Orthogonal Surfaces and Dimension

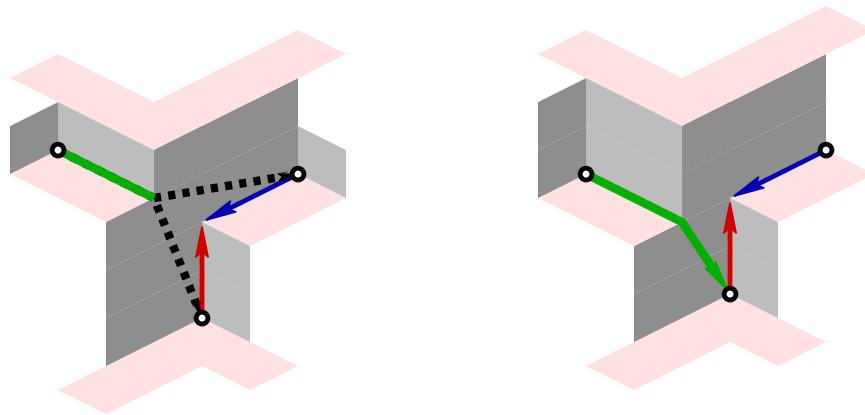
Orthogonal Surfaces



Orthogonal Surfaces and Dimension

A surface S_X is **rigid** iff

$e = (x, y) \in G_S$ and $z \neq x, y$ implies $z \not\leq e = x \vee y$.



Remark. A rigid orthogonal surface supports a unique Schnyder wood of a 3-connected planar graph.

Orthogonal Surfaces and Dimension

Theorem [Miller '02].

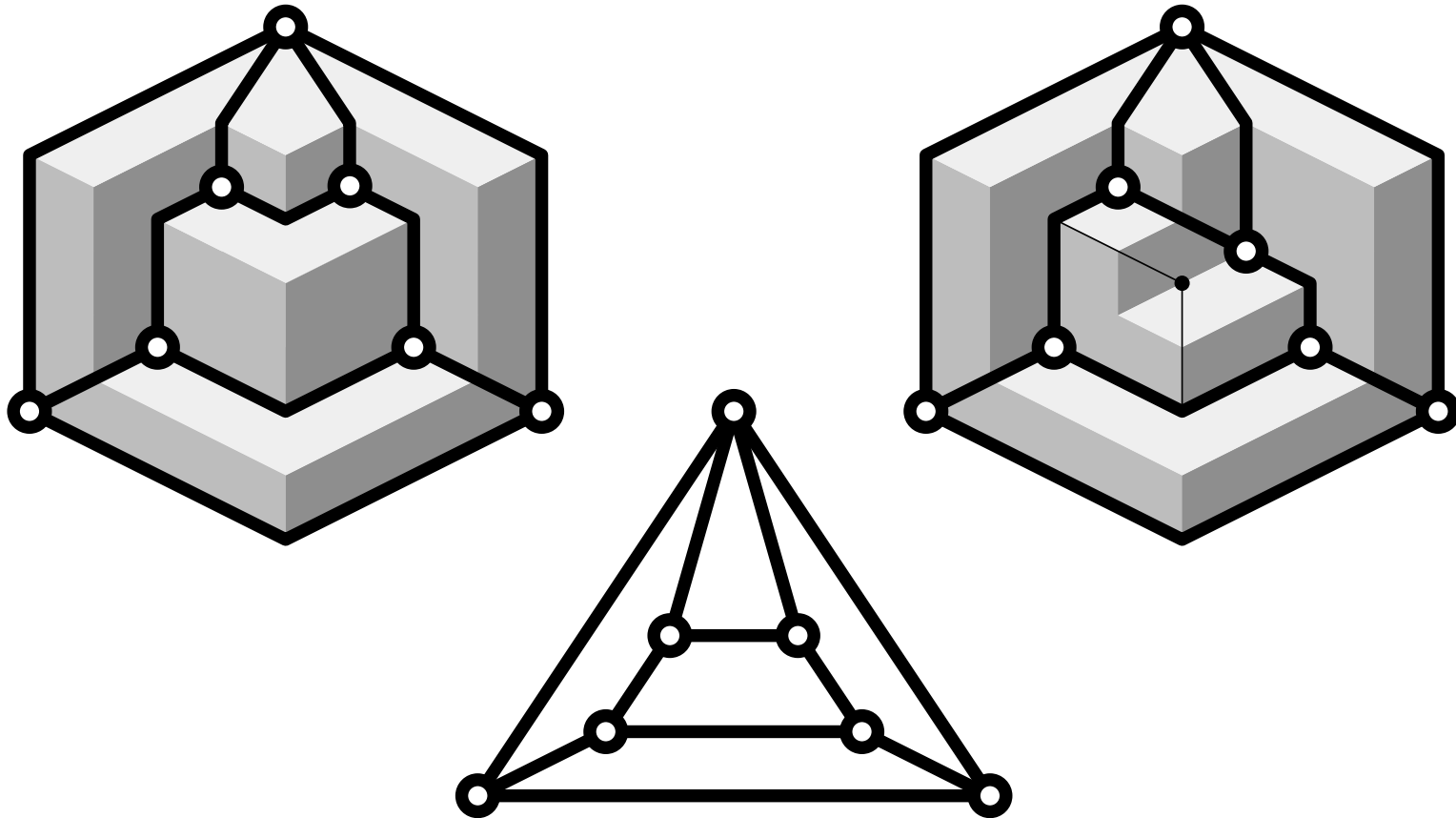
If G is the graph of a rigid orthogonal surface and P a 3-polytope with graph G , then $\dim(\mathcal{F}_P \setminus F_\infty) = 3$.

Theorem [Brightwell+Trotter '92].

If P is a 3-polytope and F is any face of P , then

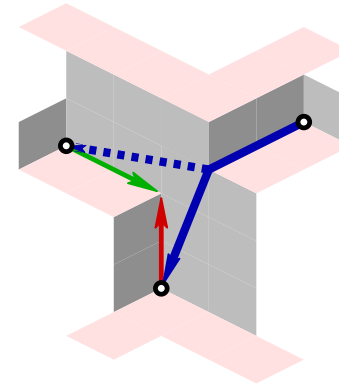
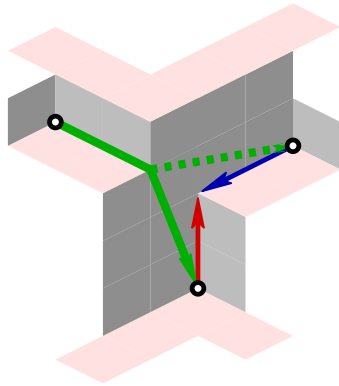
- $\dim(\mathcal{F}_P \setminus F) = 3$
- $\dim(\mathcal{F}_P) = 4$

Rigidity and the BT Theorem

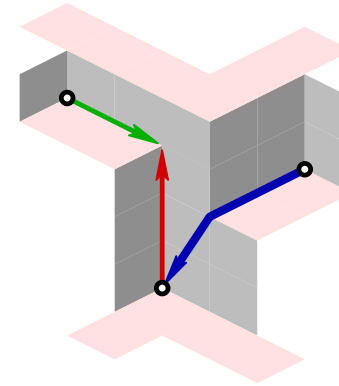
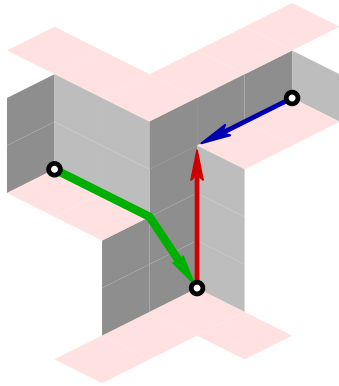


A New Approach *

Two types of non-rigid edges



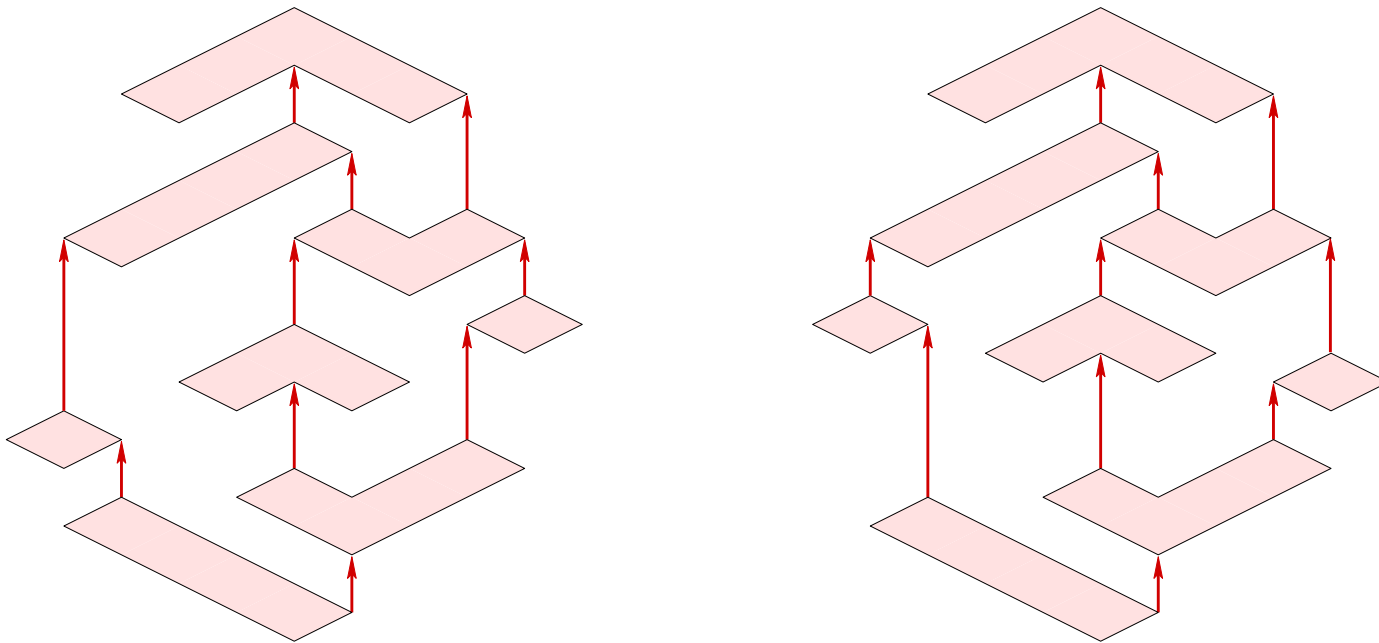
resolved by shifting red flats.



* Joint work with Florian Zickfeld.

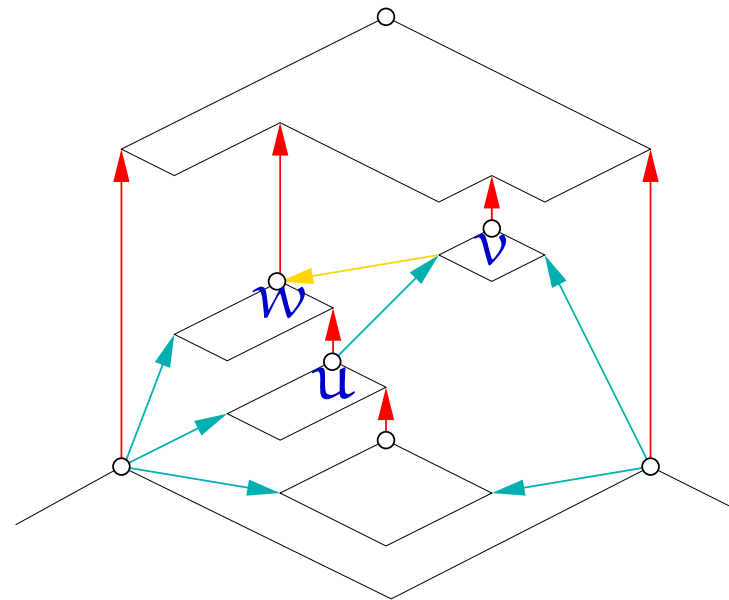
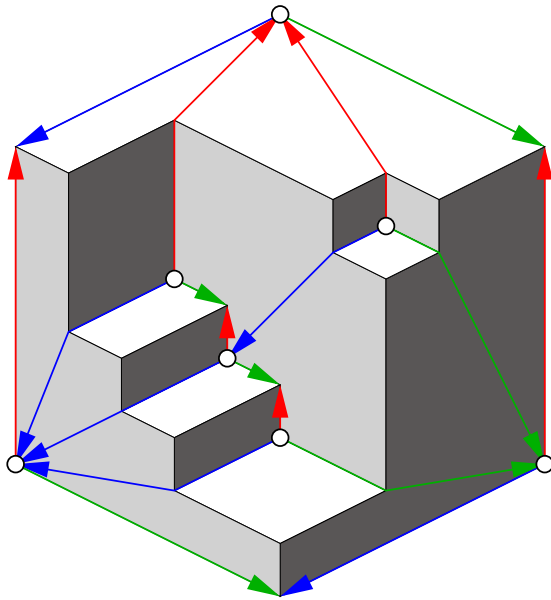
Shifting Flats

Red flats form a 2-dimensional order \mathcal{P}_r . Compatible height functions of red flats correspond to order preserving maps $\phi : \mathcal{P}_r \rightarrow \mathbb{R}$.



Shifting Flats

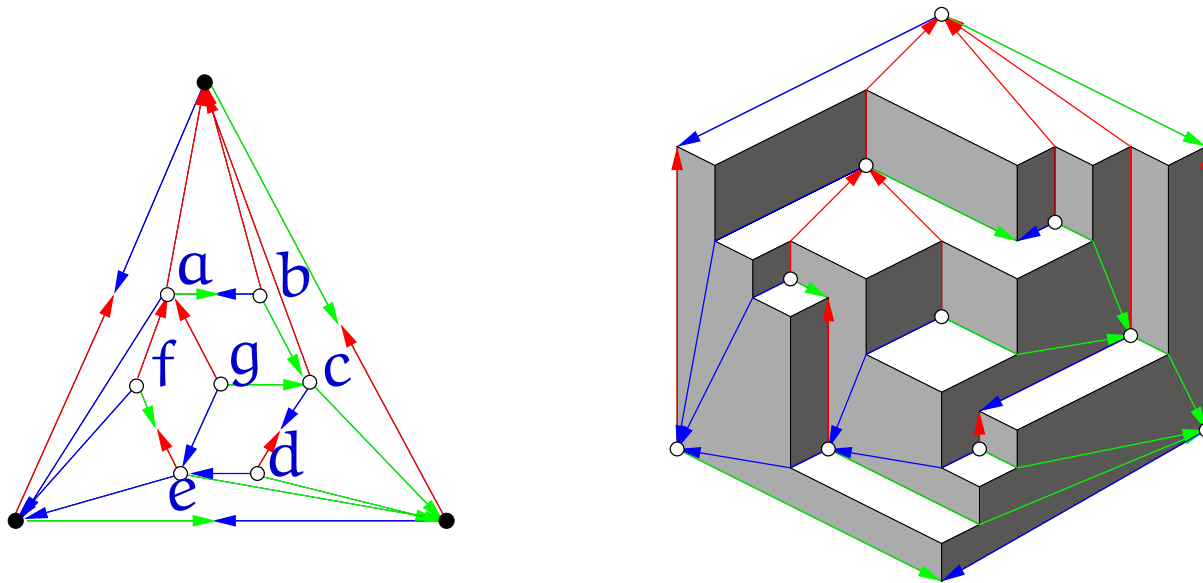
Split bidirected edges to prevent non-rigid. This doesn't spoil acyclicity.



Theorem. There are compatible assignments for the heights of flats of all three colors which together yield a rigid orthogonal surface supporting the original Schnyder wood.

Rigidity and Coplanarity

A Schnyder Wood on a rigid, but not coplanar surface



Coplanarity means that $v_1 + v_2 + v_3 = c = w_1 + w_2 + w_3$ for all $v, w \in V$, hence, $v_i = w_i$ implies $v_{i-1} - w_{i-1} = w_{i+1} - v_{i+1}$. In the given instance rigidity requires $f_1 > g_1$, $b_2 > g_2$ and $d_3 > g_3$.

THE END



Thank you.